

Middleware for Gossip Protocols

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Motivation

- Gossip protocols are highly robust
- Problematic when an error does occur
 - E.g. Amazon S3 – 6 hours to fix an otherwise simple problem
 - Want to offer a way to fix such problems without having to take down the entire system

Contributions

Design, implementation, and analysis of gossip middleware that supports rapid code updating

Talk Outline

- Versions and Deployments
- Architecture
- Evaluation
- Conclusion and Future Work

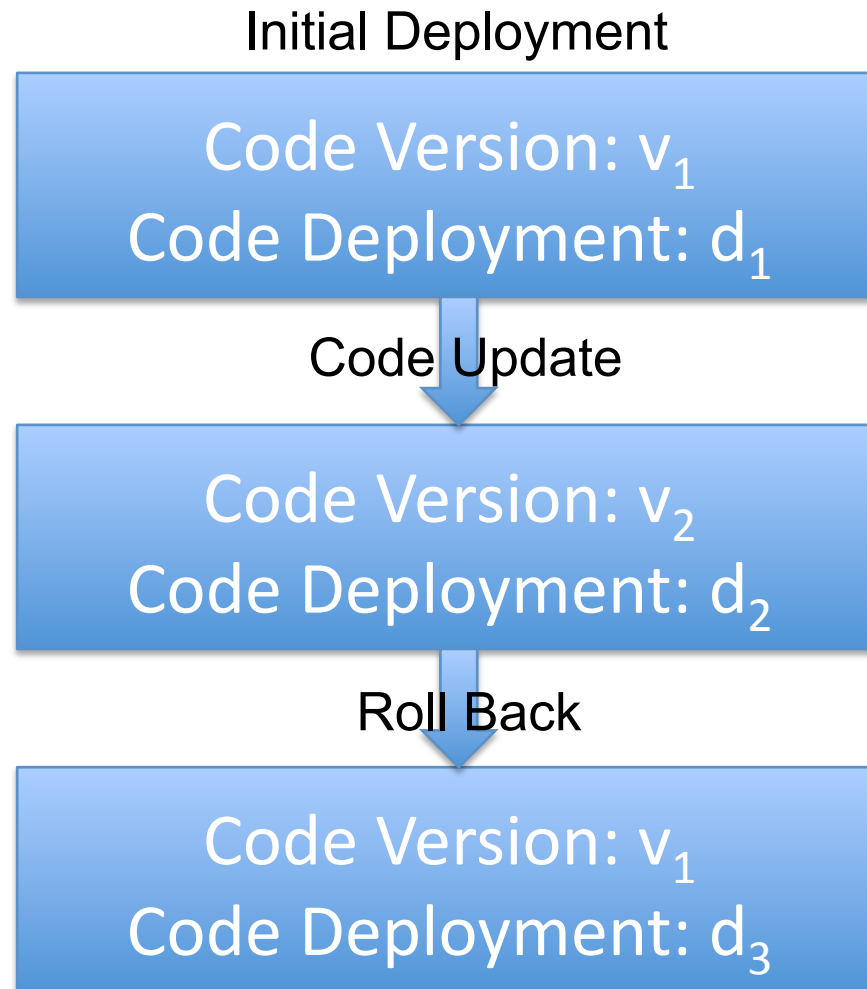
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Versions and Deployments

- *Modules*: Gossip application instances
- Each module assigned a *Deployment Number*
 - Identifies originating node and time of deployment
 - Used to determine whether or not nodes are running the correct version of the application
 - Does not correspond with code version

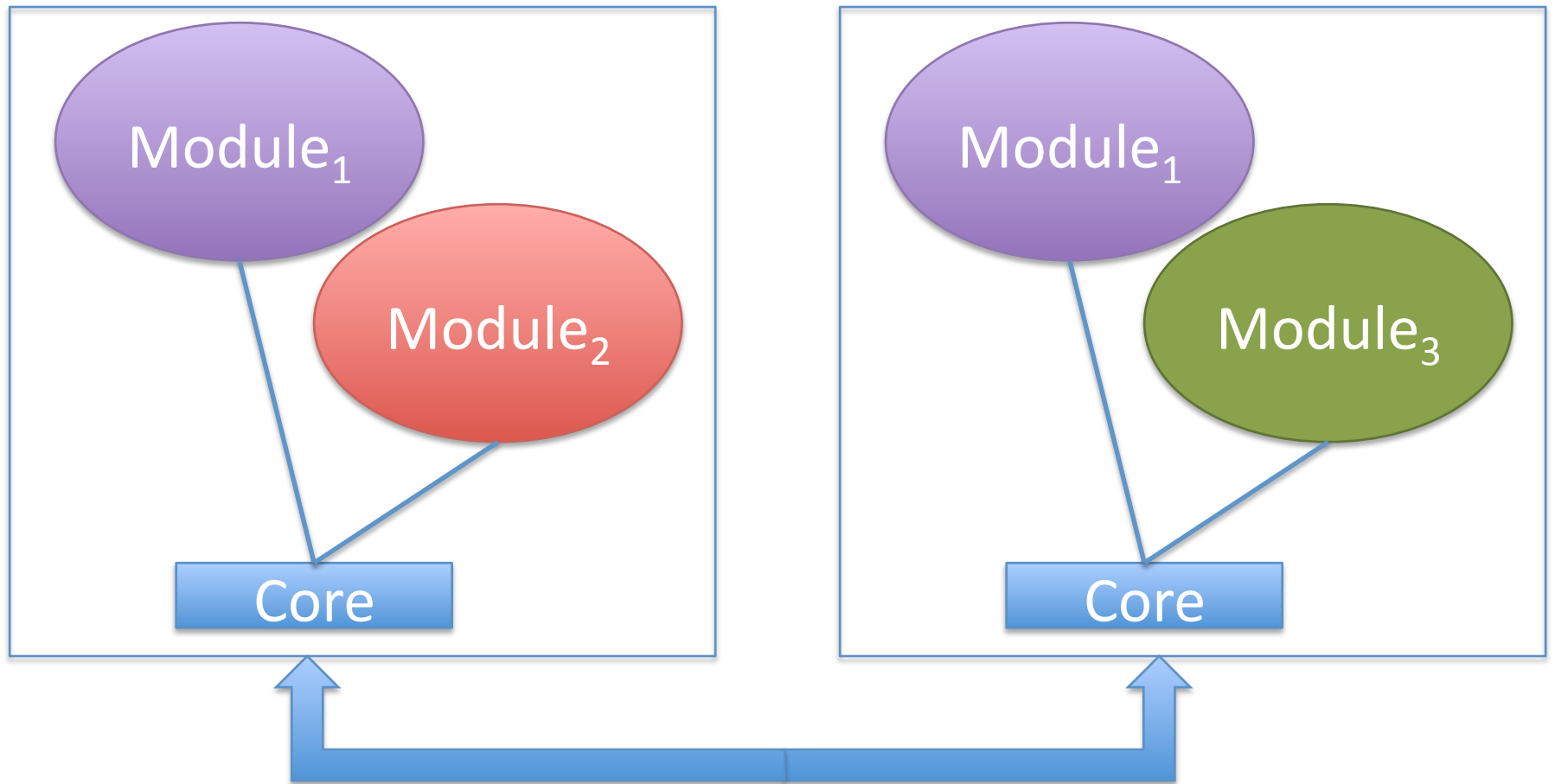
Versions and Deployments



Talk Outline

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Architecture



Core

- Provides Module Management and Updating
- Core gossips deployment numbers and corresponding code versions
- Core itself cannot be updated this way
- *Challenge: keep core small*
- Approach: core leverages ongoing gossip between modules

Module Management

- Core maintains a configuration file that contains:
 - List of Modules and current versions (identified by hash codes of the class files)
 - Deployment Number
- Keeps track of which modules and corresponding versions are currently running
- Cores gossip Configuration files

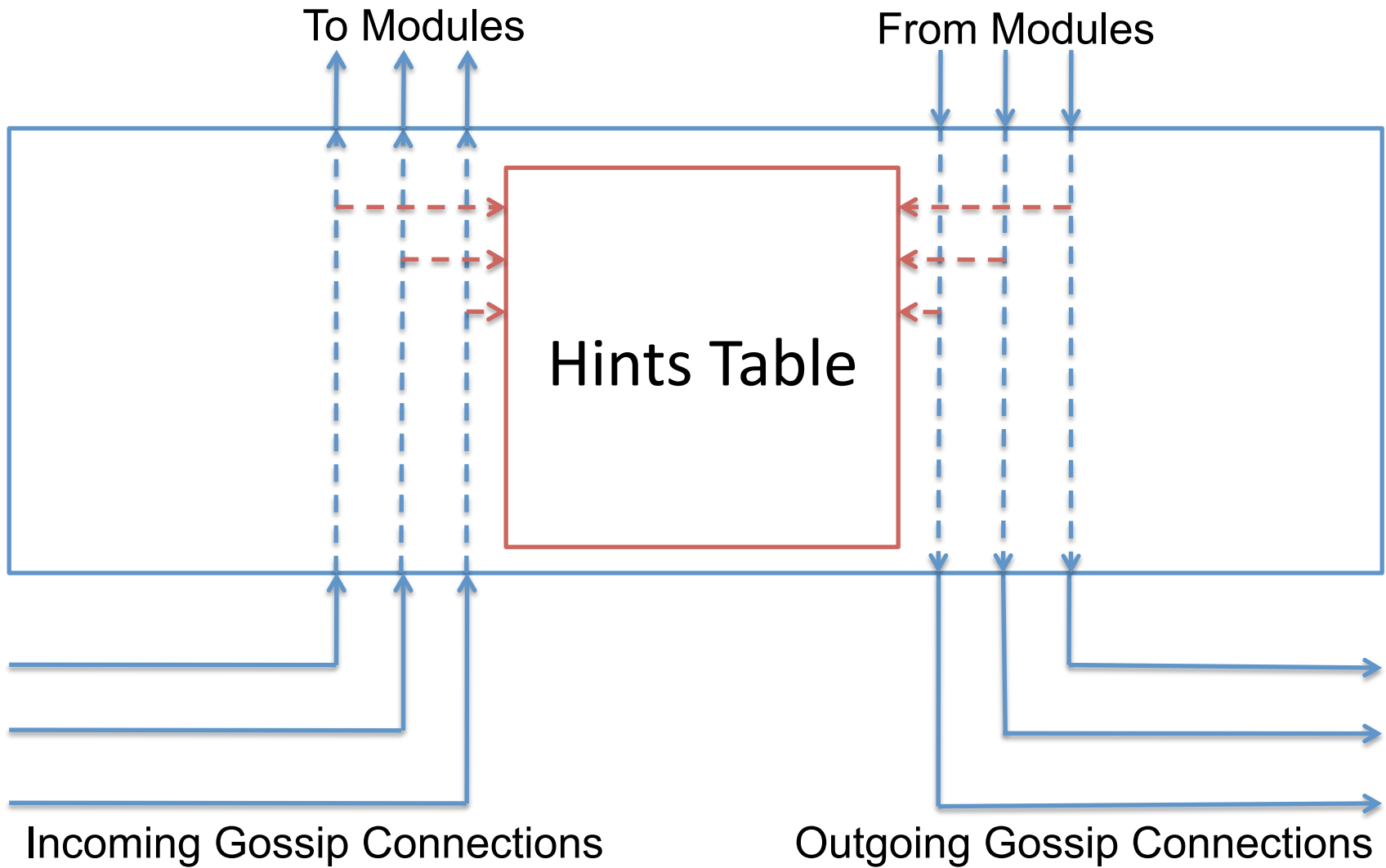
Gossip Mediation

- Core mediates gossip between modules
- Two advantages
 1. Core piggybacks module deployment number on existing gossip traffic which keeps core simple
 2. Core uses HTTP to minimize problems with firewalls

Backup Gossip

- Cores need to be able to update code even if all modules have failed
- Cores implement a rudimentary but robust gossip protocol
 - Static list of rendezvous nodes
 - Intercepted membership hints from module gossip

Core

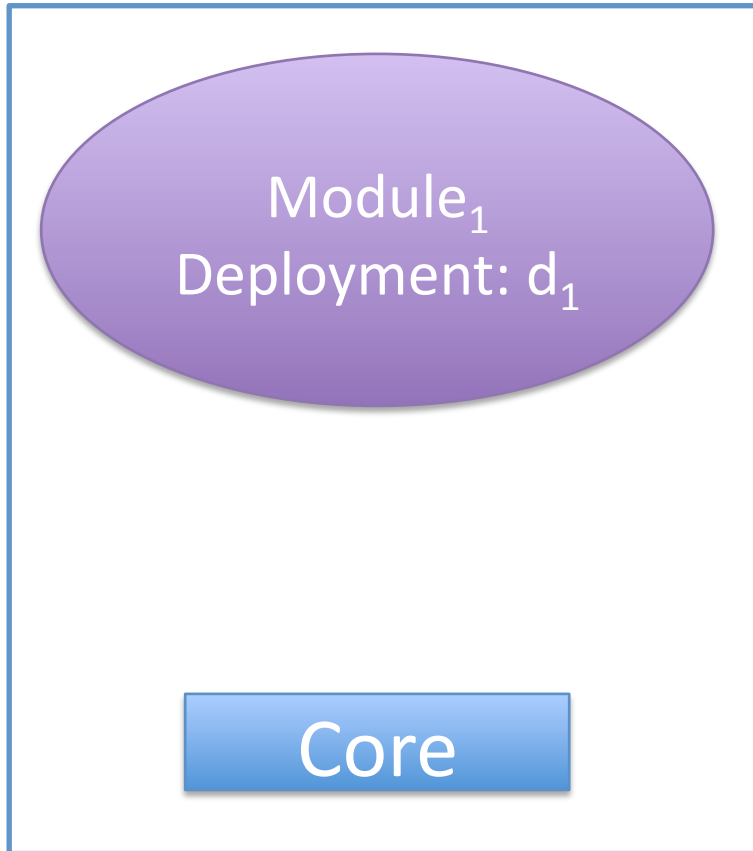


Examples of gossip interactions

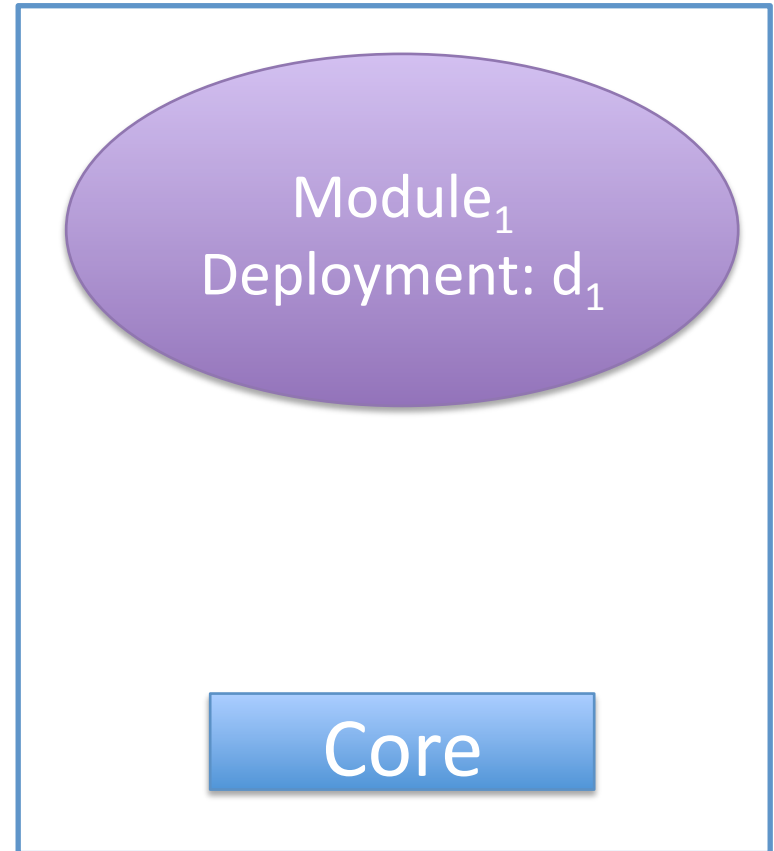
- Normal case: core piggybacks deployment numbers and checks for matched modules
- Mismatched deployment numbers: core initiates code update
- Modules fail to gossip usefully: core gossips configuration information

Normal Case

Node A

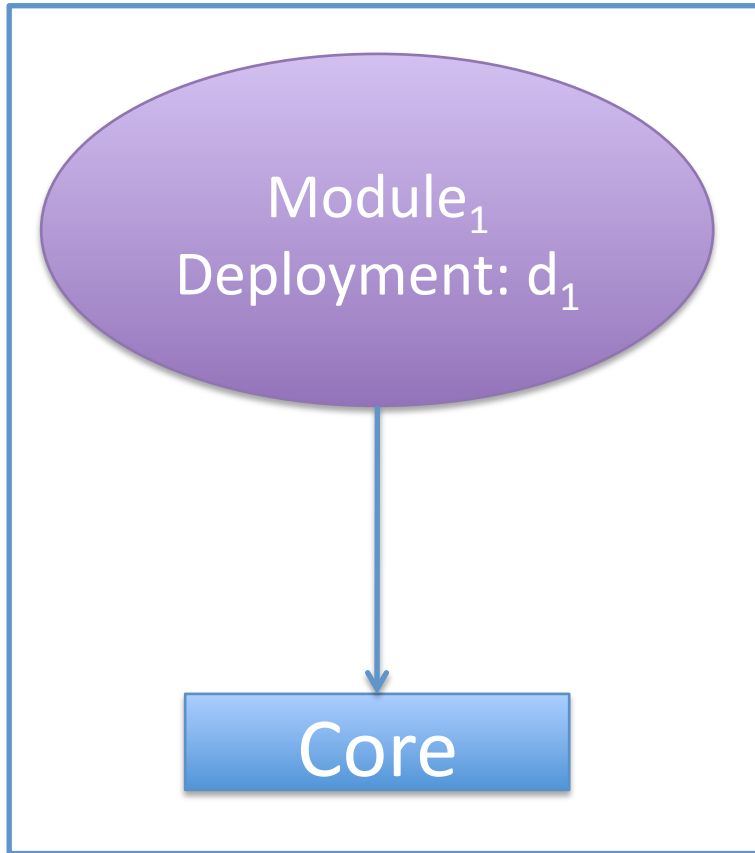


Node B

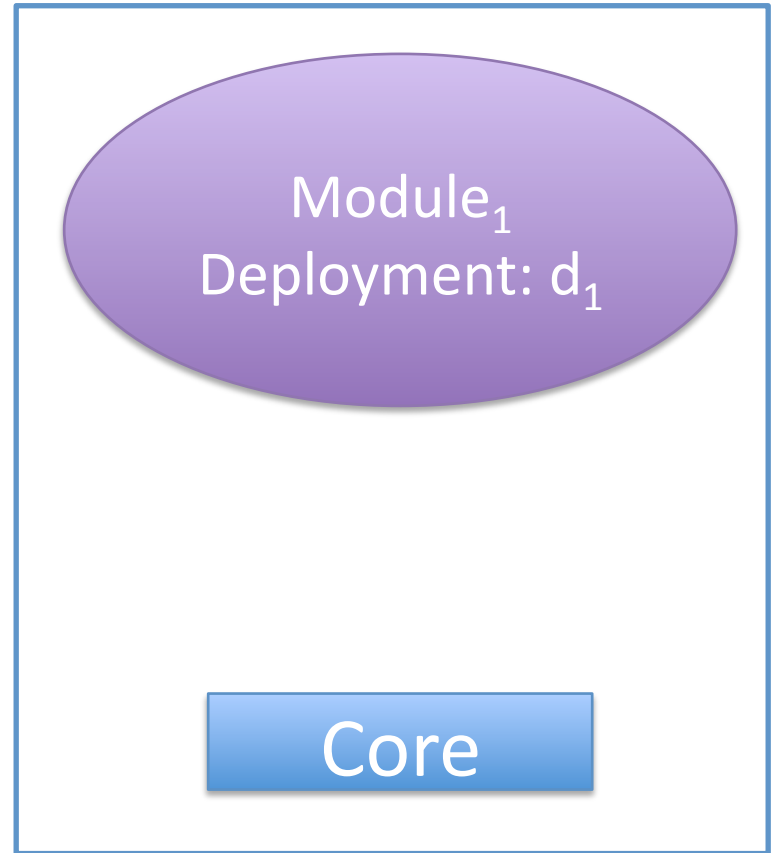


Normal Case

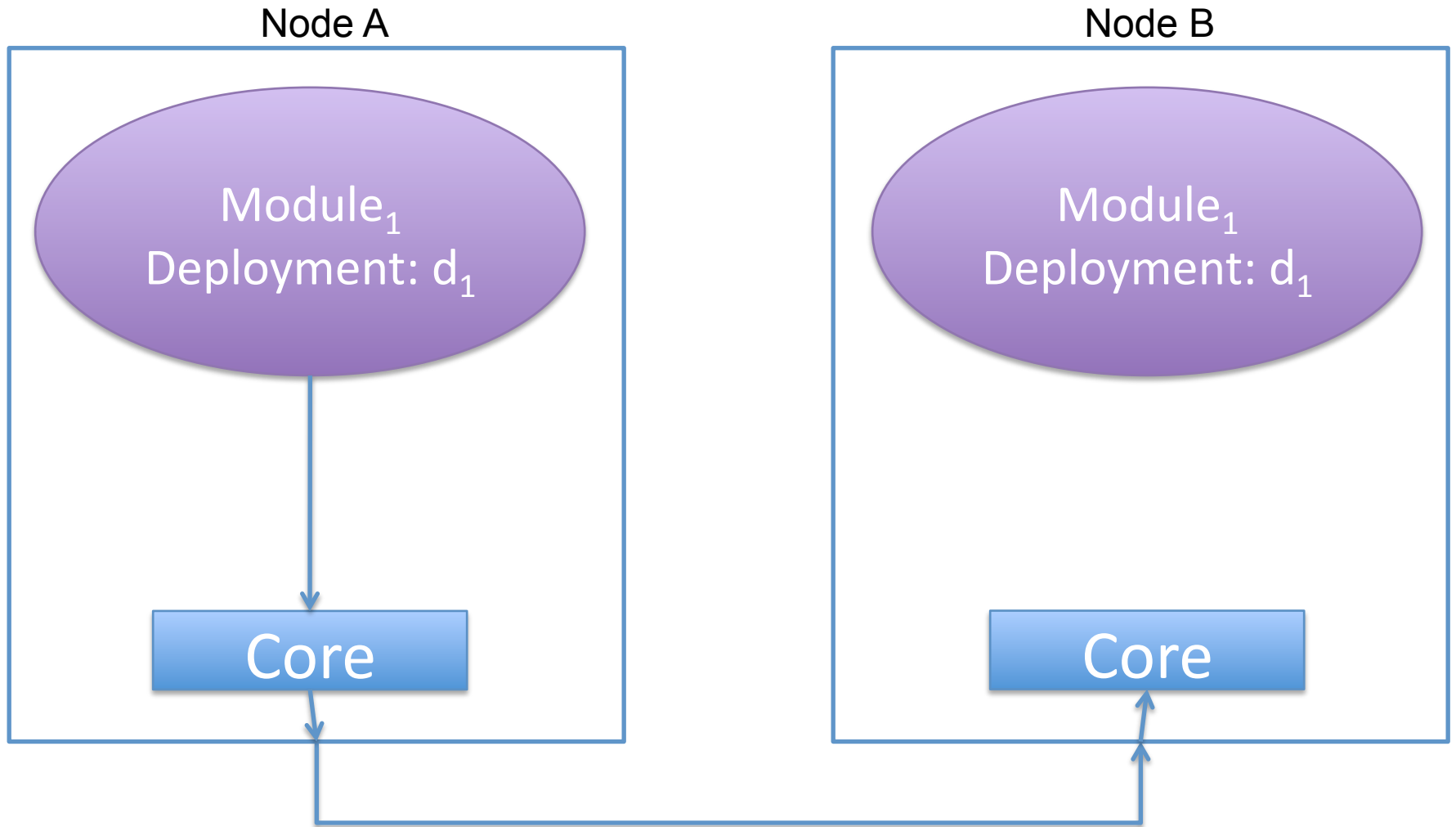
Node A



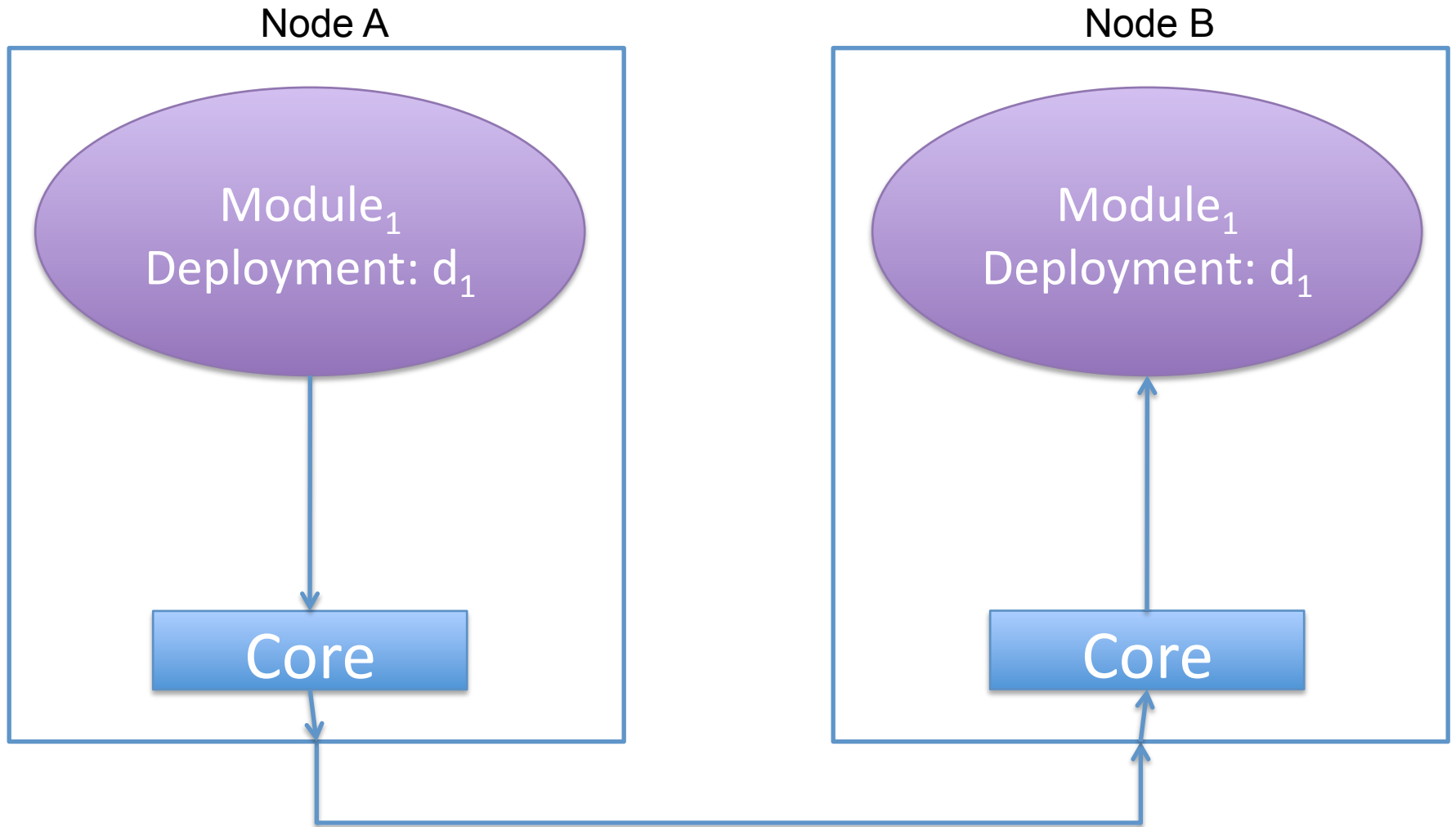
Node B



Normal Case

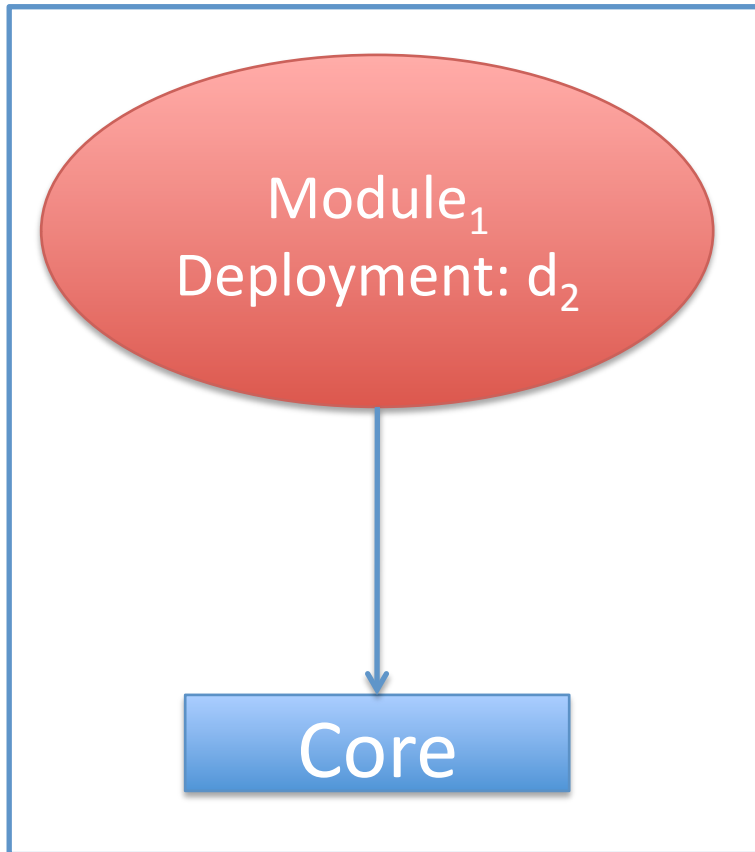


Normal Case

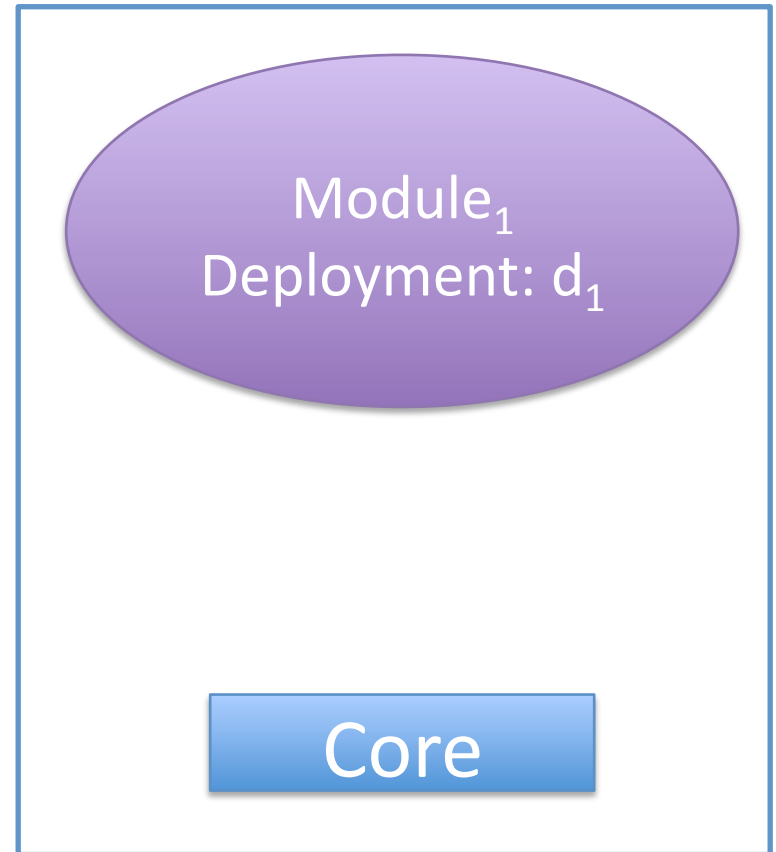


Mismatched Deployment Numbers

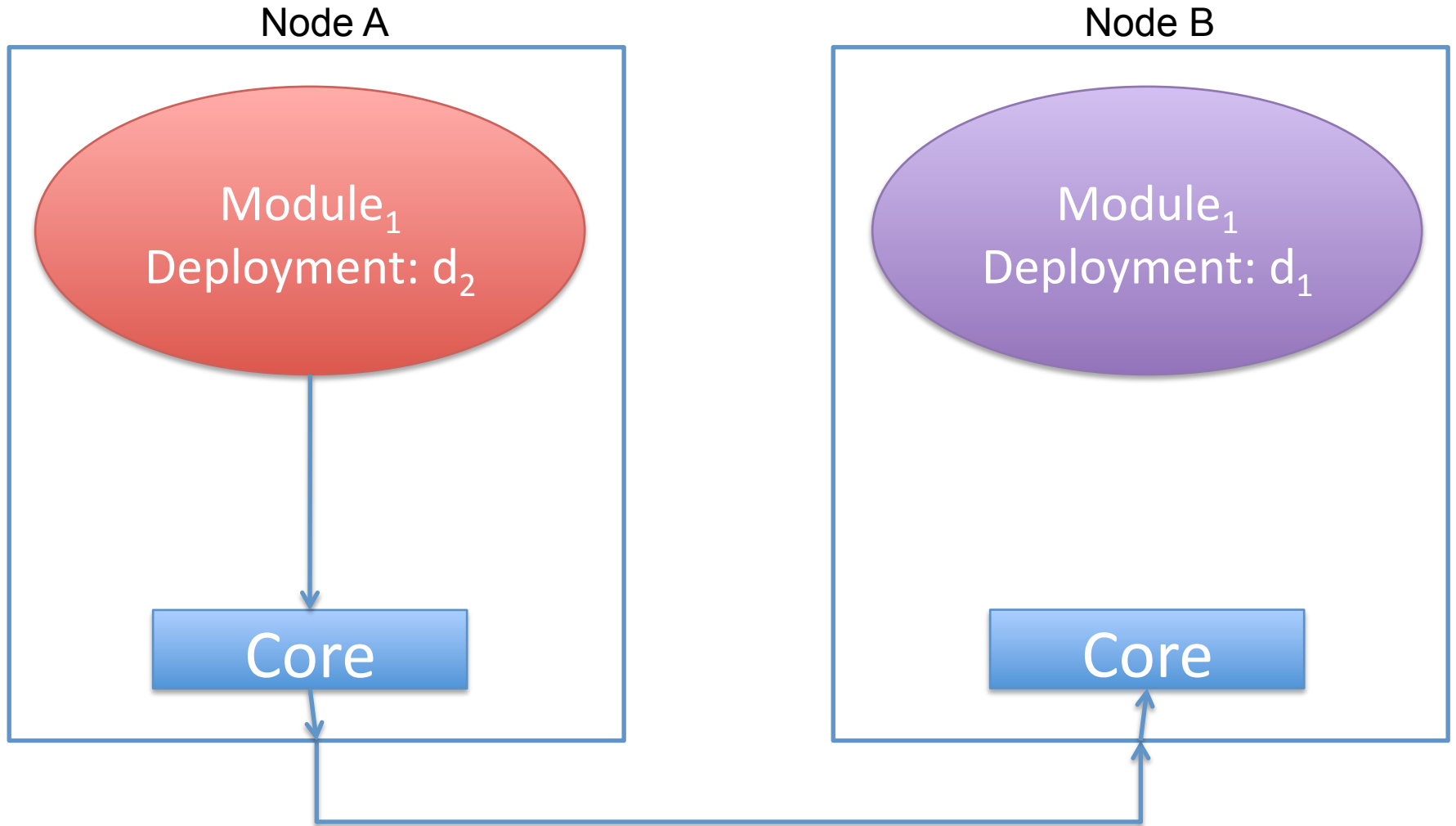
Node A



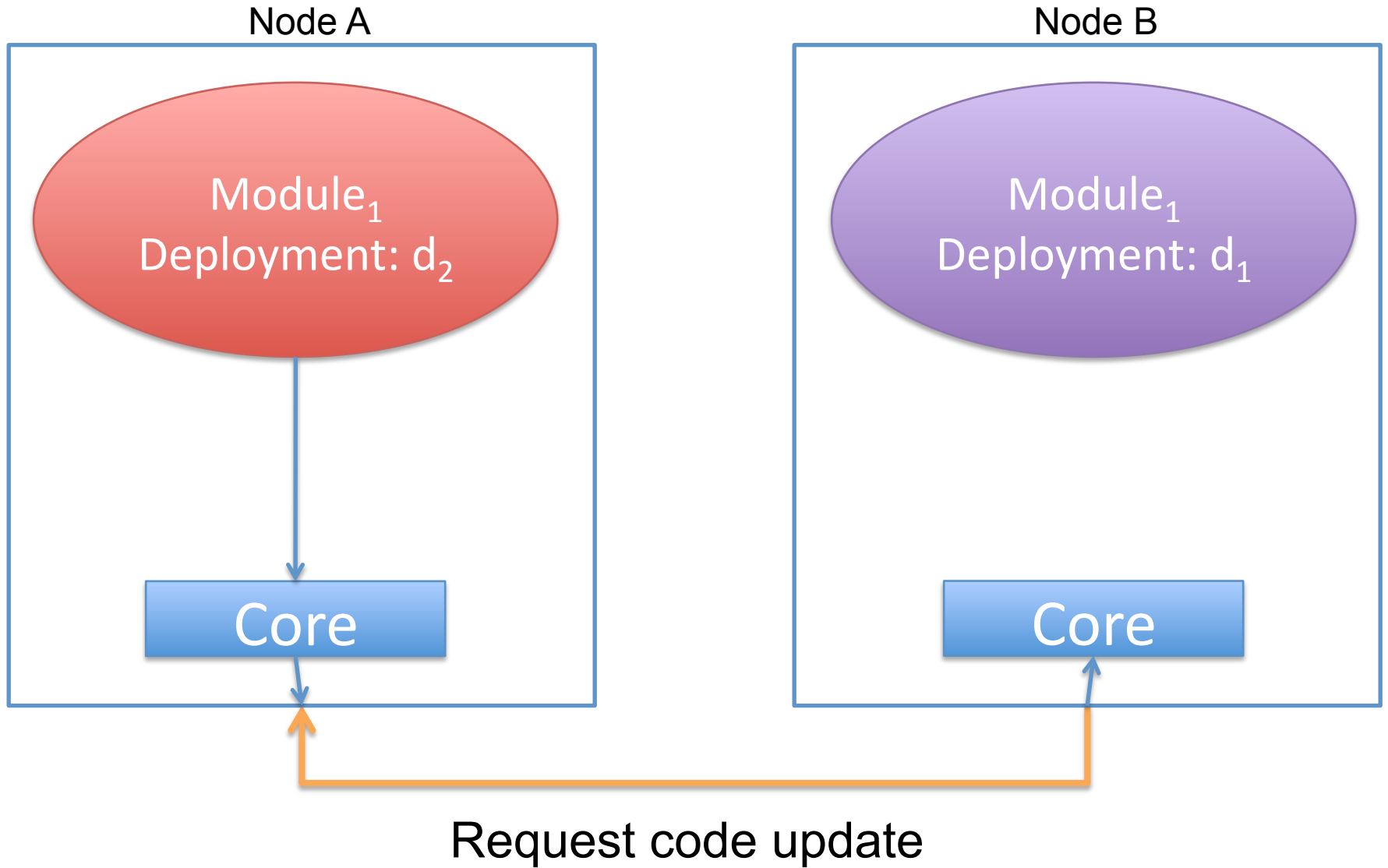
Node B



Mismatched Deployment Numbers

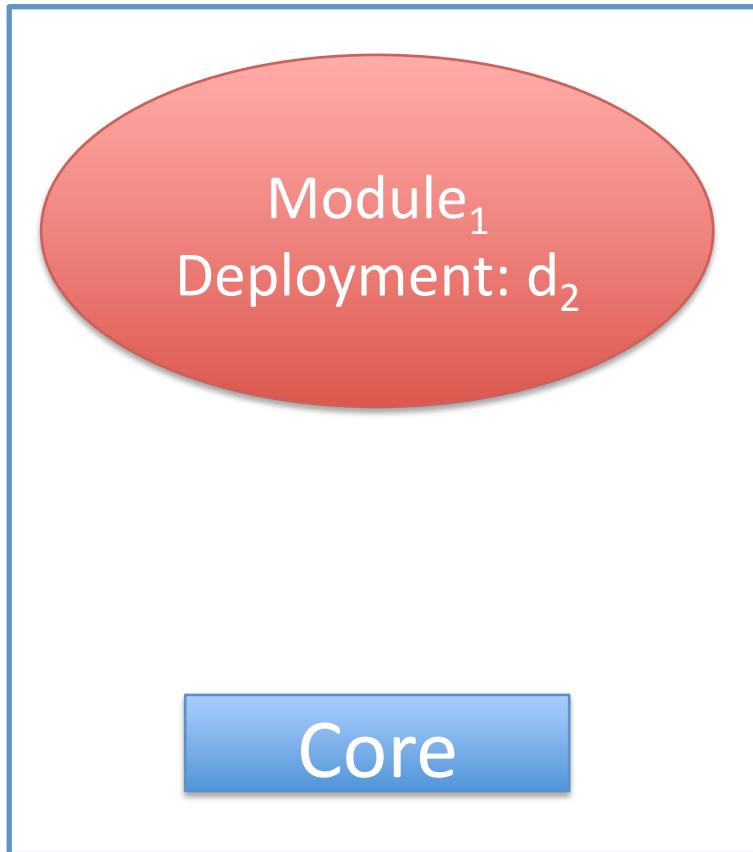


Mismatched Deployment Numbers

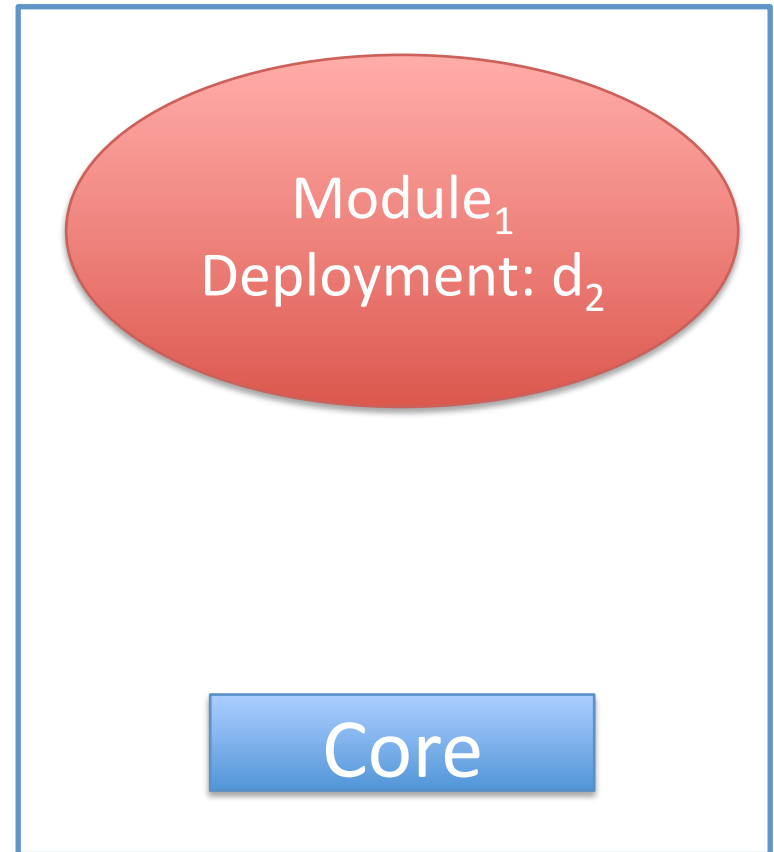


Mismatched Deployment Numbers

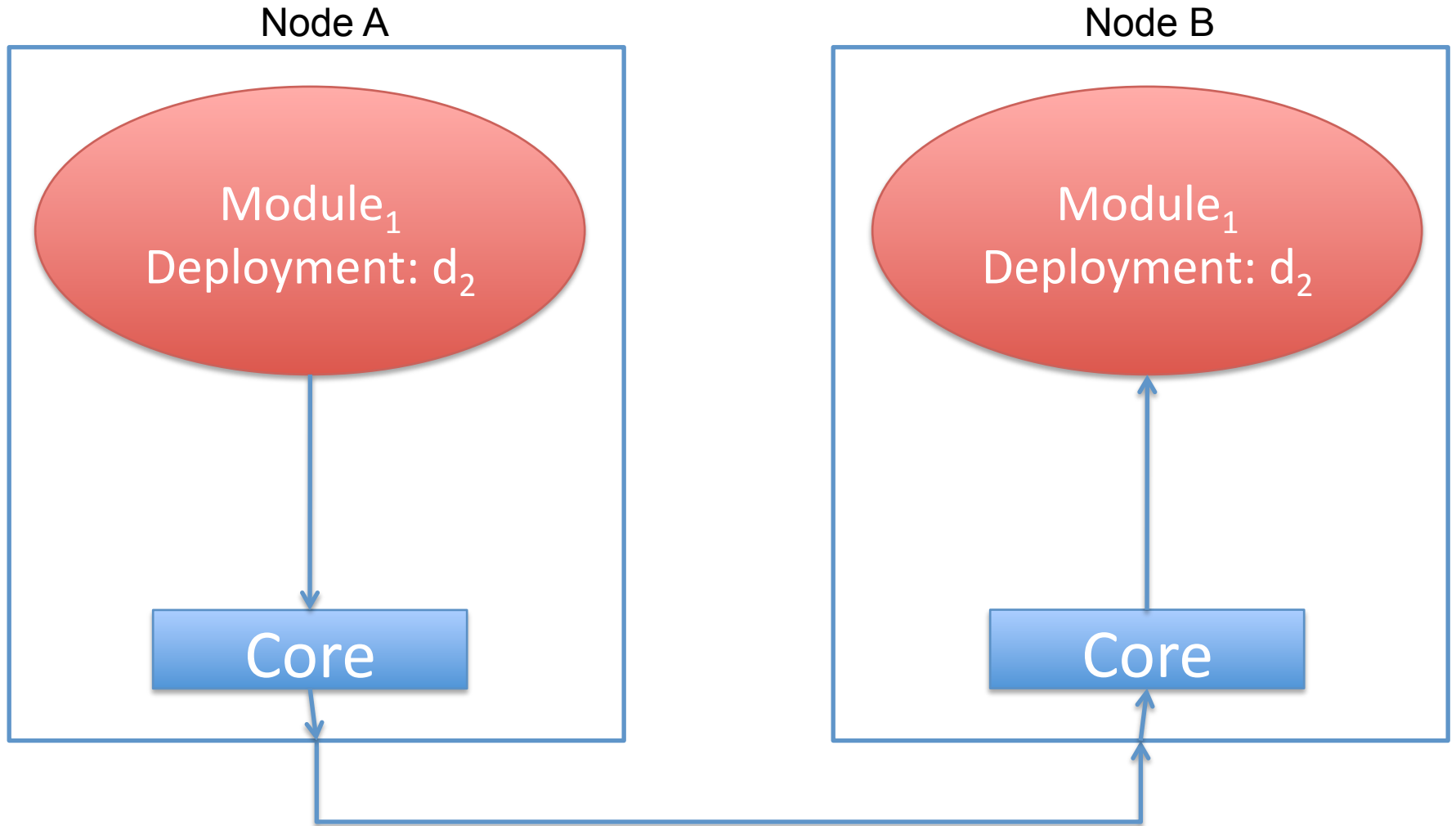
Node A



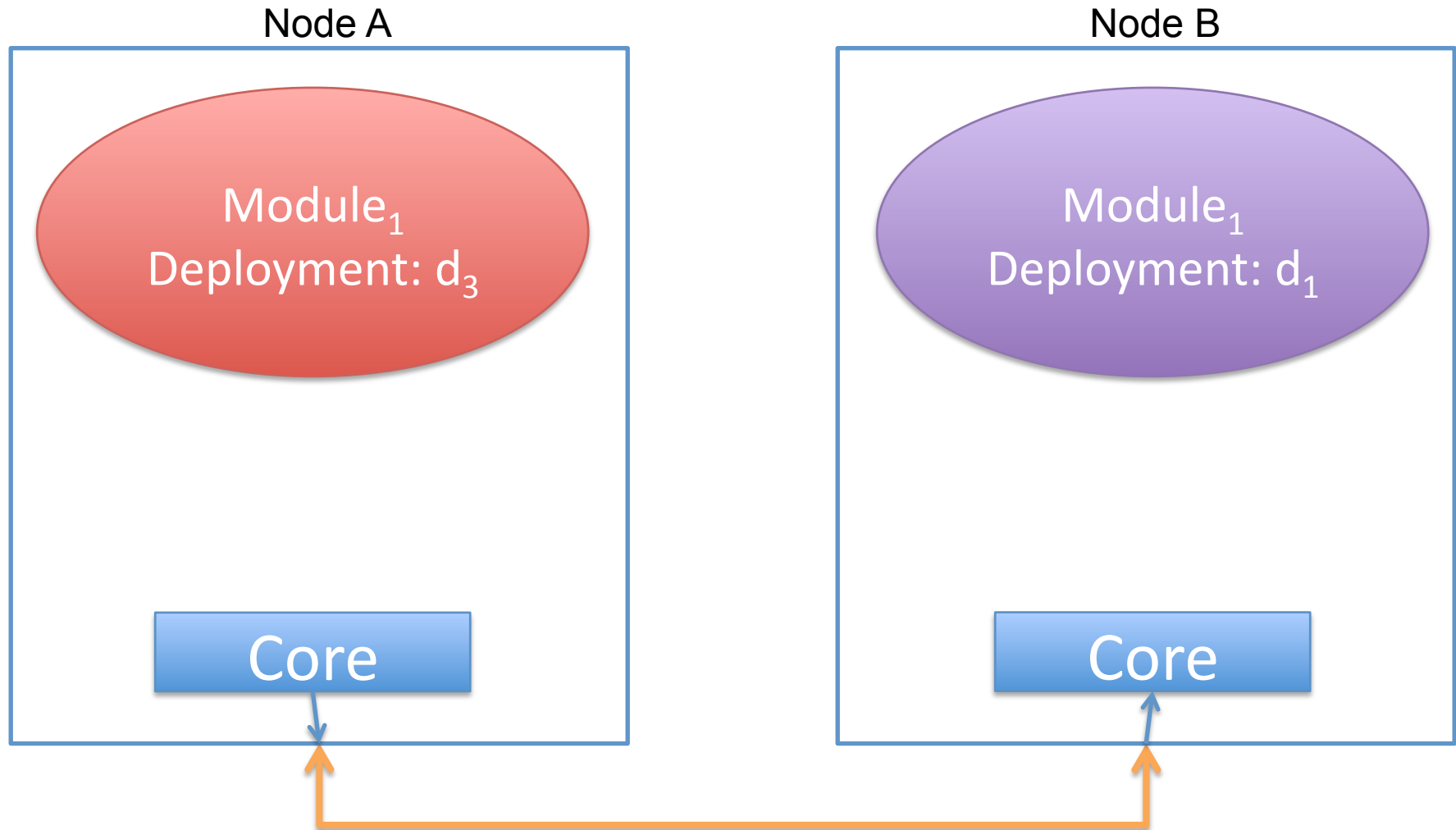
Node B



Mismatched Deployment Numbers

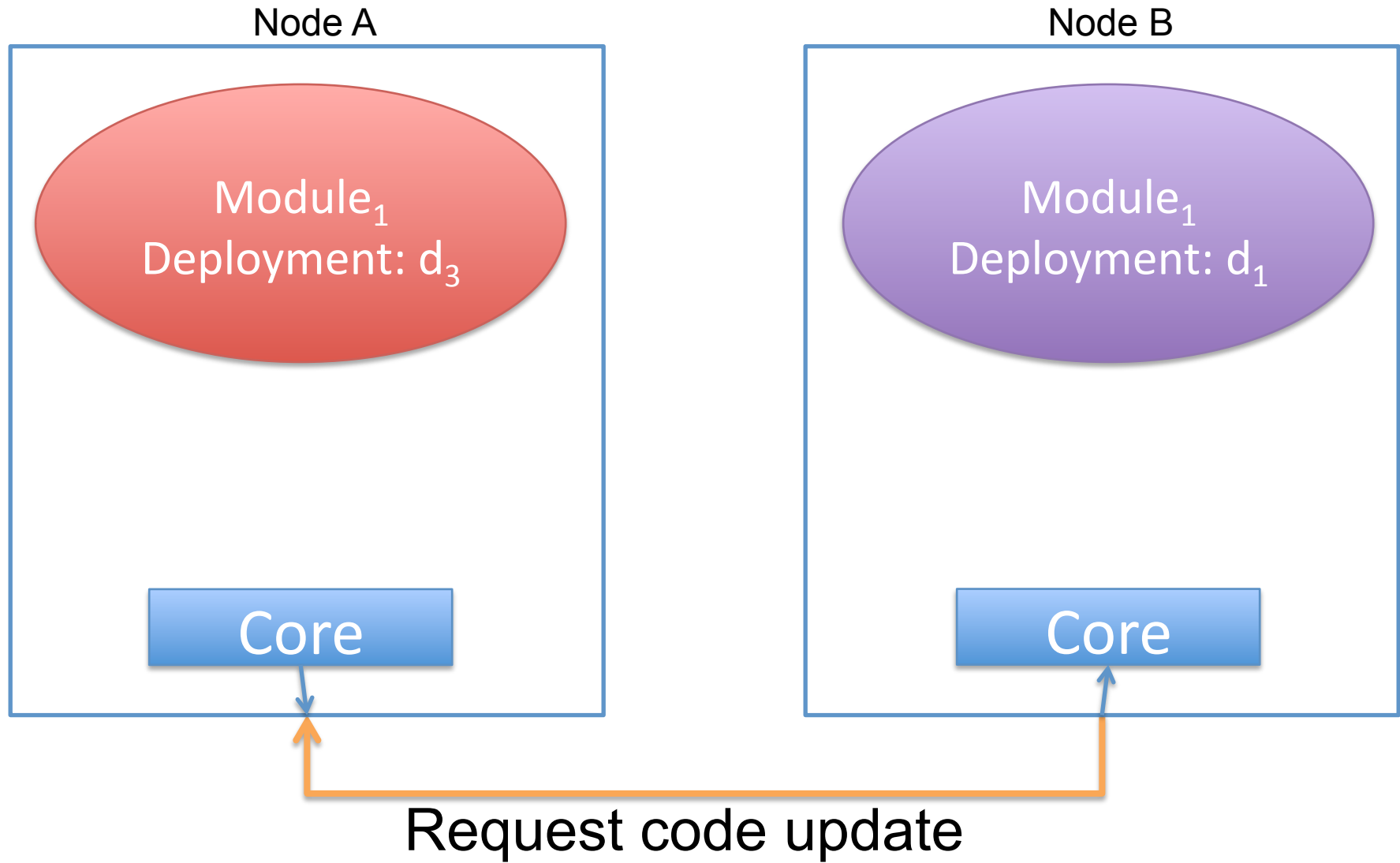


Failure to Gossip usefully

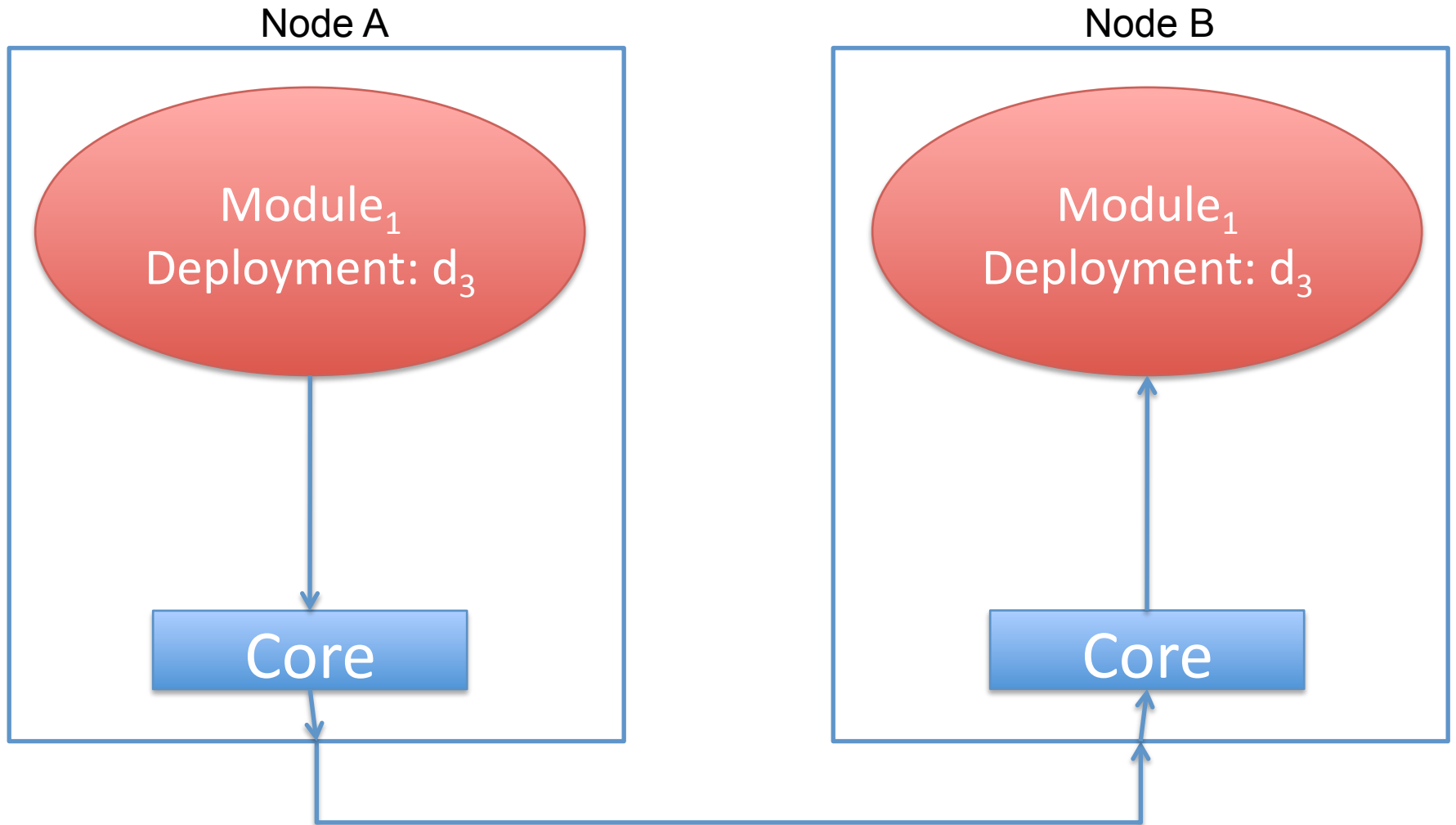


Exchange configuration deployment number

Failure to Gossip usefully



Failure to Gossip usefully



Talk Outline

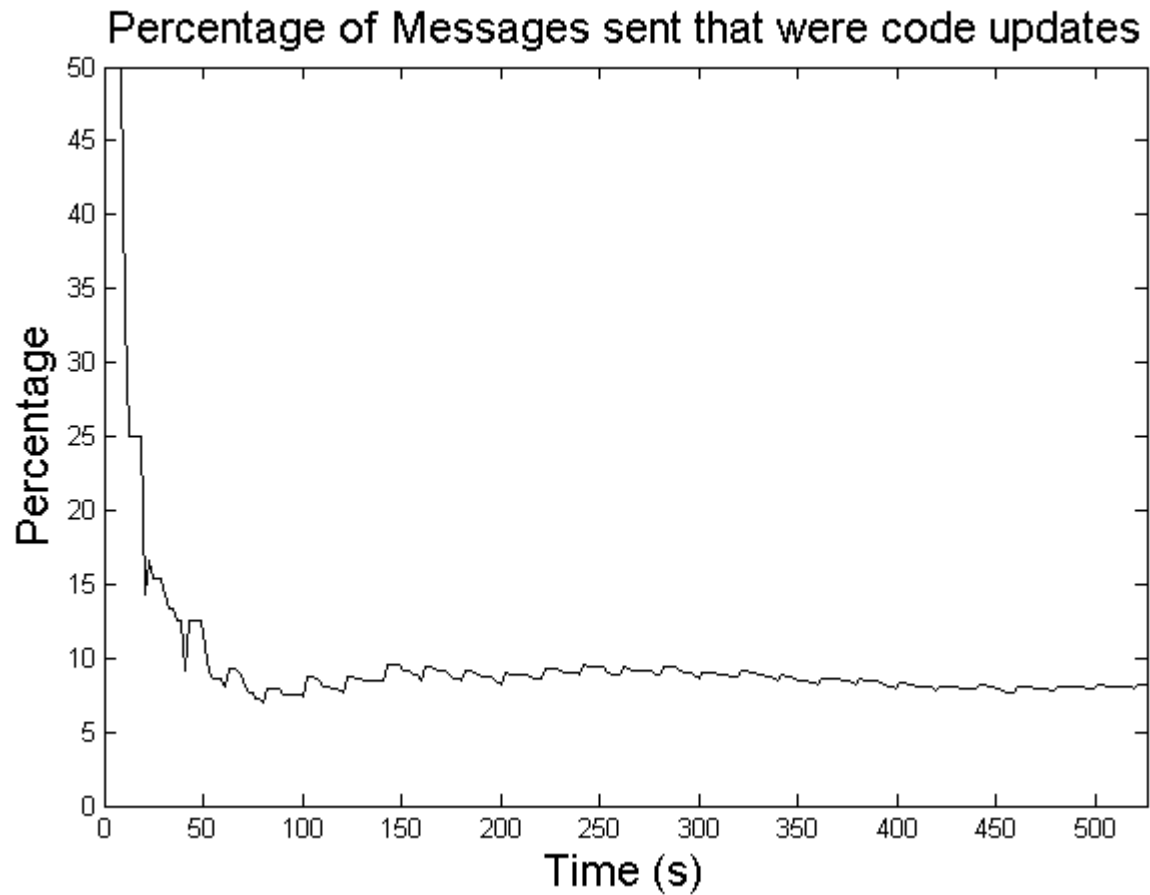
- Code Updating
- Layered Architecture
- Evaluation
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Evaluation

- Tested on 100 local instances with 10 serving as rendezvous servers
- Application: A Simple Membership Protocol

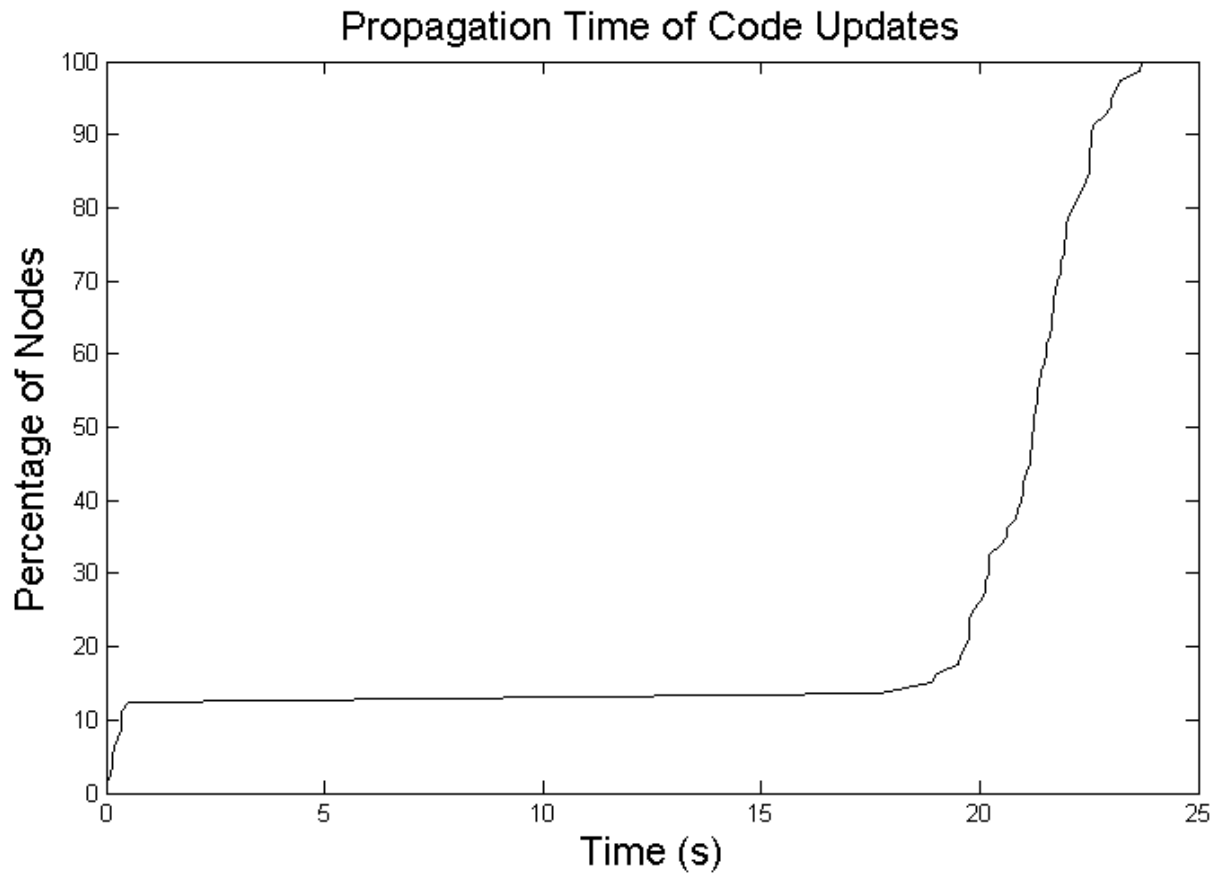
Evaluation

- How much overhead does the core add?



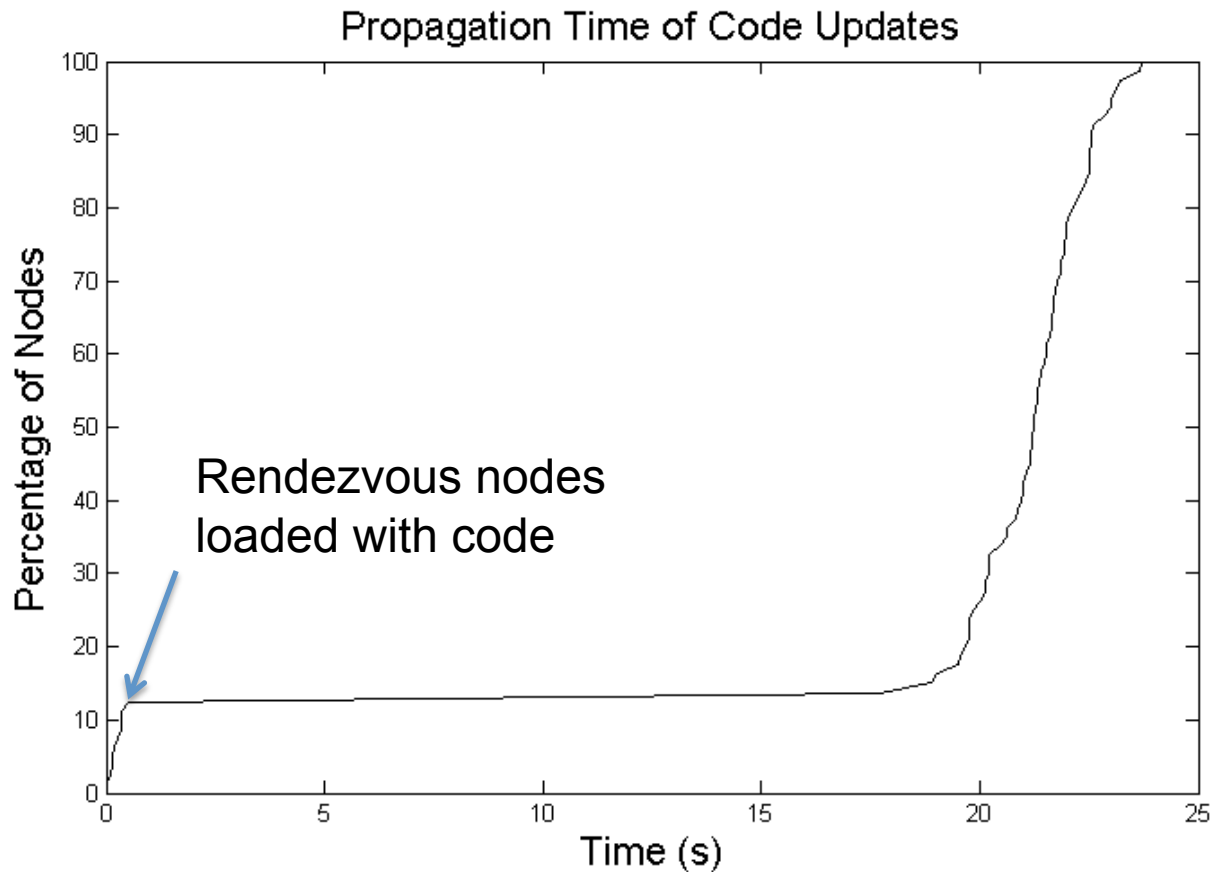
Evaluation

- How long does it take to propagate code?



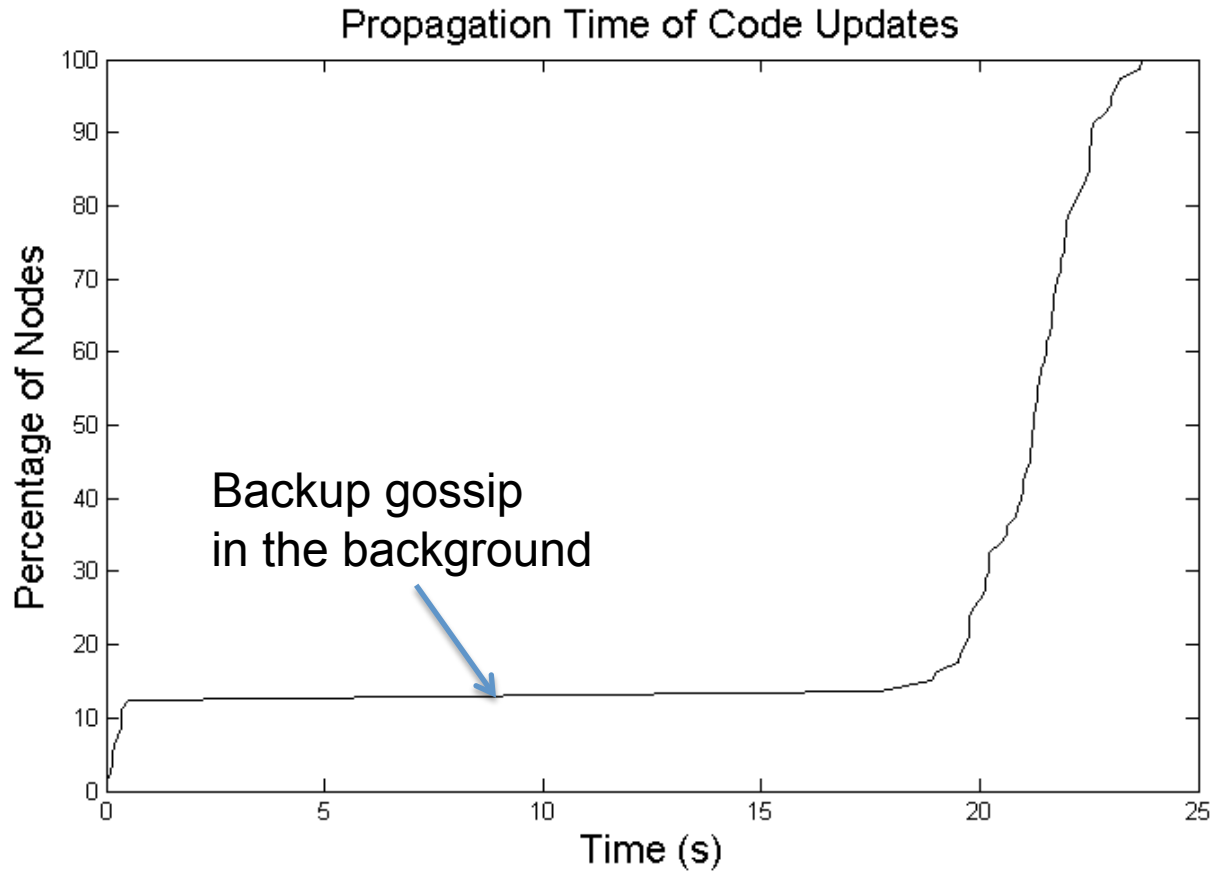
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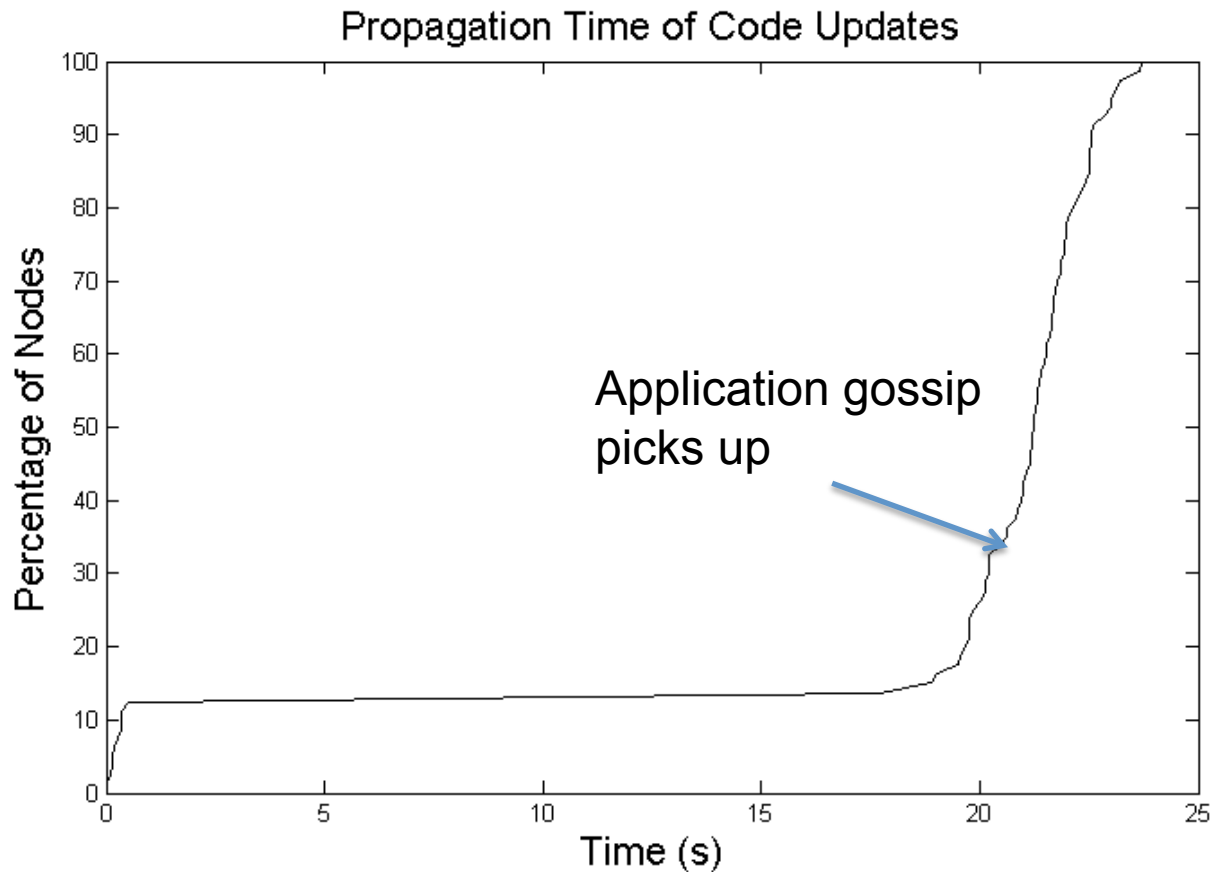
Evaluation

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Evaluation

- How long does it take to propagate code?



Conclusion and Future Work

- Can we make the core smaller?
- Can the core be updated?
- Security
- NAT Traversal as a layered service

Questions?

Module Management

- Core provides the following public methods for module updating:

```
public String transferState()  
public void acceptState()
```

