



Understanding latent sector errors and how to protect against them

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Motivation

- ▶ What is a latent sector error (LSE)?
 - ▶ Individual sectors on a drive become inaccessible (media error)
 - ▶ Prevalence?
 - ▶ 3.5% of drives experience LSE(s) [Bairavasundaram2007]
 - ▶ 7-9% for some disk models!
 - ▶ Consequence of an LSE?
 - ▶ In a system without redundancy: **data loss** 
 - ▶ In RAID-5, if discovered during reconstruction: **data loss** 
 - ▶ One of the main motivations for RAID-6
 - ▶ Growing concern with growing disk capacities
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How to protect against them?

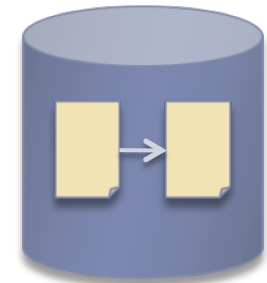
- ▶ **Periodic scrubbing**

- ▶ Proactively detect LSEs and correct them.



- ▶ **Intra-disk redundancy**

- ▶ Replicate selected metadata [e.g. FFS]
- ▶ Add parity block per file [e.g. Iron file systems]
- ▶ Add parity block per group of sectors [Dholak.08]

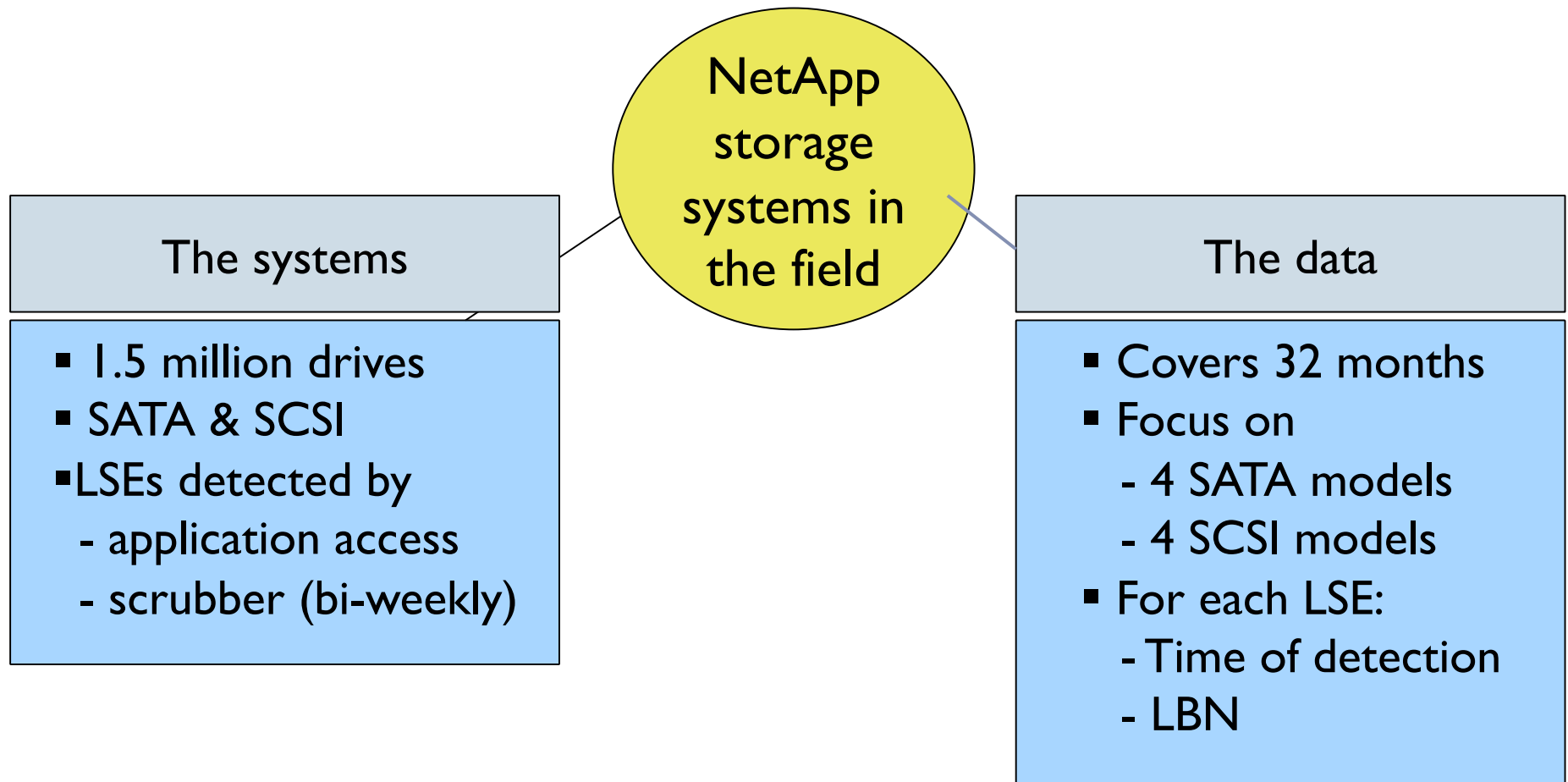


Our goal

- ▶ Understand potential of different protection schemes
 - ▶ Understand characteristics of LSEs
 - ▶ From point of view of protection

 - ▶ **How?**
 - ▶ Using real data from production machines
 - ▶ Subset of data in Bairavasundaram et al. (Sigmetrics'07)
 - ▶ **Thanks for sharing!**
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The data



How effective are protection schemes?

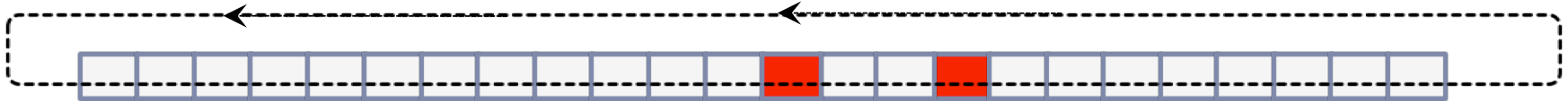
- ▶ Scrubbing
- ▶ Intra-disk redundancy

Scrubbing

- ▶ Why?
 - ▶ Detect and correct errors early
 - ▶ Reduces probability to encounter LSE during RAID reconstruction
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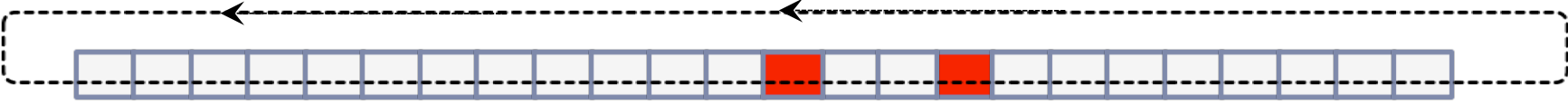
Scrubbing

- ▶ Standard sequential scrubbing

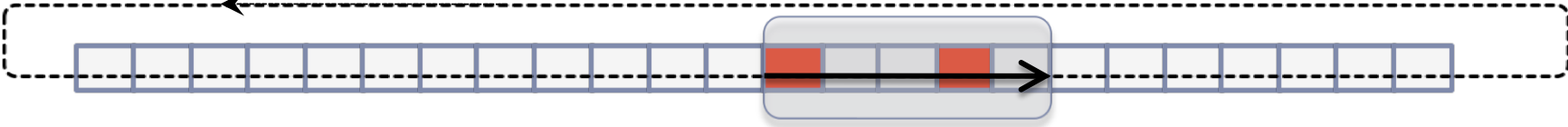


Scrubbing

▶ Standard sequential scrubbing

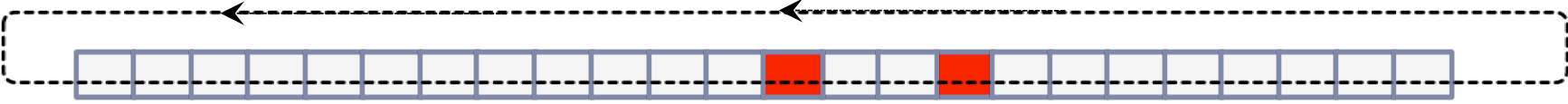


▶ Localized scrubbing

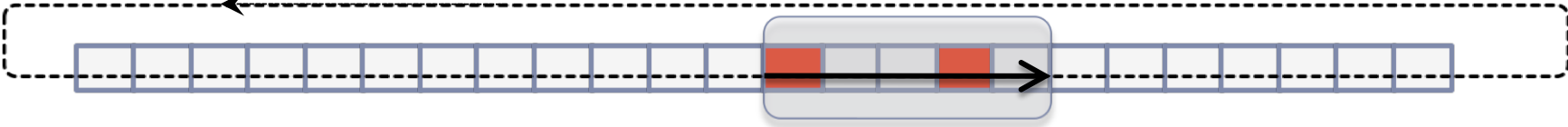


Scrubbing

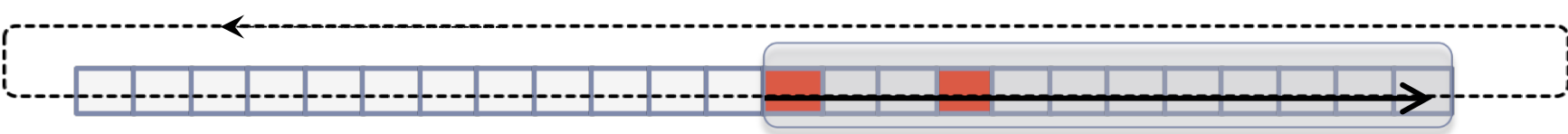
▶ Standard sequential scrubbing



▶ Localized scrubbing

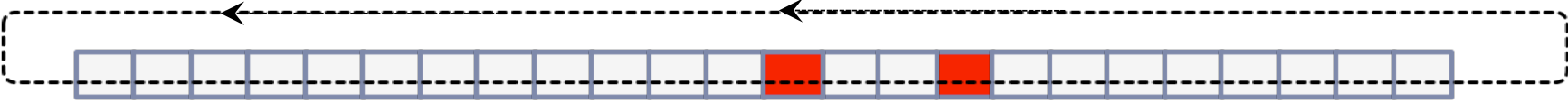


▶ Accelerated scrubbing

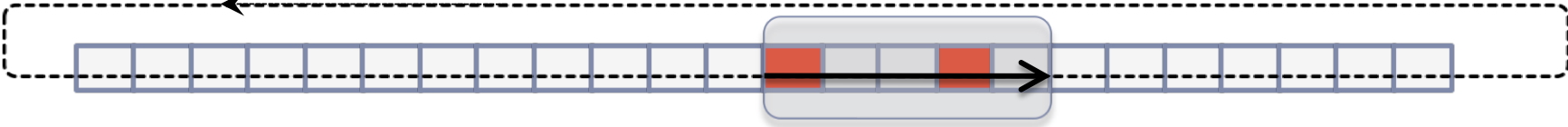


Scrubbing

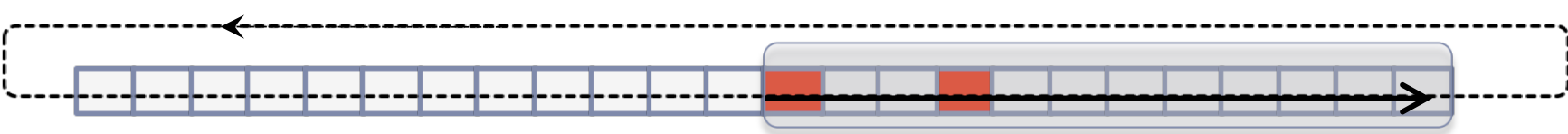
- ▶ Standard sequential scrubbing



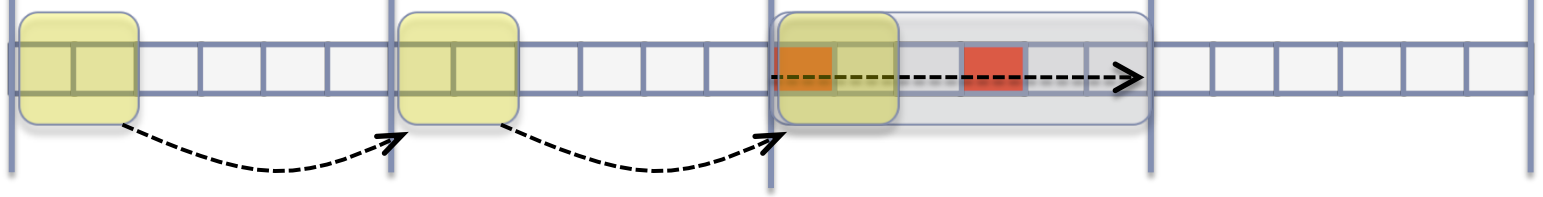
- ▶ Localized scrubbing



- ▶ Accelerated scrubbing



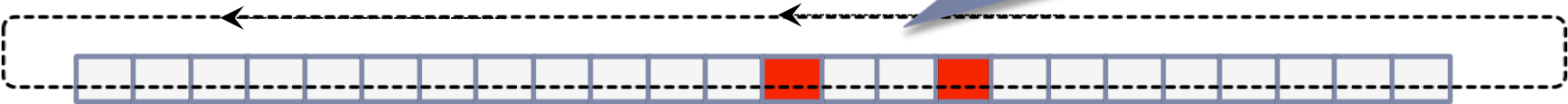
- ▶ Staggered scrubbing [Oprea et al.'10]



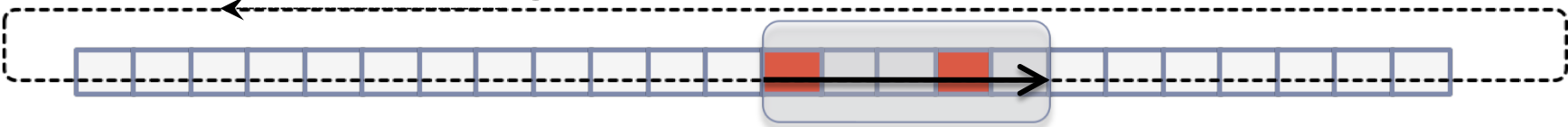
Scrubbing

How do those approaches perform in practice, i.e. on real-world data?

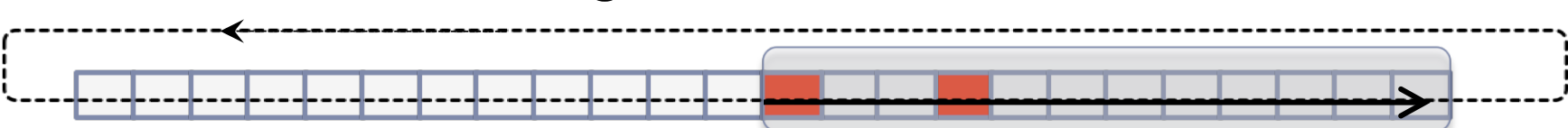
- ▶ Standard sequential scrubbing



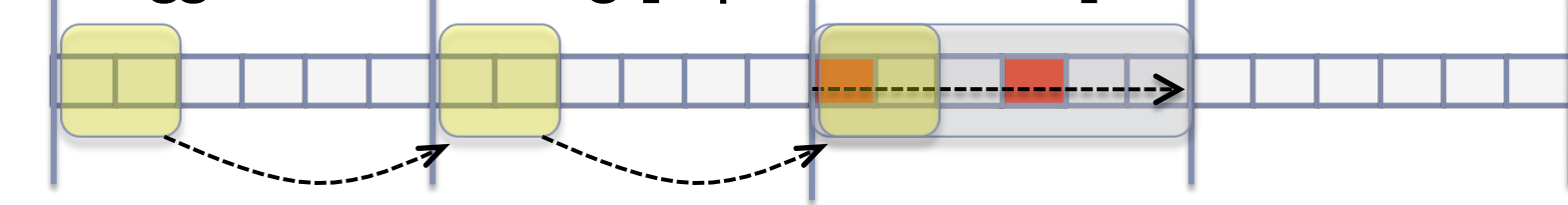
- ▶ Localized scrubbing



- ▶ Accelerated scrubbing

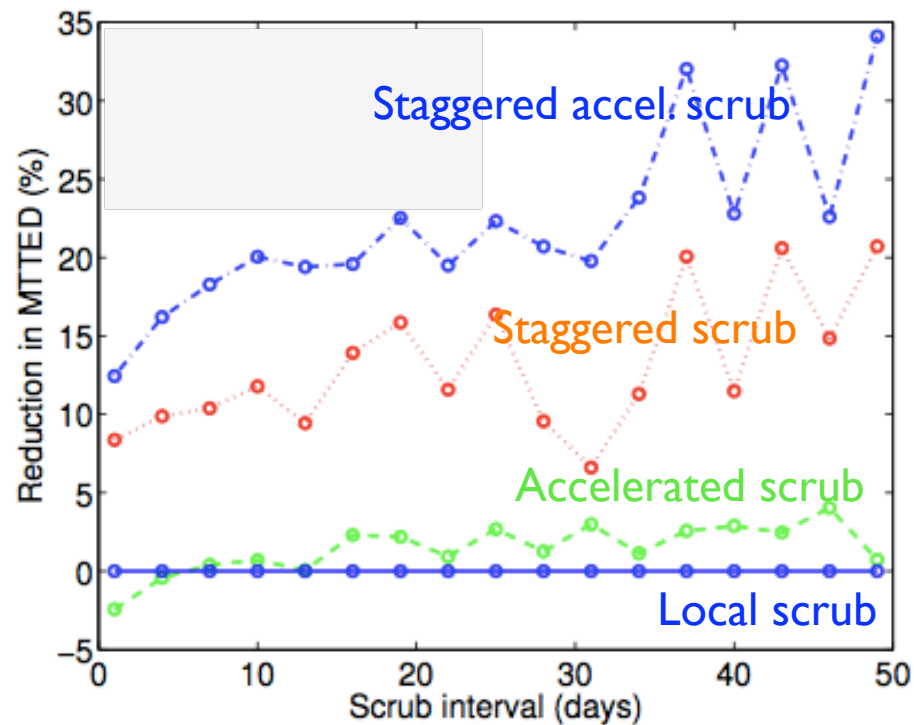


- ▶ Staggered scrubbing [Oprea et al.'10]



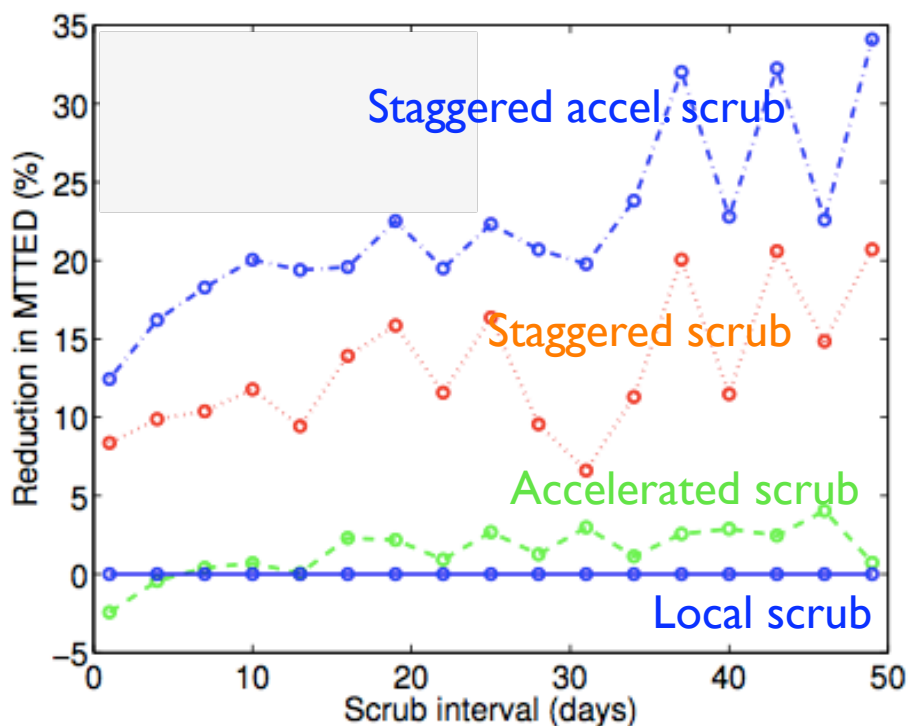
- ▶ Accelerated staggered scrubbing

Scrubbing: Evaluation on NetApp data



- ▶ No significant improvement from local & accelerated scrubs
 - ▶ They don't reduce the time to detect whether there are any errors
 - ▶ Errors are close in space, so even standard scrub finds them soon

Scrubbing: Evaluation on NetApp data



- ▶ 10-35% improvement with staggered scrubs!
 - ▶ Even better than the original paper claims!
 - ▶ Without introducing any overheads or additional reads
 - ▶ Relatively insensitive to choice of parameters

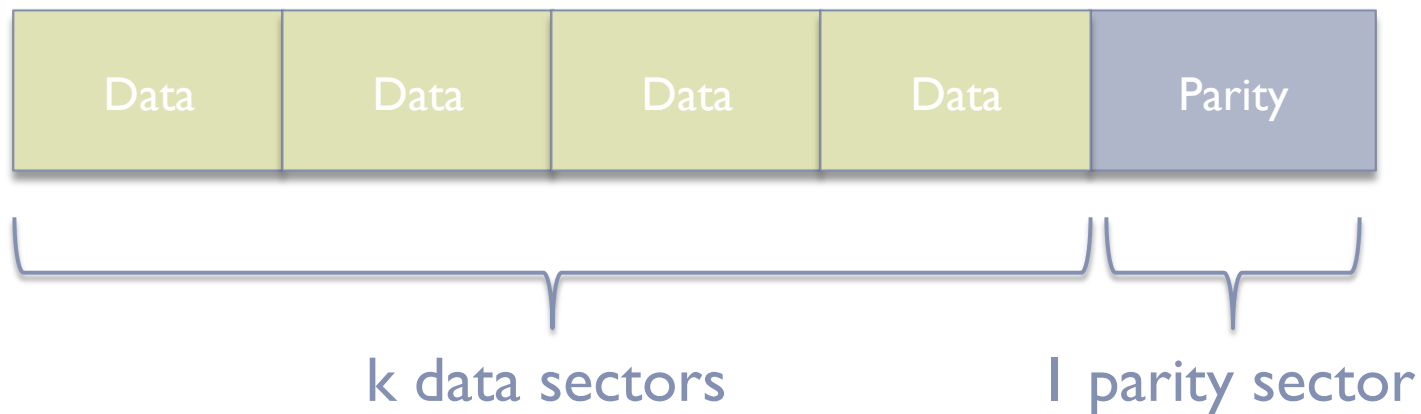
Intra-disk redundancy

- ▶ Why?
 - ▶ Recover LSEs in systems without redundancy
 - ▶ Recover LSEs during reconstruction in RAID-5

 - ▶ Goal:
 - ▶ Evaluate potential protection
 - ▶ What fraction of errors could be recovered
 - ▶ Qualitative discussion of overheads
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Intra-disk redundancy

- ▶ Simplest scheme: Single Parity Check (SPC)
- ▶ Can recover up to one LSE per parity group



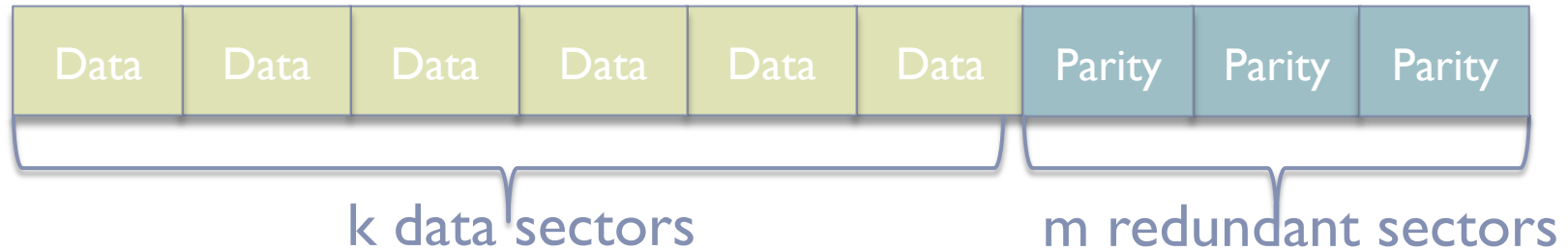
- ▶ Results from evaluation on Netapp data:
 - ▶ 25-50% of drives have errors that SPC cannot recover
- ▶ Consider stronger schemes?

Stronger schemes?

- ▶ Additional parity => additional overhead in updating parity
 - ▶ When would that be interesting?

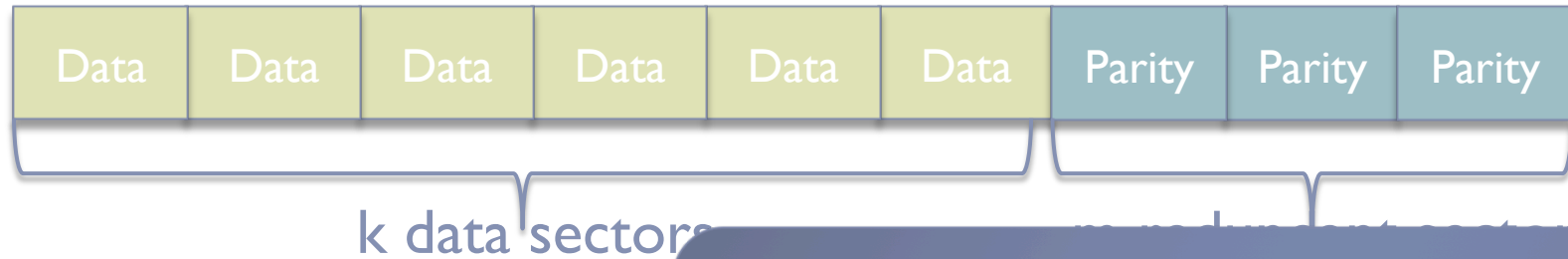
 - ▶ In environments
 - ▶ ... like archival systems, that don't have updates and don't like scrubs since they require powering up the system
 - ▶ ... with read-mostly workloads, i.e. parity updates are rare
 - ▶ ... for selected critical data on a drive, such as meta-data
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Inter-leaved Parity Check (IPC) [Dholakia08]



- ▶ Requires only 1 parity update per data update
 - ▶ Can tolerate up to m consecutive errors
-

Inter-leaved Parity Check (IPC) [Dholakia08]



- ▶ **Claim:** Achieves protection
- MDS=Maximum distance separable
- Expensive, but can tolerate

- Results differ from [Dholakia08]
- Importance of real-world data.
- Paper provides models & parameters

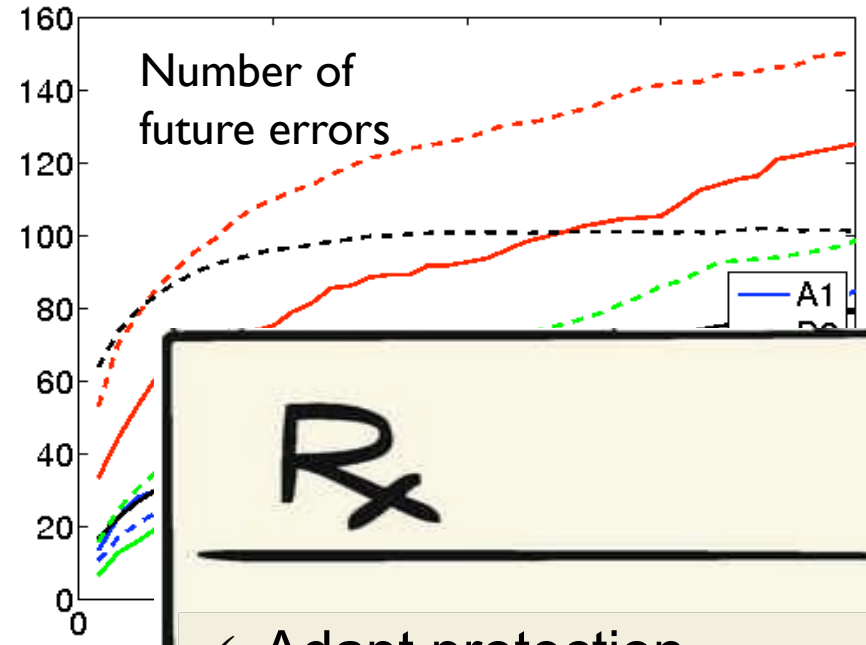
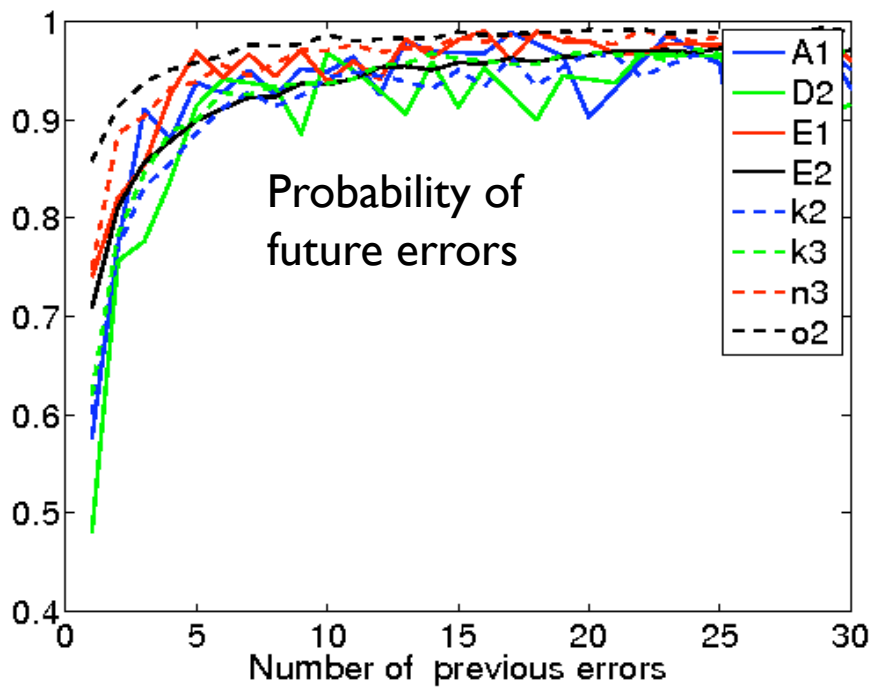
- ▶ **Results:** (from evaluation on NetApp data)
 - ▶ Far weaker than MDS!
 - ▶ Not significantly better than SPC

- ▶ **Implications**
 - ▶ Need different ideas for improving on SPC
 - ▶ Maybe reuse ideas from RAID-6? (see paper for details & results)

Questions unanswered ...

- ▶ **What level of protection to use when?**
 - ▶ E.g. what is the right scrub frequency?
 - ▶ Depends on error probability at a given time
-

Do previous errors predict future?

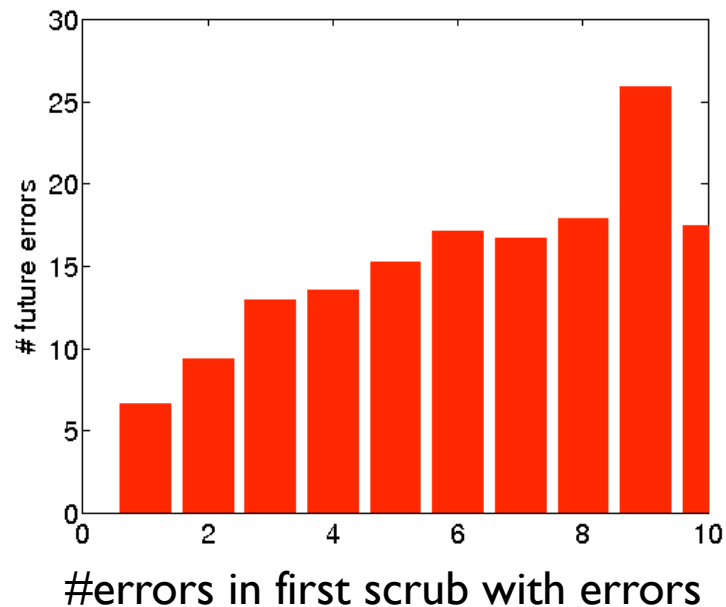


- ▶ Many previous errors
 - => higher chance of future errors
 - => higher number of future errors
- ▶ Big differences between models

R_x

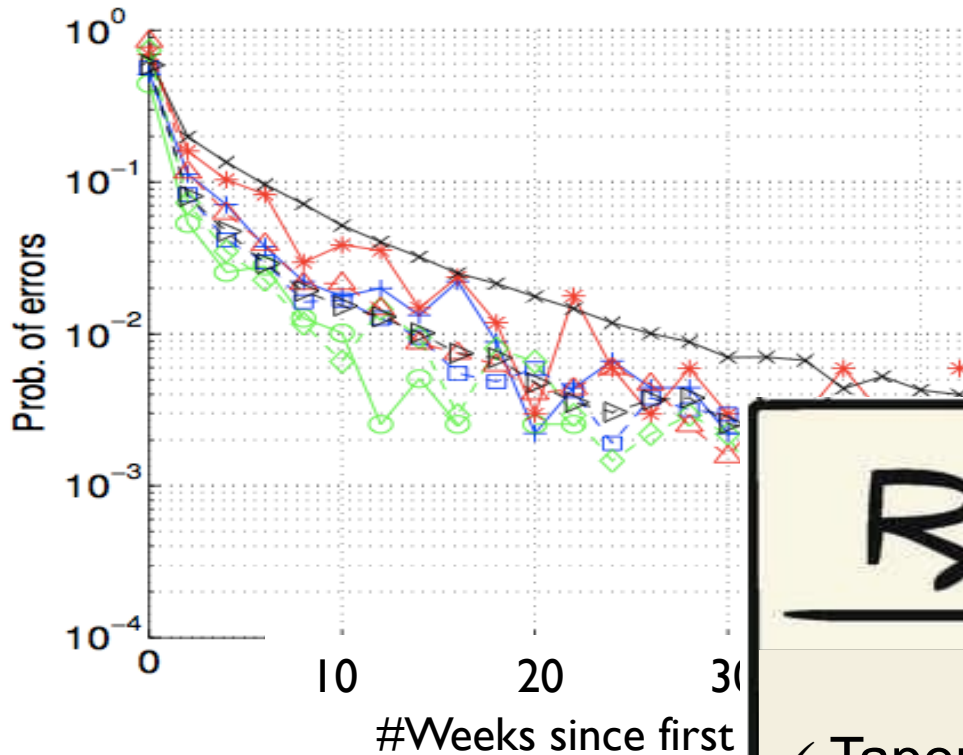
- ✓ Adapt protection based on previous errors
- ✓ Know your patient ..

Does *first* error interval predict future?



- ▶ Number of errors in first error interval:
 - Do increase expected number of future errors
 - Don't significantly increase probability of future occurrence
-

For how long are probabilities increased?




- ▶ Exponential drop-off, but still significant
- ▶ Independent of number of errors in

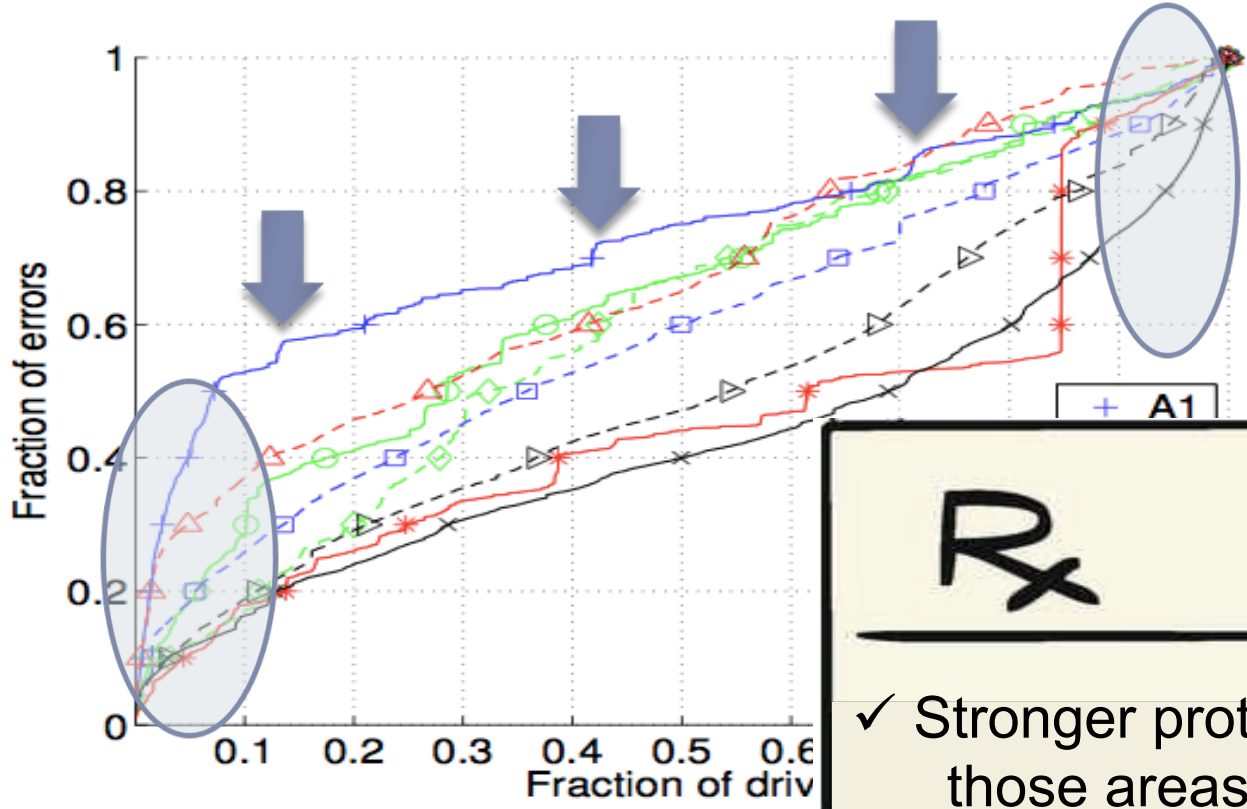
R_x

- ✓ Taper off added protection over time, e.g. reduce scrub rate

Questions unanswered ...

- ▶ **What level of protection to use when?**
 - ▶ What is the error probability at a given time?
 - ▶ **What level of protection to use where?**
 - ▶ Are all areas of the drive equally likely to develop errors?
-

Where on the drive are errors located?




- ▶ Up to 50% of errors concentrated in 1
- ▶ Also increased probability in some other

R_x

- ✓ Stronger protection for those areas
- ✓ Don't use for important data

Questions unanswered ...

- ▶ **What level of protection to use when?**
 - ▶ What is the error probability at a given time?
- ▶ **Same protection scheme across entire drive?**
 - ▶ Are all parts equally likely to develop errors?
- ▶ **Scrubbing potentially harmful?**
 - ▶ Do additional read operations increase error rate?

Does utilization affect LSEs?


- ▶ Collected data in Google data center (10,000 servers)
 - ▶ Number of LSEs
 - ▶ Number of reads & number of writes
- ▶ Results:
 - ▶ No correlation between #reads and LSEs
 - ▶ No correlation between #writes and LSEs

R_x

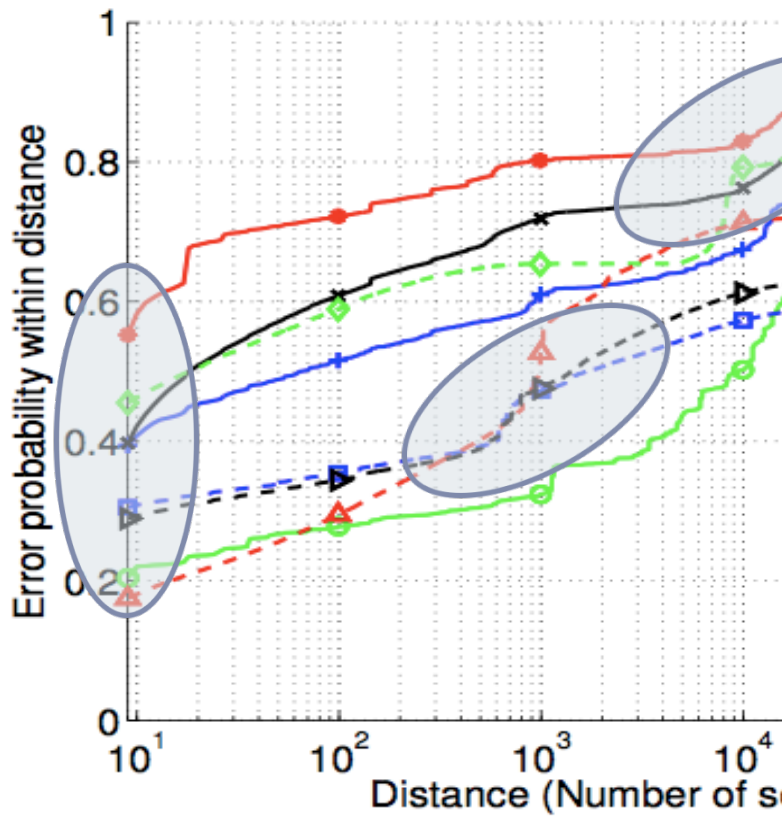
✓ Maybe need not worry about scrubs introducing new errors?

- ▶ Needs further investigation (future work).

Questions unanswered ...

- ▶ **What level of protection to use when?**
 - ▶ What is the error probability at a given time?
- ▶ **Same protection scheme across entire drive?**
 - ▶ Are all parts equally likely to develop errors?
- ▶ **Scrubbing potentially harmful?**
 - ▶ Do additional read operations increase error rate?
- ▶ **What is the common distance between errors ...**
 - ▶ Important for example for replica placement

How far are errors spaced apart?




- ▶ 20-60% of errors have a neighbor within distance
- ▶ Probability concentration (bumps)


R_x

- ✓ Avoid placing replicas at certain distances
- ✓ Explains why single parity scheme not always helpful

Questions unanswered ...

- ▶ What level of protection to use when?
 - ▶ What is the error probability at a given time?
- ▶ Different protection for different parts of the drive?
 - ▶ Are all parts equally likely to develop errors?
- ▶ Scrubbing potentially harmful?
 - ▶ Do additional read operations increase error rate?
- ▶ What is the common distance between errors ...
 - ▶ Important for replica placement
- ▶ Are errors that are close in space also close in time?
 - ▶ Yes!

Questions unanswered ...

- ▶ What level of protection to use when?
 - ▶ What is the error probability at a given time?
 - ▶ Different protection for different parts of the drive?
 - ▶ Are all parts equally likely to develop errors?
 - ▶ Scrubbing potentially harmful?
 - ▶ Do additional read operations increase error rate?
 - ▶ What is the common distance between errors ...
 - ▶ Important for replica placement
 - ▶ Are errors that are close in space also close in time?
 - ▶ Yes!
 - ▶ And many other questions – see paper!
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Conclusion

- ▶ **Evaluated potential of different protection schemes**
 - ▶ **Scrubbing**
 - ▶ Simple new scheme (staggered scrubbing) performs very well!
 - ▶ **Intra-disk redundancy**
 - ▶ Single parity can recover LSEs in 50-75% of the drives
 - ▶ Need to look at more complex schemes for coverage beyond that
 - ▶ **Looked at characteristics of LSEs**
 - ▶ And how to exploit them for reliability
 - ▶ **Many characteristics not captured well by simple models**
 - ▶ Provided parameters for models
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▶ **Thanks!**

- ▶ To NetApp for sharing the data
- ▶ To you for listening

▶ **Questions?**
