A geometric model for hadleystem on-line social networks and of the second

Anthony Bonato

Ryerson University

dahowlet

chincheffe

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WOSN'10

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Geometric model for OSNs



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chrisse

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Complex Networks

 web graph, social networks, biological networks, internet networks, ...



On-line Social Networks (OSNs) Facebook, Twitter, LinkedIn, MySpace...



Properties of OSNs

• observed properties:

(Kumar et al,06):

- power law degree distribution, small world
- community structure
- densification power law and shrinking distances



Figure 6: Average and effective diameter of the giant component of Flickr and Yahoo! 360 timegraphs, by week.

Why model complex networks?

- uncover and explain the generative mechanisms underlying complex networks
- predict the future
- nice mathematical challenges
- models can uncover the hidden reality of networks

Many different models



Models of OSNs

- relatively few models for on-line social networks
- goal: find a model which simulates many of the observed properties of OSNs
 must evolve in a natural way...

"All models are wrong, but some are more useful." – G.P.E. Box

Transitivity



Iterated Local Transitivity (ILT) model (Bonato, Hadi, Horn, Prałat, Wang, 08)

- key paradigm is transitivity: friends of friends are more likely friends
- start with a graph of order n
- to form the graph G_{t+1} for each node x from time t, add a node x', the clone of x, so that xx' is an edge, and x' is joined to each node joined to x

 $G_0 = C_4$



Properties of ILT model

- densification power law
- distances decrease over time
- community structure: bad spectral expansion (Estrada, 06)

... Degree distribution



Geometry of OSNs?

- OSNs live in social space: proximity of nodes depends on common attributes (such as geography, gender, age, etc.)
- IDEA: embed OSN in 2-, 3or higher dimensional space



Dimension of an OSN

- dimension of OSN: minimum number of attributes needed to classify or group users
- like game of "20 Questions": each question narrows range of possibilities
- what is a credible mathematical formula for the dimension of an OSN?

Random geometric graphs



- nodes are randomly placed in space
- each node has a constant sphere of influence
- nodes are joined if their sphere of influence overlap

Simulation with 5000 nodes



Spatially Preferred Attachment (SPA) model (Aiello, Bonato, Cooper, Janssen, Prałat, 08)

- volume of sphere of influence proportional to in-degree
- nodes are added and spheres of influence shrink over time
- asymptotically almost surely (a.a.s.) leads to power laws graphs



Protean graphs

(Fortunato, Flammini, Menczer,06), (Łuczak, Prałat,06), (Janssen, Prałat,09)

- parameter: α in (0,1)
- each node is ranked 1,2, ..., n by some function r
 - 1 is best, n is worst
- at each time-step, one new node v is born, one randomly node chosen dies (and ranking is updated)
- link probability r^{-α}
- many ranking schemes a.a.s. lead to power law graphs: random initial ranking, degree, age, etc.

- we consider a geometric model of OSNs, where
 - nodes are in mdimensional hypercube in Euclidean space
 - volume of sphere of influence variable: a function of ranking of nodes



Geometric Protean (GEO-P) Model (Bonato, Janssen, Prałat, 10)

- parameters: α , β in (0,1), $\alpha+\beta < 1$; positive integer m
- nodes live in m-dimensional hypercube
- each node is ranked 1,2, ..., n by some function r
 we use random initial ranking
- at each time-step, one new node v is born, one randomly node chosen dies (and ranking is updated)
- each existing node u has a sphere of influence with volume $r^{-\alpha}n^{-\beta}$
- add edge uv if v is in the region of influence of u

Notes on GEO-P model

- models uses both geometry and ranking
- number of nodes is static: fixed at n
 - order of OSNs at most number of people (roughly...)
- top ranked nodes have larger regions of influence

Simulation with 5000 nodes



Simulation with 5000 nodes



random geometric

GEO-P

Properties of the GEO-P model (Bonato, Janssen, Prałat, 2010)

- a.a.s. the GEO-P model generates graphs with the following properties:
 - power law degree distribution with exponent

 $b = 1 + 1/\alpha$

- average degree d = $(1+o(1))n^{(1-\alpha-\beta)}/2^{1-\alpha}$
 - densification
- diameter $D = O(n^{\beta/(1-\alpha)m} \log^{2\alpha/(1-\alpha)m} n)$
 - small world: constant order if m = Clog n

Degree Distribution

 for m < k < M, a.a.s. the number of nodes of degree at least k equals

$$(1+O(\log^{-1/3} n))\left(\frac{\alpha}{\alpha+1}\right)n^{(1-\beta)/\alpha}k^{-1/\alpha}$$

- $m = n^{1 \alpha \beta} \log^{1/2} n$
 - m should be much larger than the minimum degree
- $M = n^{1 \alpha/2 \beta} \log^{-2\alpha 1} n$
 - for k > M, the expected number of nodes of degree k is too small to guarantee concentration

Density

- $i^{-\alpha}n^{-\beta}$ = probability that new node links to node of rank i
- average number of edges added at each time-step

$$\sum_{i=1}^{n} i^{-\alpha} n^{-\beta} \approx \frac{1}{1-\alpha} n^{1-\alpha-\beta}$$

- parameter β controls density
- if β < 1 α, then density grows with n (as in real OSNs)

Diameter

- eminent node:
 - old: at least n/2 nodes are younger
 - highly ranked: initial ranking greater than some fixed R
- partition hypercube into small hypercubes
- choose size of hypercubes and R so that
 - a.a.s. each hypercube contains at least log²n eminent nodes
 - sphere of influence of each eminent node covers each hypercube and all neighbouring hypercubes
- choose eminent node in each hypercube: backbone
- show a.a.s. all nodes in hypercube distance at most 2 from backbone



Spectral properties

- the spectral gap λ of G is defined by the difference between the two largest eigenvalues of the adjacency matrix of G
- for G(n,p) random graphs, λ tends to 0 as order grows
- in the GEO-P model, λ is close to 1
- bad expansion/big spectral gaps in the GEO-P model found in social networks but not in the web graph (Estrada, 06)
 - in social networks, there are a higher number of intrarather than inter-community links

Dimension of OSNs

- given the order of the network n, power law exponent b, average degree d, and diameter D, we can calculate m
- gives formula for dimension of OSN:

$$m = \frac{\log\left(\frac{n}{\frac{b-1}{2d^{\frac{b-1}{b-2}}}}\right)}{\log D}$$

Uncovering the hidden reality

- reverse engineering approach
 - given network data (n, b, d, D), dimension of an OSN gives smallest number of attributes needed to identify users
- that is, given the graph structure, we can (theoretically) recover the social space



6 Dimensions of Separation

OSN	Dimension
YouTube	6
Twitter	4
Flickr	4
Cyworld	7

Research directions

- fitting GEO-P model to data
 - is theoretical estimate of log n dimension accurate?
 - find similarity measures (see PPI literature)
- community detection
 - first map network in social space?
- spread of influence
 - SIS, SIR models
 - Graph theory: firefighting, Cops and Robbers

preprints, reprints, contact: search: "Anthony Bonato"





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