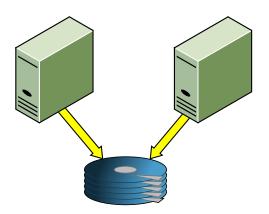
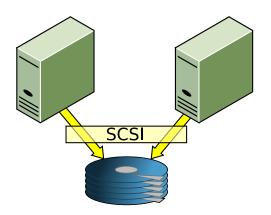
Decentralized Deduplication in SAN Cluster File Systems

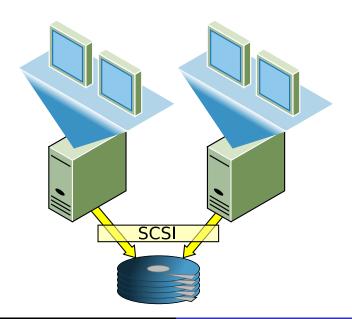
Austin T. Clements*

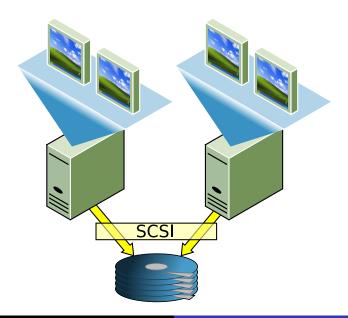
Irfan Ahmad Murali Vilayannur Jinyuan Li

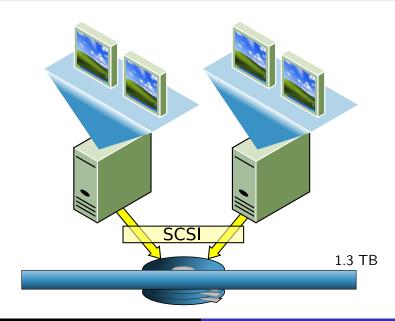
VMware, Inc. *MIT CSAIL

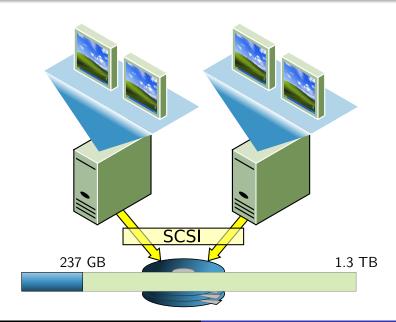


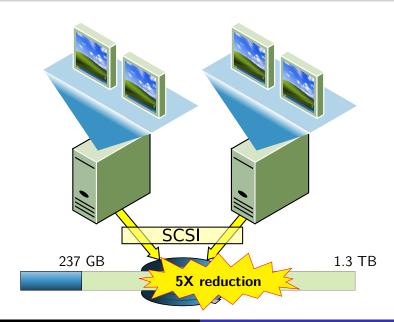


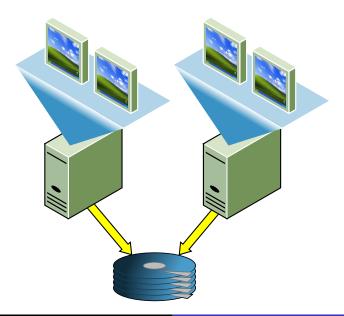


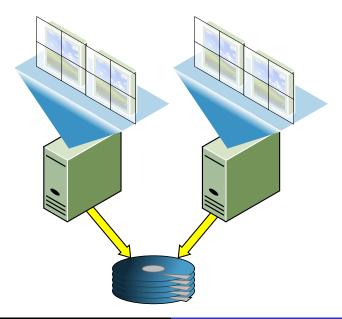


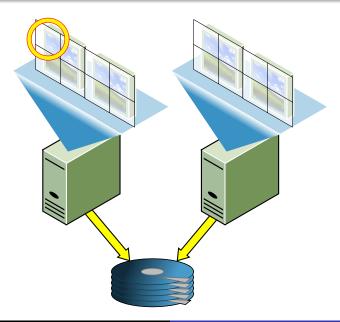


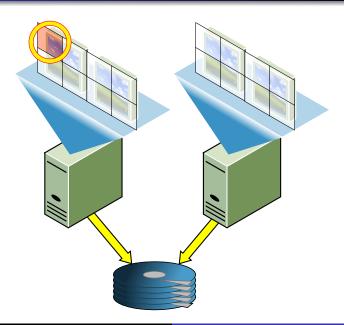


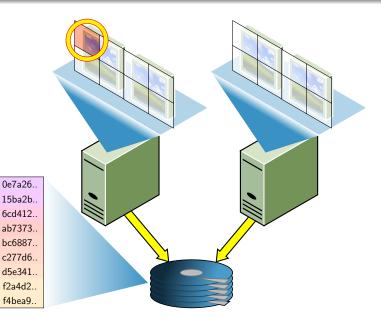


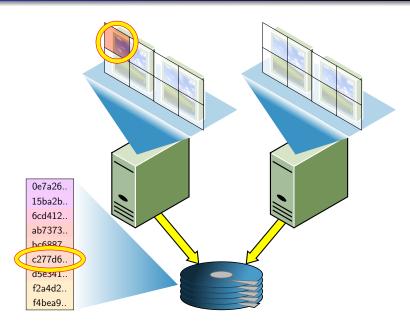


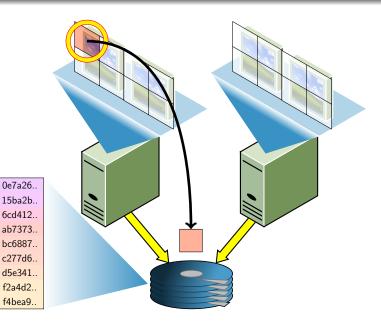


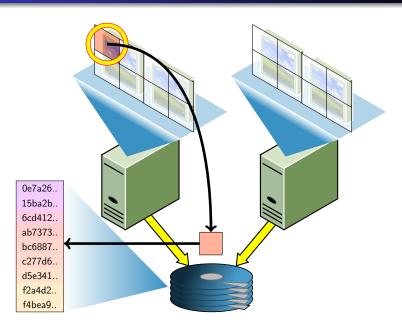












The Problem

Slow — Lots of IO in the write path. Can't cache the index.

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Very Slow — Writes require allocation and thus coordination. No hope of disk locality.

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Hopelessly Slow — Multi-host lock contention on shared index.

Decentralized Deduplication

DeDe

DeDe

Write Monitoring Local Deduplication

DeDe

Write Monitoring Local Deduplication

- Out-of-band deduplication of live, primary storage
- Process duplicates efficiently, in large batches
- Minimize contention on the index
- Resilient to stale index information
- Unique blocks remain mutable and sequential
- → No overhead for blocks that don't benefit from deduplication

DeDe

Write Monitoring Local Deduplication





Write Monitoring Local Deduplication Cross-Host Deduplication



Write Monitoring Local Deduplication



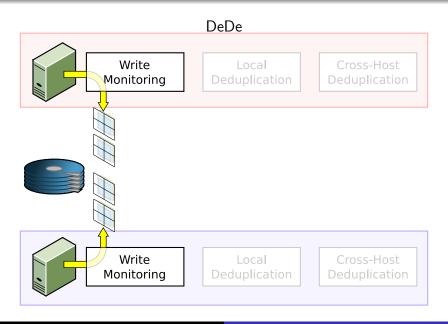


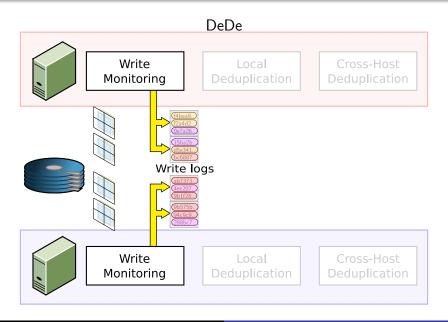
Write Monitoring Local Deduplication Cross-Host Deduplication

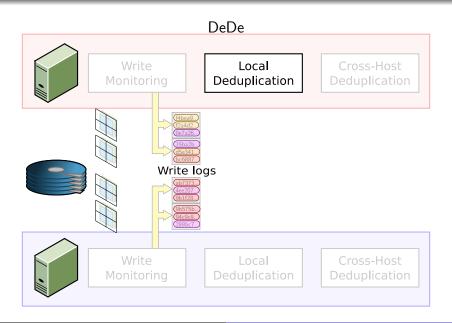


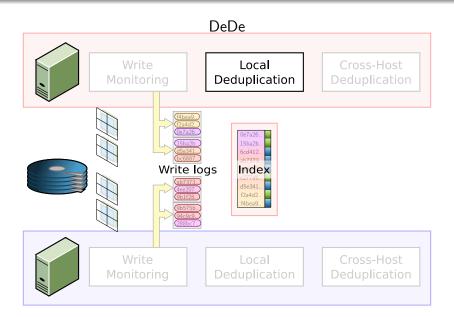


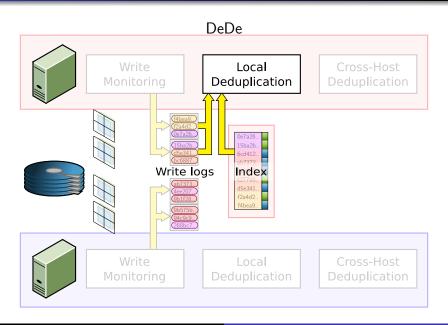
Write Monitoring Local Deduplication

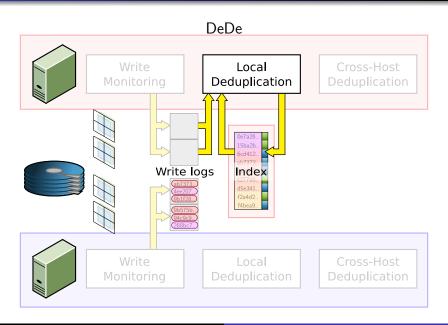


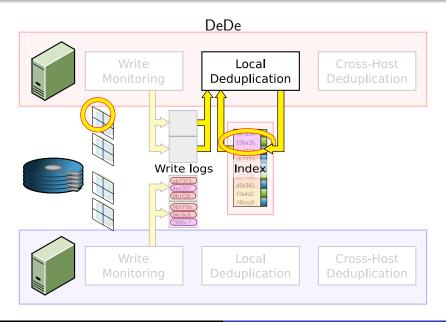


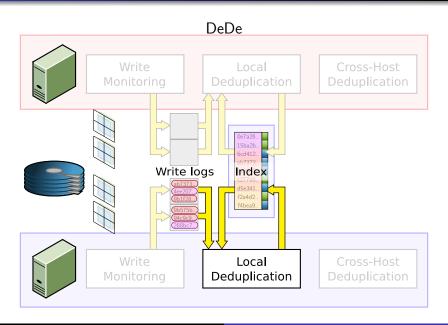


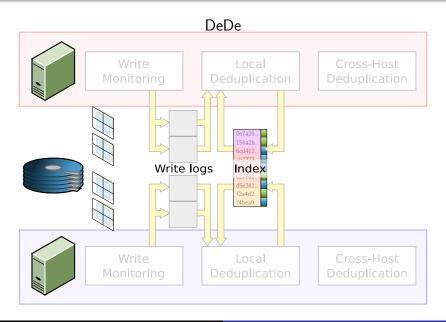


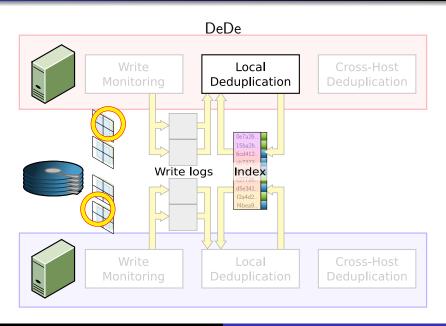




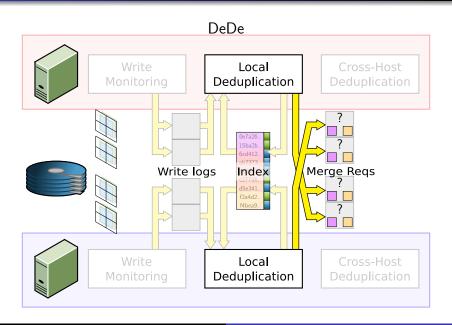




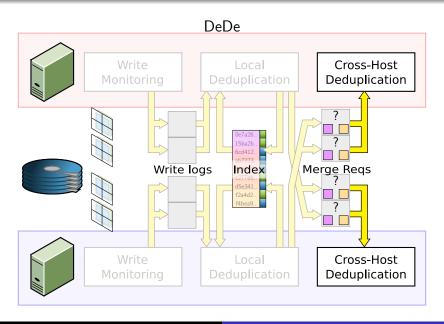




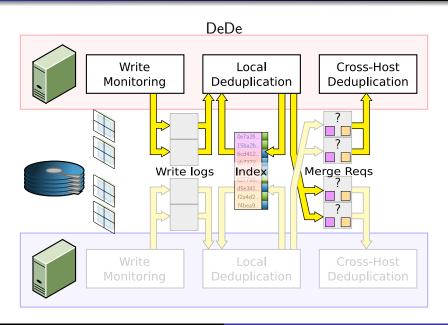
Three-Stage Deduplication



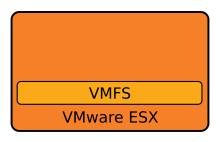
Three-Stage Deduplication

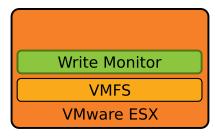


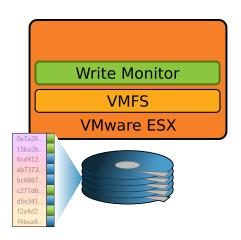
Three-Stage Deduplication

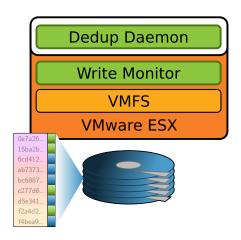




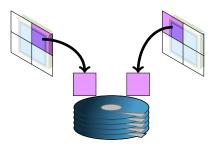




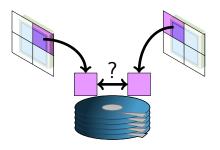




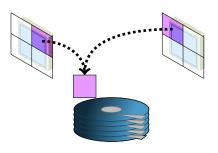
Compare-and-share



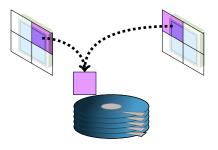
Compare-and-share



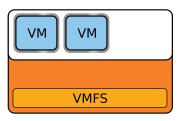
Compare-and-share



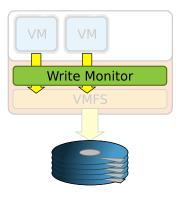
Compare-and-share



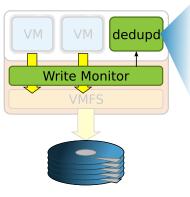
DeDe finds duplicates. VMFS eliminates them.





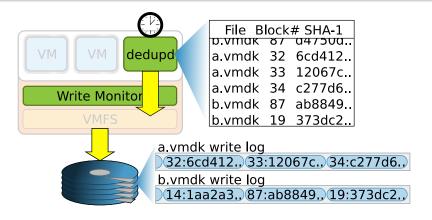


• A lightweight kernel module monitors writes, computes hashes

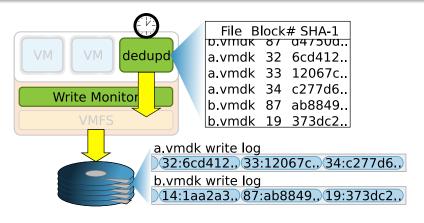


File Block# SHA-1 b.vmak 87 u4750a.. a.vmdk 32 6cd412.. a.vmdk 33 12067c.. a.vmdk 34 c277d6.. b.vmdk 87 ab8849.. b.vmdk 19 373dc2..

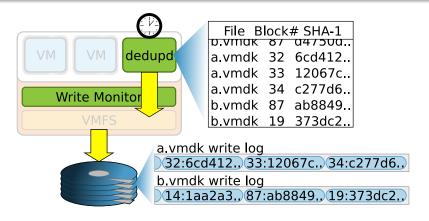
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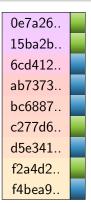
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- ullet 150 MB of regular writes ightarrow 1 MB sequential log write

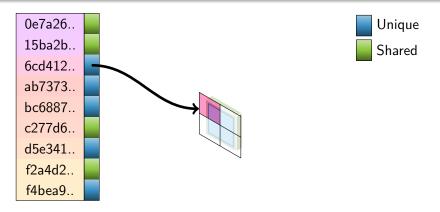
```
0e7a26..
15ba2b..
6cd412...
ab7373...
bc6887...
c277d6..
d5e341..
f2a4d2...
f4bea9..
```

• Map from hashes to block locators, list sorted by hash

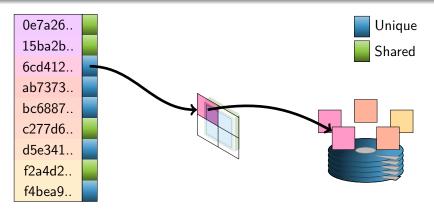




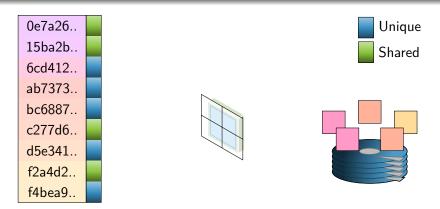
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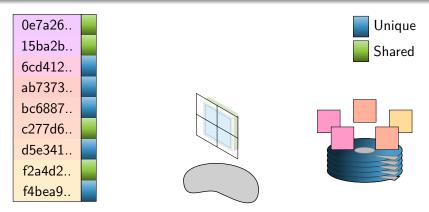
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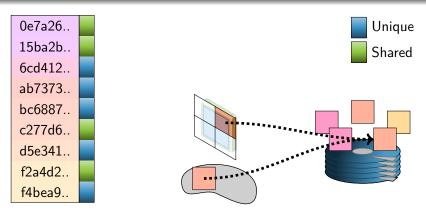
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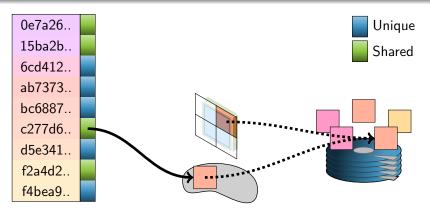
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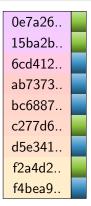
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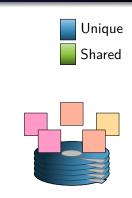


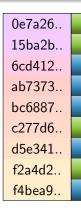
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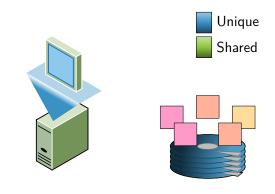


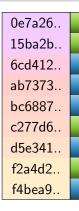
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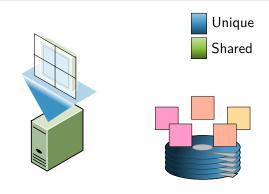


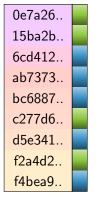


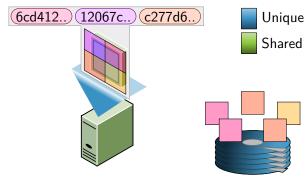


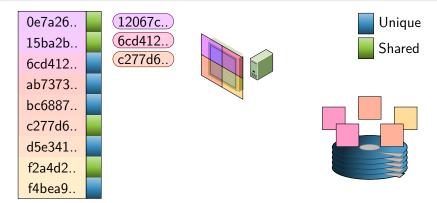


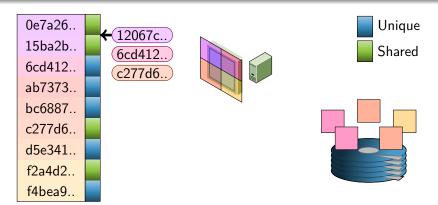




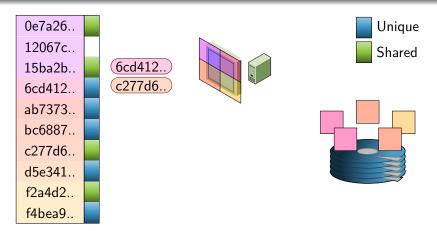


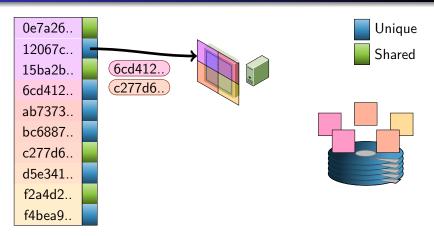


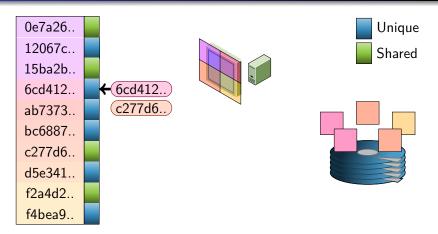


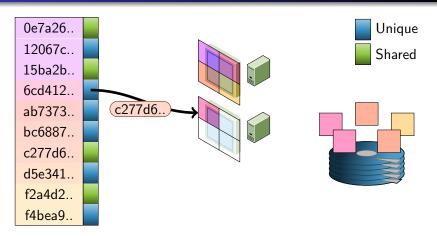


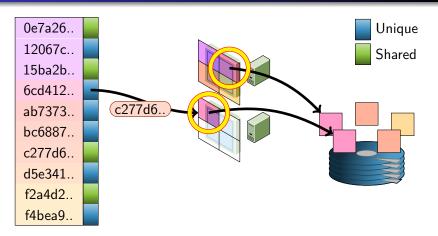


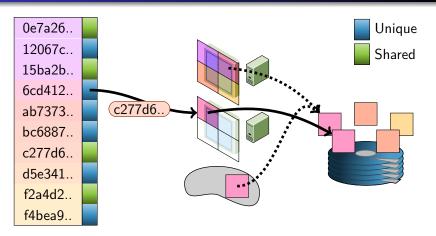


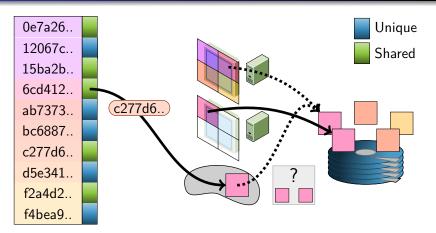


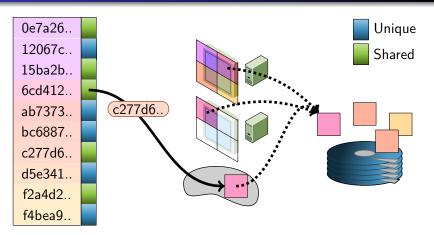


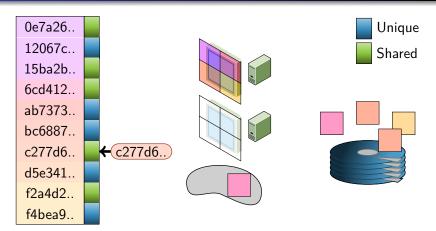


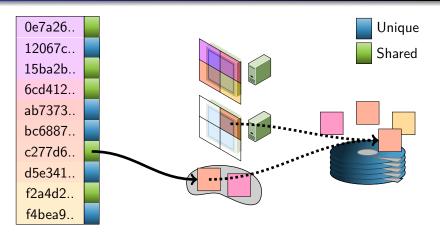


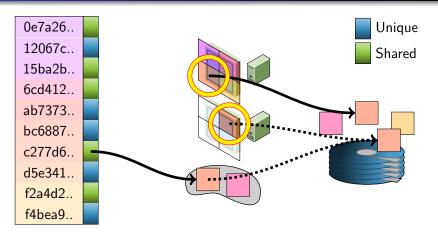


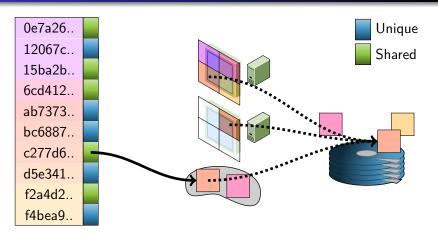


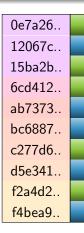


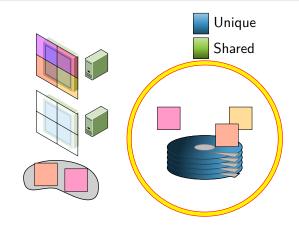












Phew.



Evaluation

How much space does DeDe save?

Evaluation

How much space does DeDe save?

How much overhead does DeDe introduce?

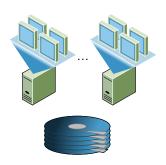
Evaluation

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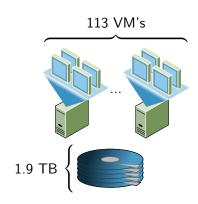
How much overhead does DeDe introduce?

How fast can DeDe deduplicate?

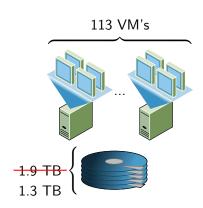
- Corporate Virtual Desktop Infrastructure cluster
- Desktop XP VM's
- 6–12 months of active use
- Originally cloned from small number of base images



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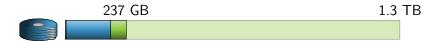


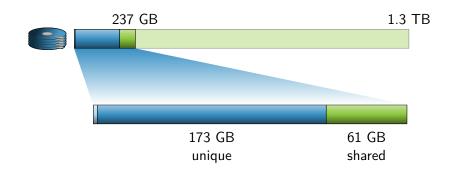
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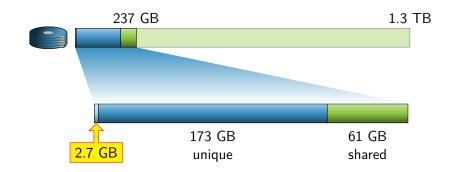


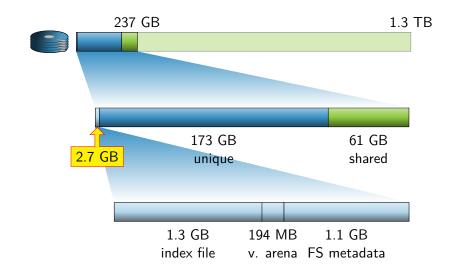
1.3 TB









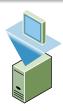


Runtime Effects

- Write monitoring
- Disk array caching

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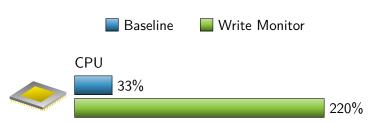
- Write monitoring
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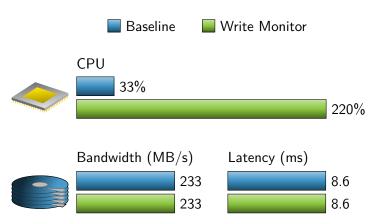
EMC CLARiiON CX3-40

Worst-case benchmark: 100% sequential write IO, No computation

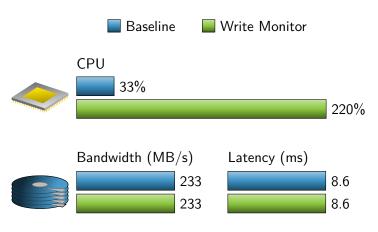
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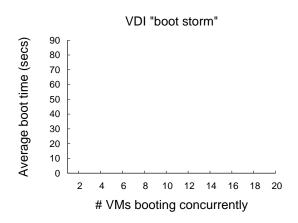
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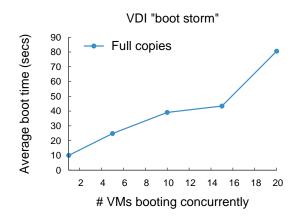


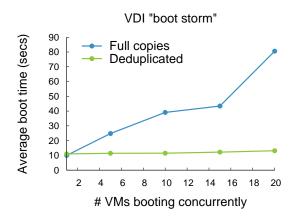
Worst-case benchmark: 100% sequential write IO, No computation



(See paper for database application benchmark)







Index scan
COW sharing

Index scan
COW sharing

→ Virtually no cost for unique blocks

6.6 GB/sec

 $\begin{array}{ll} \text{Index scan} & 6.6 \text{ GB/sec} \\ \text{COW sharing} & 2.6 \text{ MB/sec} \end{array}$

→ Virtually no cost for unique blocks

```
Index scan 6.6 GB/sec
COW sharing 2.6 MB/sec
(It's a prototype!)
```

⇒ Virtually no cost for unique blocks

```
Index scan 6.6 GB/sec
COW sharing 2.6 MB/sec
(It's a prototype!)
```

→ Virtually no cost for unique blocks
 → 9 GB of new shared blocks per hour
 (And provisioning can be special-cased)

Related Work

- Centralized archival
 - Venti
 - Data Domain
 - Foundation
- Centralized primary storage
 - NetApp ASIS
 - Microsoft Single Instance Store
- Distributed
 - Farsite
- SAN with Coordinator
 - DDE

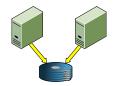
Decentralized, out-of-band, live file system deduplication.

Decentralized, out-of-band, live file system deduplication.

Deduplication is effective.

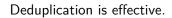
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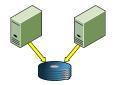
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Deduplication is hard.

Decentralized, out-of-band, live file system deduplication.



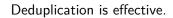


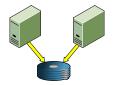
Deduplication is hard.

Write Local Cross-Host Deduplication

Three-stage deduplication has only modest performance overhead.

Decentralized, out-of-band, live file system deduplication.





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Three-stage deduplication has only modest performance overhead.

Thank you.