

Capo: Recapitulating Storage for Virtual Desktops

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The World According to Gartner

- 60% of all enterprises will deploy network computers by 2001
- 5 – 30 million Windows terminals sold per year by 2005
- 40% of desktops – 49 million – will be virtualized by 2013



Sooner or later they're bound to be right

Virtual Desktop Infrastructure (VDI)



- IT admins love it
 - Centrally administered
 - Reduced hardware and maintenance costs
- Users will embrace it (hopefully)
 - Familiar personal computing environment
 - *Performance (latency) is critical*

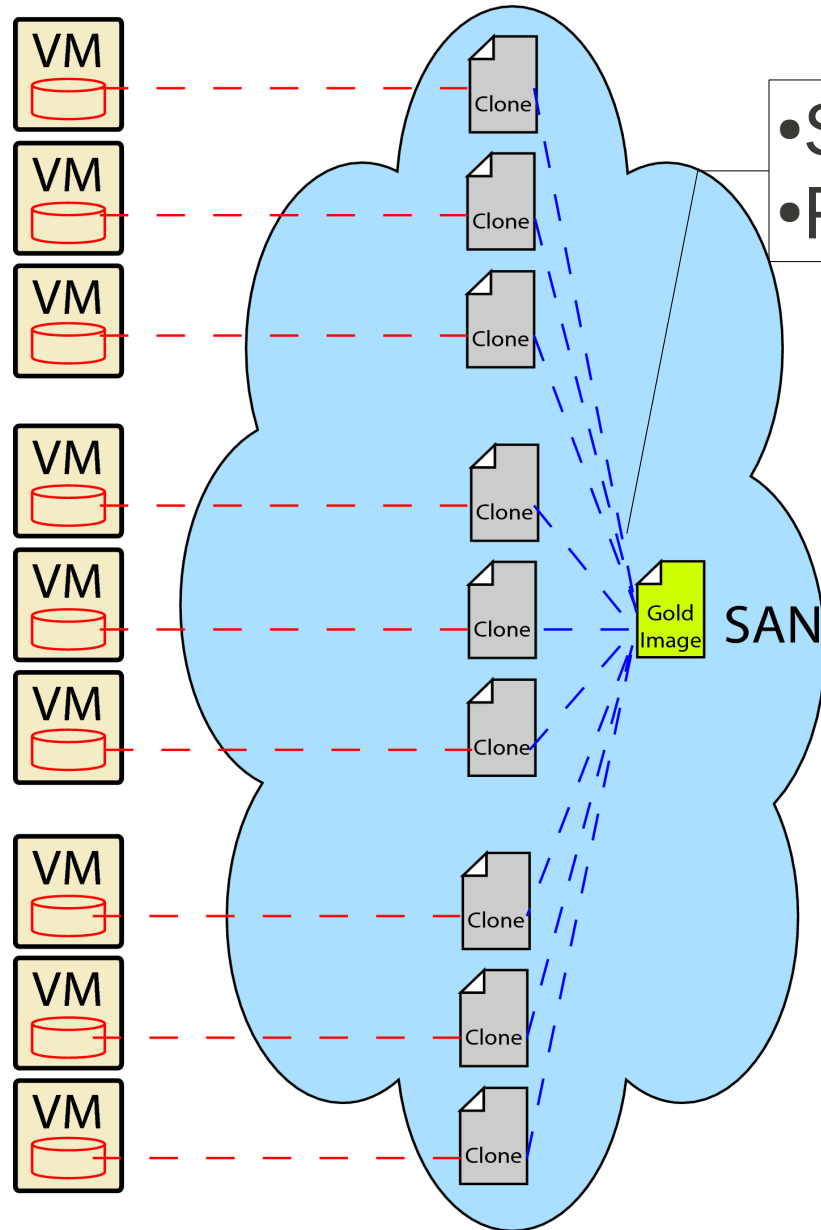
Improving VDI

- How can we:
 - Reduce the cost of VDI deployments?
 - Improve the user experience?

Outline

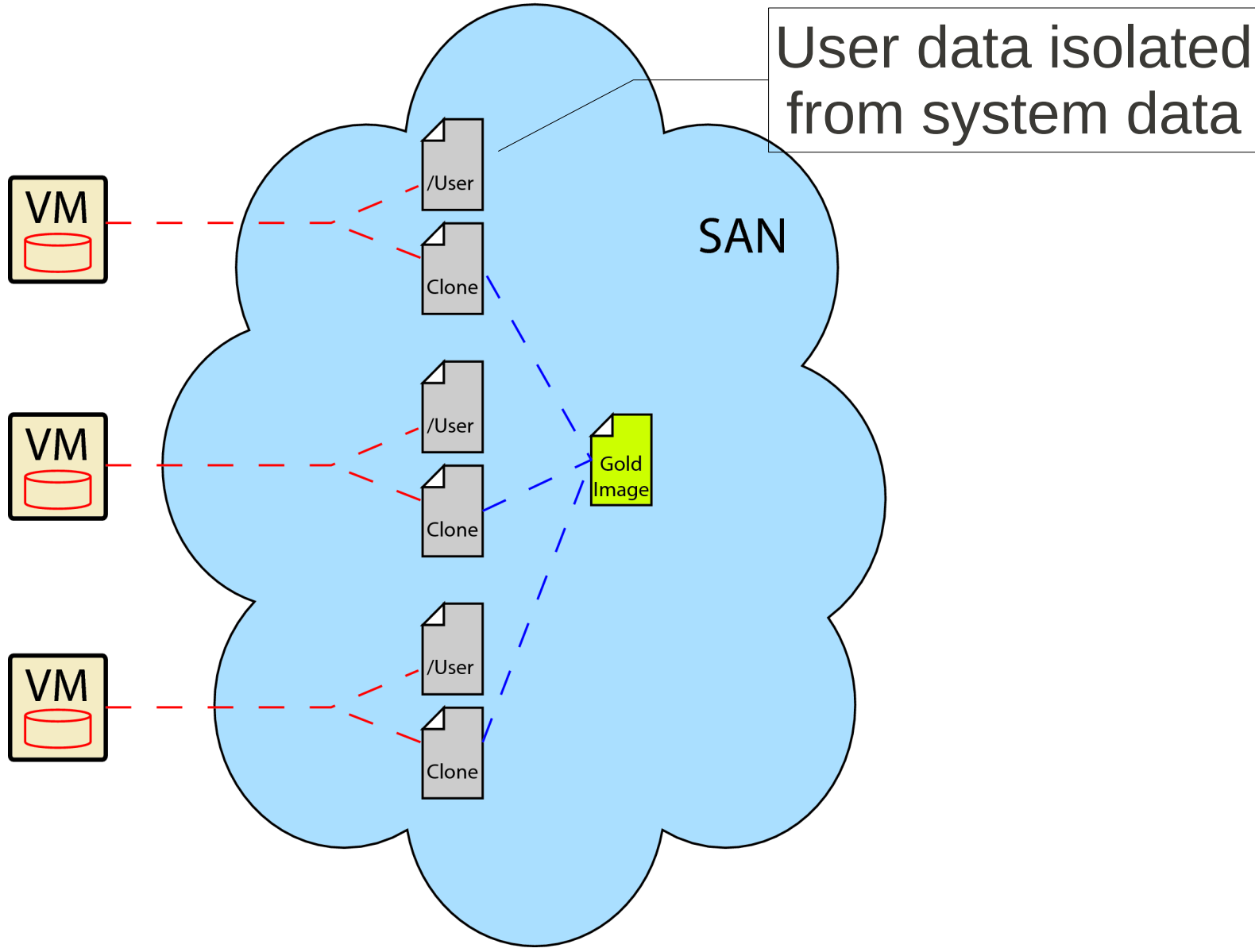
- Background
- How VDI Works
- UBC Workload Analysis
- Design and Implementation
- Evaluation
- Conclusion

How VDI Works

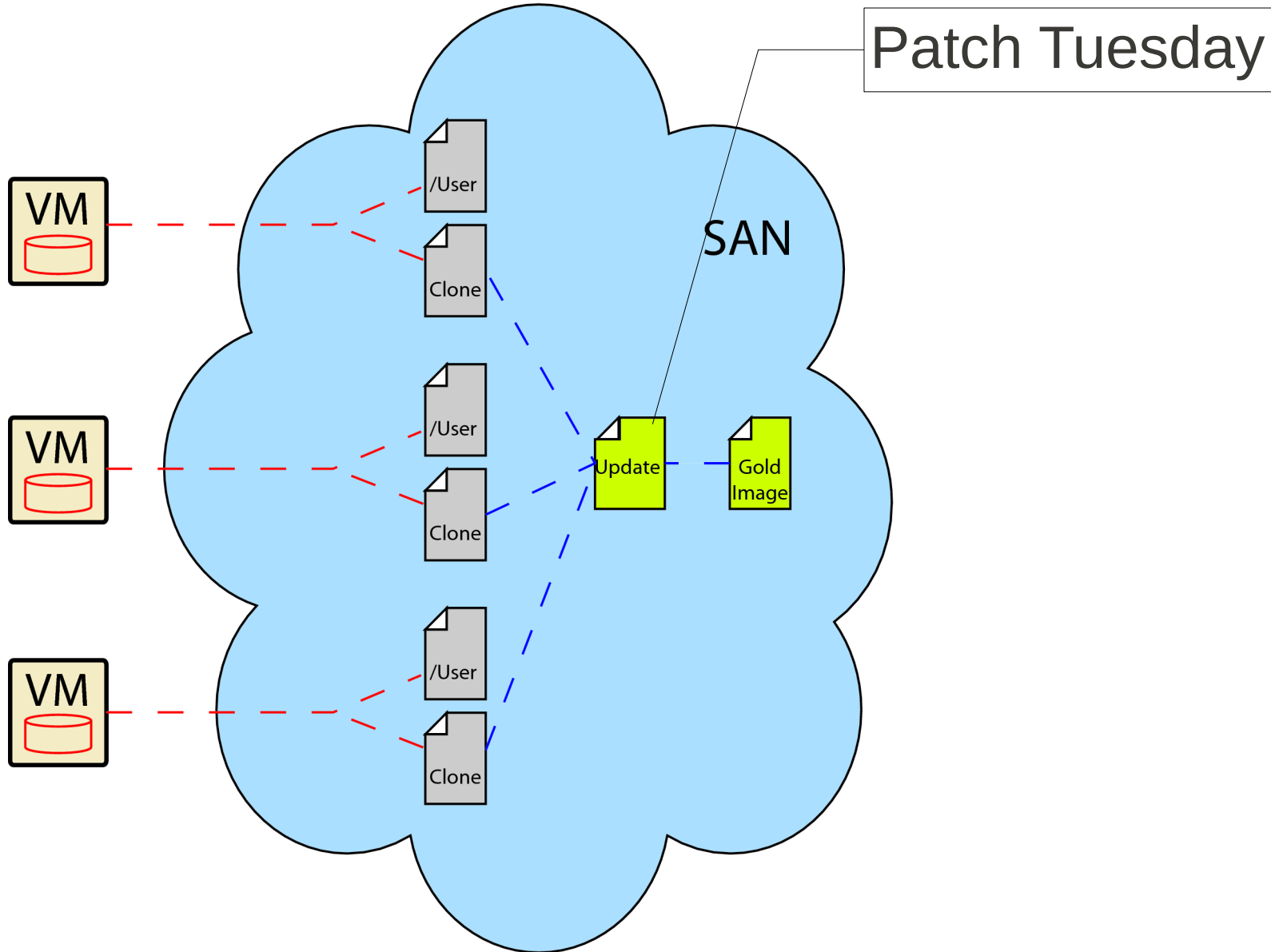


- Sparse allocation
- Fast clone

How VDI Works



How VDI Works



UBC Workload Analysis



15506-41DG
'Office: 9am' Disc
© Jupiterimages

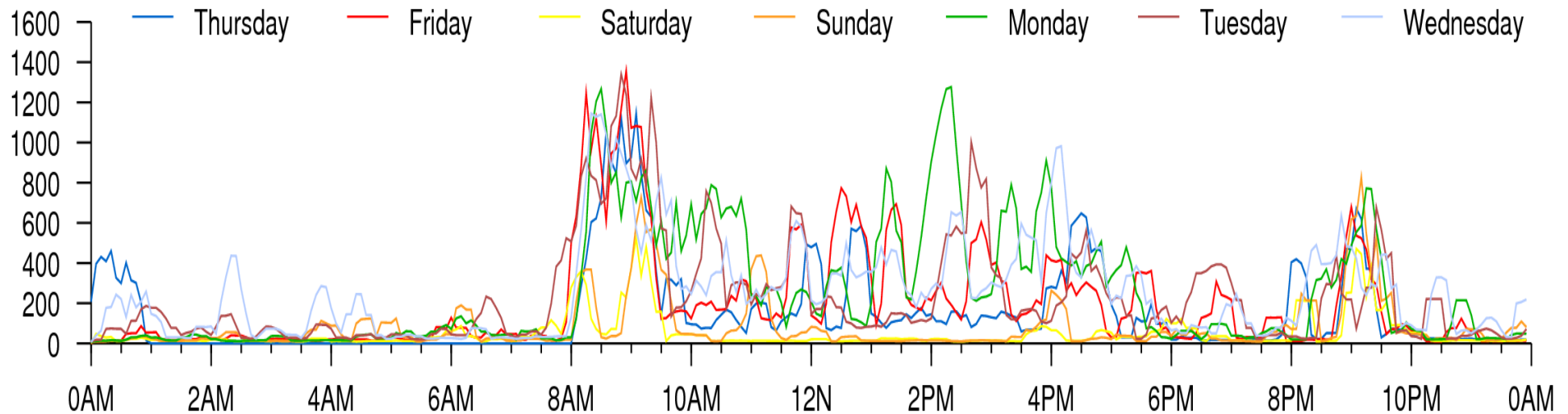
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UBC Workload Analysis

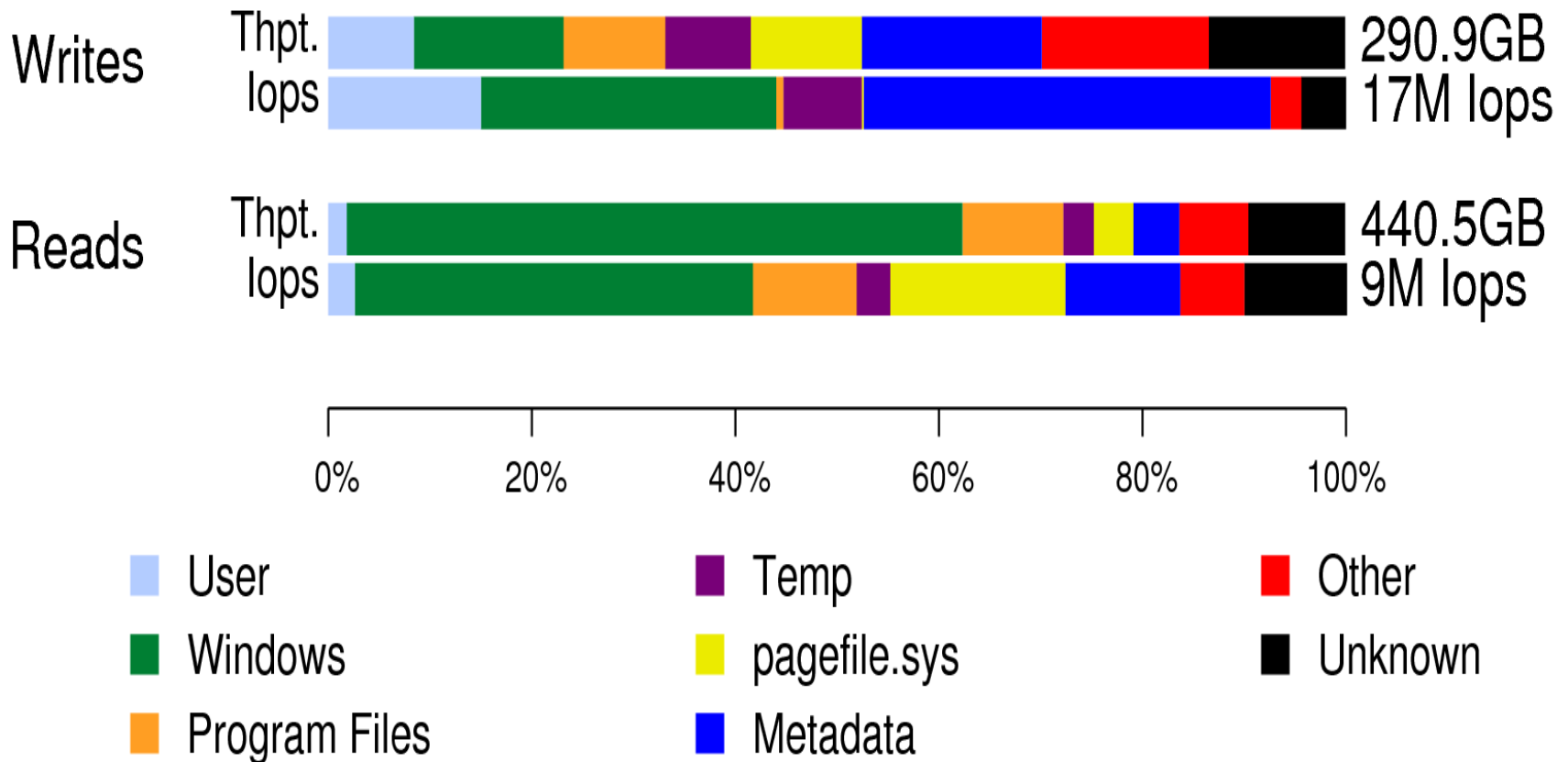
- We profiled 55 Windows Vista desktops
 - Administrative offices at UBC
 - Installed profiler during regularly-scheduled weekly update
 - Captured file- and block-level accesses
 - Collected 75 GB of compressed, binary logs

Workload: Day-to-Day Activity

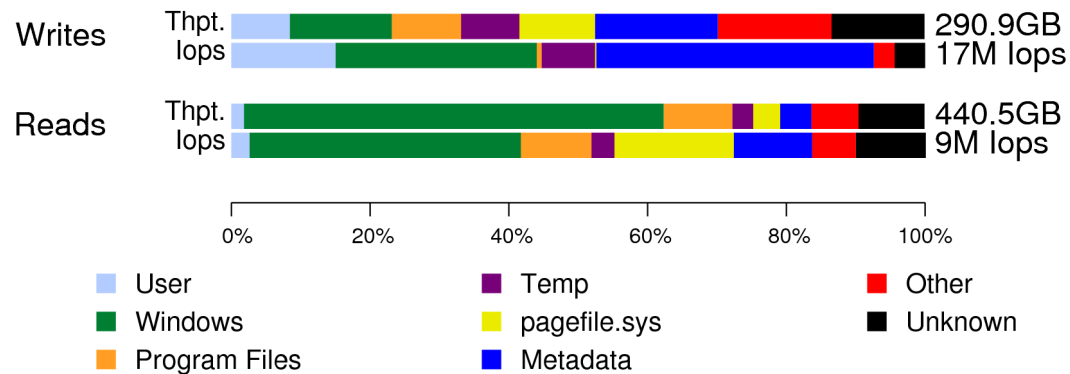


- Workload is quite bursty
- What do the requests look like?

Workload: IO Requests



Workload: IO Requests



- Most accesses to system-controlled objects (\Windows, \Program Files)
- Metadata-heavy workload
- IOps: 65% writes; throughput: 65% reads
- What do these writes look like?

Workload: Write Requests

- Fairly high churn rate
 - 8% of bytes re-written in 10 seconds
 - 50% of bytes re-written in 24 hours
- Average divergence of 1GB after about an hour
 - A large portion of this is from pagefile.sys and other Windows files

Workload Summary

- VDI workloads are bursty
- Significant sharing among VMs
- High churn rate for hot data
- Namespace not accessed uniformly

VDI Storage Scalability



How can we improve VDI storage scalability?



Local Persistent Cache

- Goal: *offload IOps to local disks*
- Library in dom0 interposes on access to network files
- Cached files stored on local file system
 - Bitmaps track sparse files
- Supports write-through and write-back with adjustable window

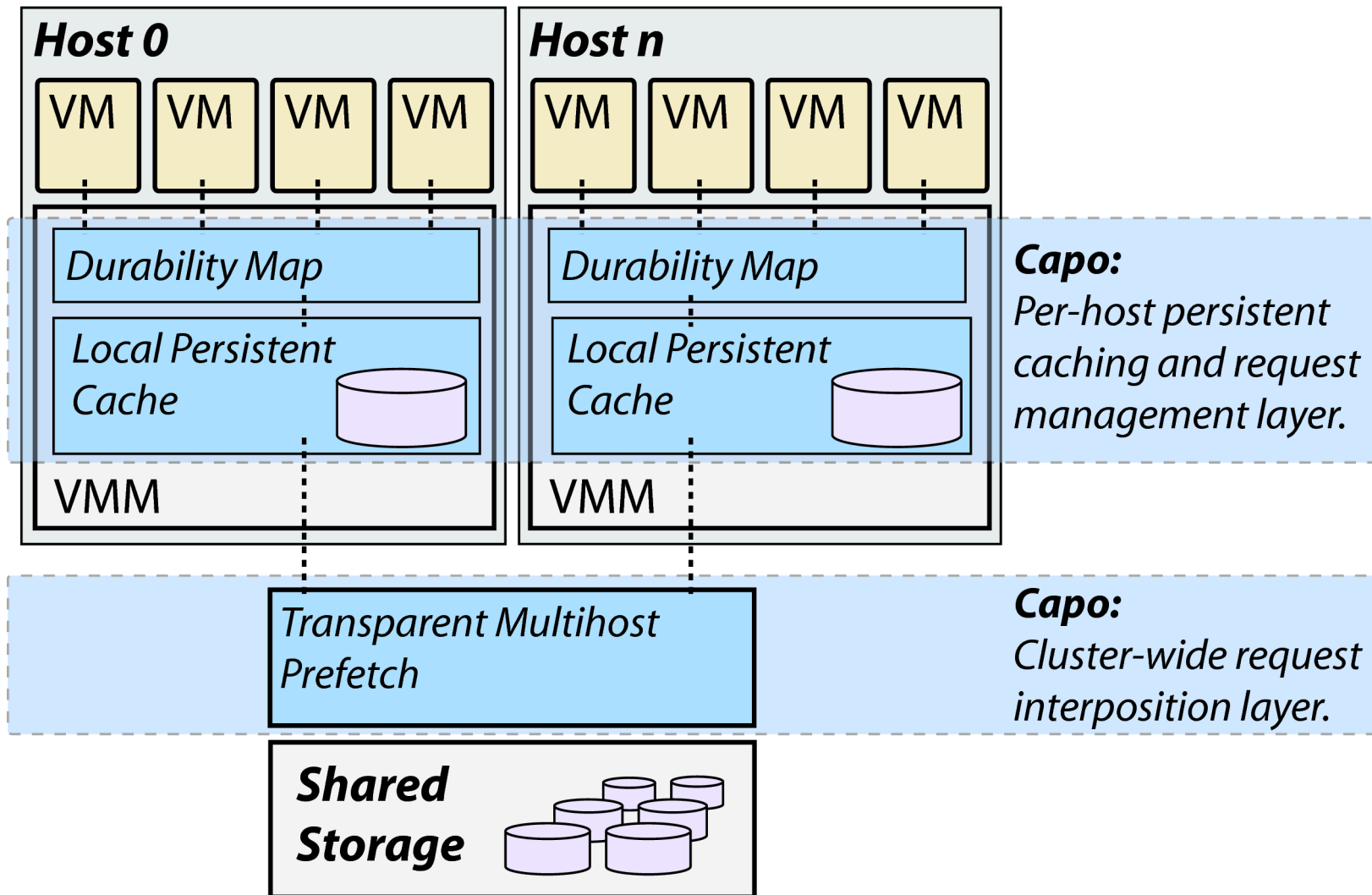
Multihost Preloader

- Goal: *share local caches among all hosts*
- NFS proxy snoops requests to cached files
- Shared data is multicasted to all subscribed hosts
- Basic congestion control and/or isolated network required

Differential Durability

- Goal: *reduce writeback burden*
- Data categorized according to value
 - User-created data, installer data, temporary files, pagefile.sys
- Gold Image disk partitioned
 - Valuable data placed on a disk with an aggressive write-through policy
 - Expendable data stored on cheap local disks

Capo Architecture Diagram



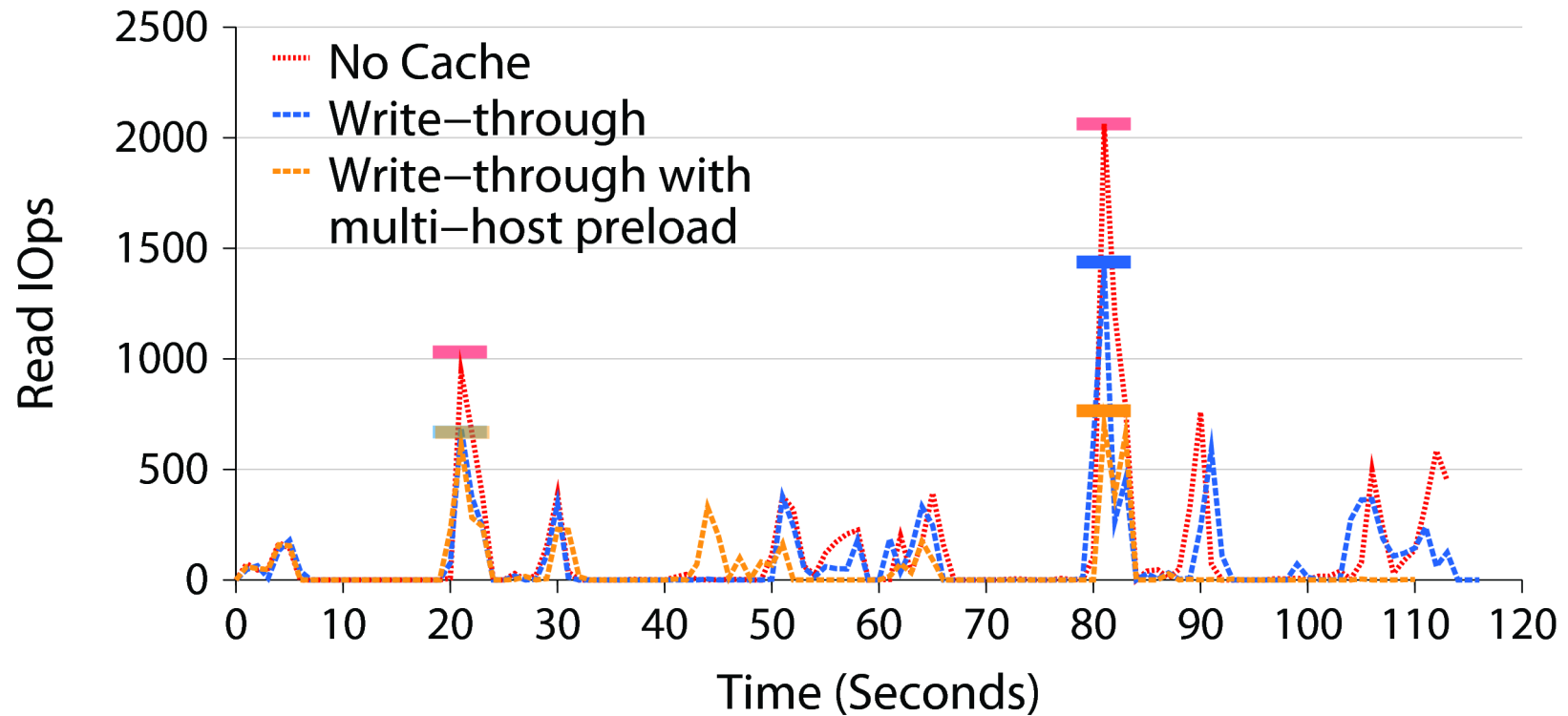
Implementation

- Local cache prototyped during summer internship
 - 7,000 lines of C code
- Prefetcher implemented three (3) times
 - Packet capture, unfsd, RPC proxy
- IntelliCache™ now available in latest XenServer releases

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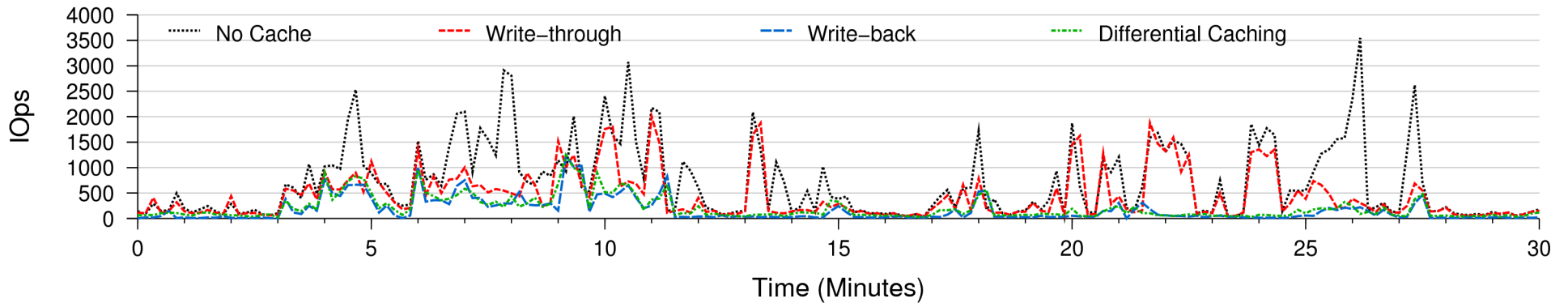
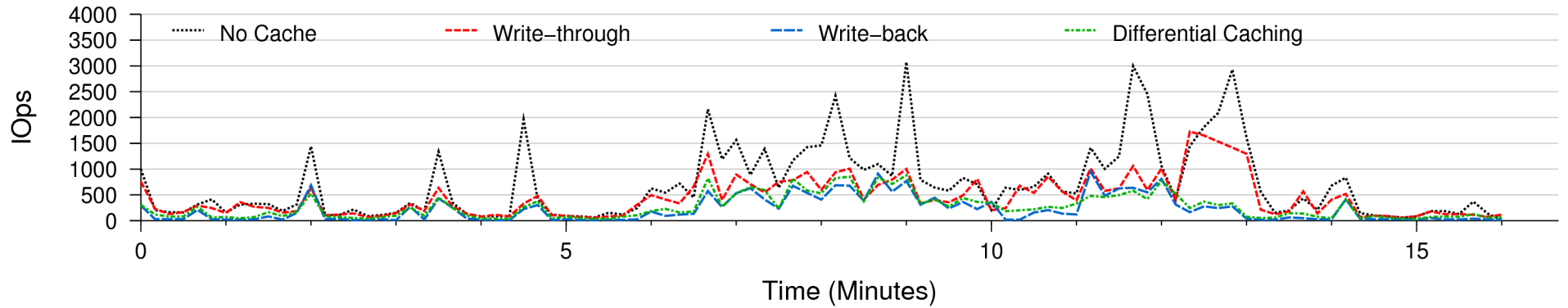
Microbenchmarks: Preloading



Trace Replay: Methodology

- Original environment:
 - VMware
 - SAN
- Our lab:
 - XenServer
 - Linux NFS filer w/ 6-disk RAID 0 volume
- Replayer:
 - Simple perl application
 - Tries to match original trace request pattern

Trace Replay: Selected Peaks



Conclusion

- VDI presents new challenges for storage systems
- Central storage is a reasonable solution...
- But local caching and differential durability can help reduce costs and improve performance

Questions?