

Zurich Research Laboratory

# Improving Efficiency and Enhancing Concurrency of Untrusted Storage

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### Where is my data?













2008

# 1980



## **Untrusted Storage Service**



- Many independent clients
  Correct
  Store data on server
  Communicate only with server
  Small trusted memory
- Storage server
  - Untrusted
  - Potentially corrupted
- Clients read and write concurrently

#### How to ensure consistent view of data to all clients?



### **Consistent Access to Untrusted Storage**

- Loose synchronization and concurrency pose a new problem
- Suppose clients sign data with digital signatures: Server cannot forge any values ...
  - $\rightarrow$  But answer with outdated value ("replay attack")
  - $\rightarrow\,$  Or send different values to different clients
- Server may present different views to clients
  - $\rightarrow\,$  "Fork" their views of history
  - $\rightarrow$  Clients cannot prevent this
- Fork linearizability [MS02], provided by SUNDR [LKMS04]
  - If server forks the views of two clients once, then
  - → their views are forked ever after
  - → they never again see any updates of each other
- Forks are easier to detect than subtle data modifications
  - $\rightarrow$  Using a separate channel for detection



## **Fork-linearizability**





After  $C_1$  writes u,

C<sub>2</sub> reads x:

 $\rightarrow$  C<sub>2</sub> forked from C<sub>1</sub>C<sub>3</sub>

After C<sub>1</sub> reads y:

 $\rightarrow$  C<sub>1</sub> forked from C<sub>3</sub>



## **New Results**

- More efficient fork-linearizable communication protocol [CSS07]
  - $\rightarrow$  Messages of size O(n) instead of O(n<sup>2</sup>) with n clients
- Fork-linearizable protocols are not wait-free [CSS07]
  - $\rightarrow\,$  Reader must wait for writer even if server correct
- New notion: weak fork-linearizability [CKS08]
  - $\rightarrow$  New wait-free protocol, where clients need not wait for each other and messages of size O(n) only
- More impossibility results [CKS08]
  - $\rightarrow$  Fork-sequential consistency does not enable wait-free protocols
  - $\rightarrow\,$  Fork-\* consistency does not enable wait-free protocols



### References

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- [LKMS04] J. Li, M. Krohn, D. Mazières, and D. Shasha. Secure untrusted data repository (SUNDR). In Proc. Symp. Operating Systems Design and Implementation (OSDI), 2004.
- [MS02] D. Mazières and D. Shasha. Building secure file systems out of Byzantine storage. In Proc. 21st ACM Symp. Principles of Distributed Computing (PODC), 2002.