

Efficient Query Computing for Uncertain Possibilistic Databases with Provenance

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June 20, 2011

The problem we investigate:

How to compute answers to queries for uncertain data with attached “confidence values”?

Problem 1

Existing models for **uncertain** data (e.g., possibility theory) are **not closed** for SPJ queries.

Problem 2

Computing **probabilities** of SPJ answers in models that combine **uncertainty** and **provenance** is a problem with **#P complexity**.

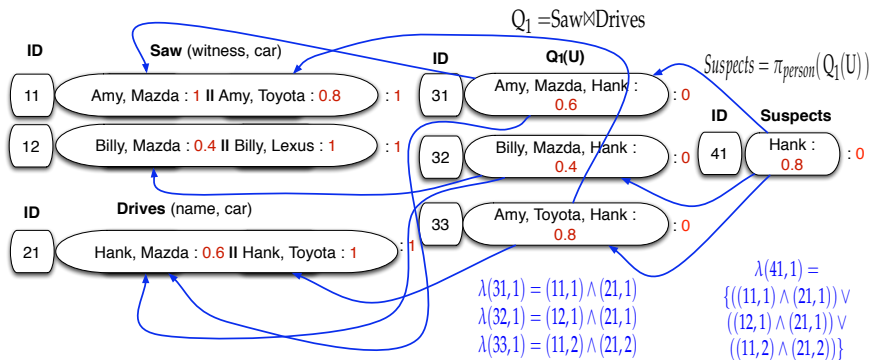
We propose:

A data model that combines **uncertainty**, **provenance** and **possibilities**.

Benefits of the proposed model:

- **Closed** for SPJ queries
- Computing **possibilities** of SPJ answers has **PTIME complexity**.

Motivating Example



Possible Worlds:



Possibilistic Databases are not closed

ID	$Q_1 = \text{Saw} \bowtie \text{Drives}$	
31	Amy, Mazda, Hank : 0.6	: 0
32	Billy, Mazda, Hank : 0.4	: 0
33	Amy, Toyota, Hank : 0.8	: 0

- Tuples 31 and 33 **cannot coexist** in any possible world.
- **Adding provenance** (lineage) makes possibilistic uncertain data model **closed** for SPJ queries.

A. Das Sarma et al. "Representing Uncertain Data: Models, Properties, and Algorithms". VLDB Journal, October, 2009.
P. Bosc and O. Pivert. "About PSJ queries addressed to possibilistic relational databases". IEEE T. Fuzzy Systems, 2005.

Computing probabilities is intractable

$$\text{Suspects} = \pi_{\text{person}}(Q_1(U))$$

ID	Suspects
41	Hank

$$\lambda(41, 1) = \{((11, 1) \wedge (21, 1)) \vee ((12, 1) \wedge (21, 1)) \vee ((11, 2) \wedge (21, 2))\}$$

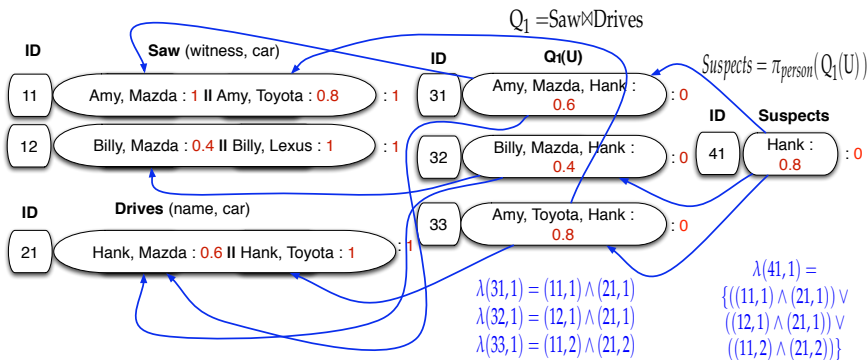
: Computing probabilities has #P complexity

N. Dalvi and D. Suciu. "Efficient query evaluation on probabilistic databases." In VLDB 2004.

Axioms of Possibilistic theory:

- $\Pi(X) = 1$
- $\Pi(\emptyset) = 0$
- $\Pi(E_1 \cup E_2) = \max(\Pi(E_1), \Pi(E_2))$
- $\Pi(E_1 \cap E_2) \leq \min(\Pi(E_1), \Pi(E_2))$
- $\Pi(E_1 \cap E_2) = \min(\Pi(E_1), \Pi(E_2))$
(for not-interactive events)
- $\max\{\Pi(E), \Pi(\bar{E})\} = 1$
- $N(E) = 1 - \Pi(\bar{E})$

Running Example



Possible Worlds:



Ongoing Work:

- Extend the query language to extensions of conjunctive queries.
- Find for which class of query languages the problem remains in polynomial time.
- Find for which class of query languages the problem becomes intractable.
- Study complexity of new query languages that can query over uncertainty and provenance.

Thank you