

Leveraging Value Locality in Optimizing NAND Flash-based SSDs

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Agenda

- **Relook at Locality**
- **Another dimension of Locality : Value Locality**
 - Value Locality and SSDs
- **CA-SSD Design**
 - Mapping Structures
 - Metadata Management
- **Evaluation**
 - CA-SSD vs Traditional SSD

Locality: The pillar of storage

➤ Temporal Locality

- If a logical address is accessed now, it is likely to be accessed again in the near future

➤ Spatial Locality

- If a logical address is accessed now, there is a high likelihood that its neighboring addresses will be accessed in the near future

➤ Pervasive : L1/L2 cache, TLB, Buffer Cache, Virtual Memory, Disk Cache, Web Cache ...

Another Dimension of Locality

➤ Value Locality

- Certain content is accessed preferentially

➤ Data deduplication using Content Addressable Storage (CAS)

➤ Use cases of Value Locality (VL) **Can we use Value Locality to address the idiosyncrasies of SSDs?**

- Network traffic reduction
- Content based Caching
- Efficient data storage (archival/backup)
 - E.g: Venti, Foundation, EMC Centera, Data Domain Storage Systems

CAS suits SSD

SSD

Writes are a
bottleneck

Read/Write
asymmetry

Block
Erases

CAS

Provides
Deduplication

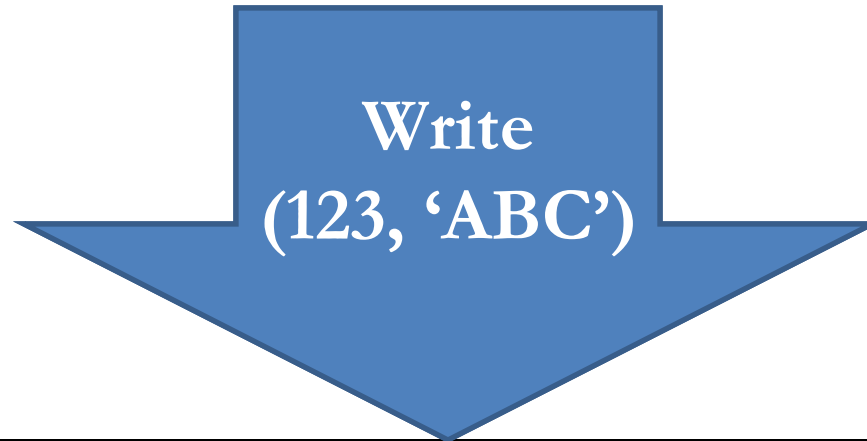
CAS and SSD: Made for each other?

CAS

SSD

**Out of Place
Updates**

Out of Place Updates in CAS



Logical Address	120	121	122	123	124
Translation					
Physical Address	120	121	122	123	124
Storage	ABC	DEF	PQR	TUV	XYZ

The diagram illustrates the mapping of logical addresses to physical addresses and storage. A write operation is performed at logical address 123 with the value 'ABC'. The translation process maps logical address 123 to physical address 120. The storage at physical address 120 is updated with 'ABC', which is an out-of-place update because the value is not at the physical address corresponding to the logical address (123).

CAS and SSD: Made for each other?

CAS

**Out of Place
Updates**

SSD

**Erase before
Write**

Problem with CAS

CAS

Loss of
Sequentiality

SSD

Fast Random
Reads

**Do real workloads exhibit
Value Locality?**

Workloads [Koller10]

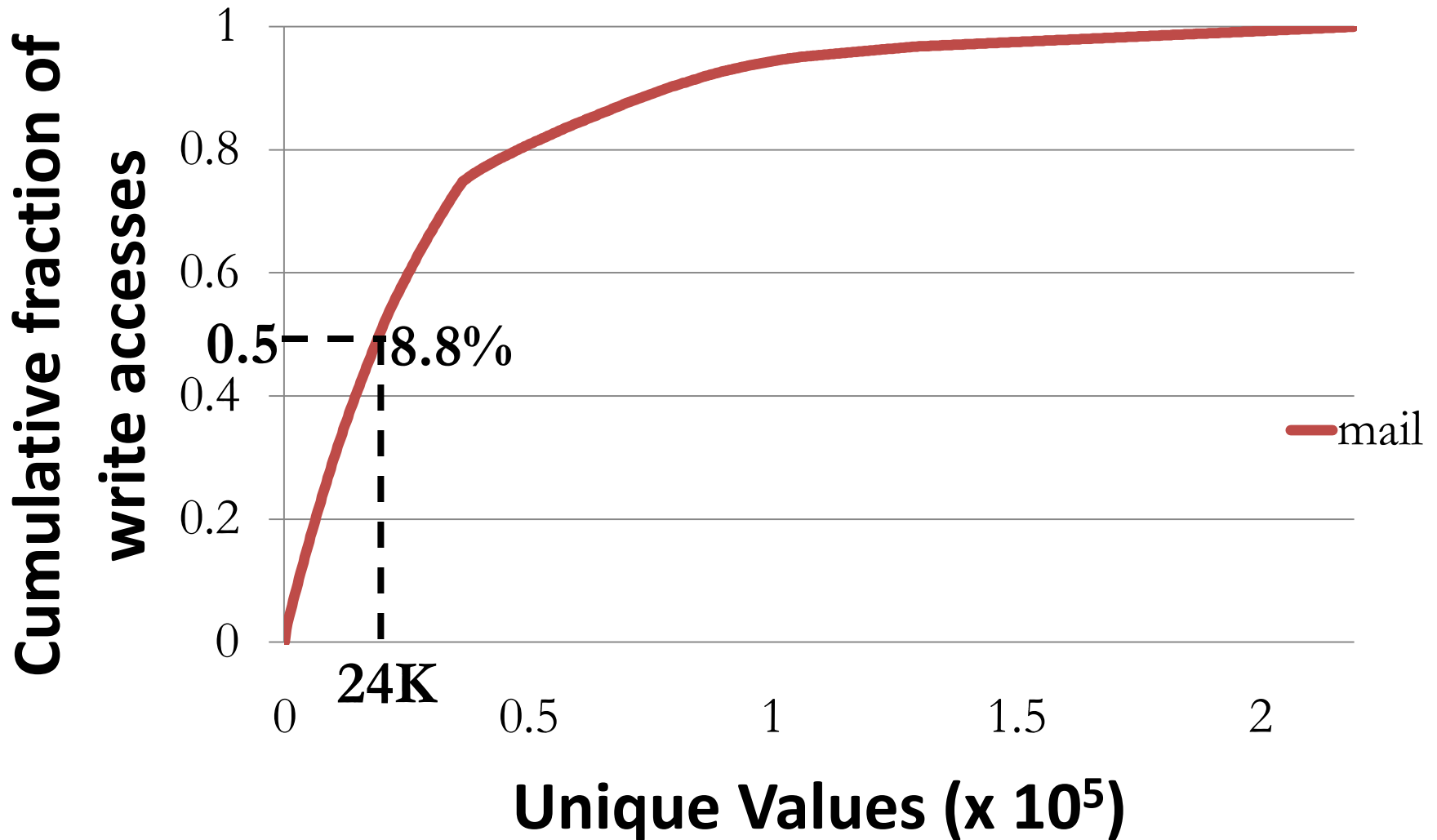
Workload	Writes (%)	Total Requests (Millions)	Unique Write Requests (%)	Unique Read Requests (%)
Write dominant				
web	77.0	Duplication	42.35	32.05
mail	77.3	3.6	7.83	80.85
homes	96.7	4.4	66.37	80.75

[Koller10] Koller, R., and Rangaswami, R. "I/O Deduplication: Utilizing Content Similarity to Improve I/O Performance." (*FAST'10*)

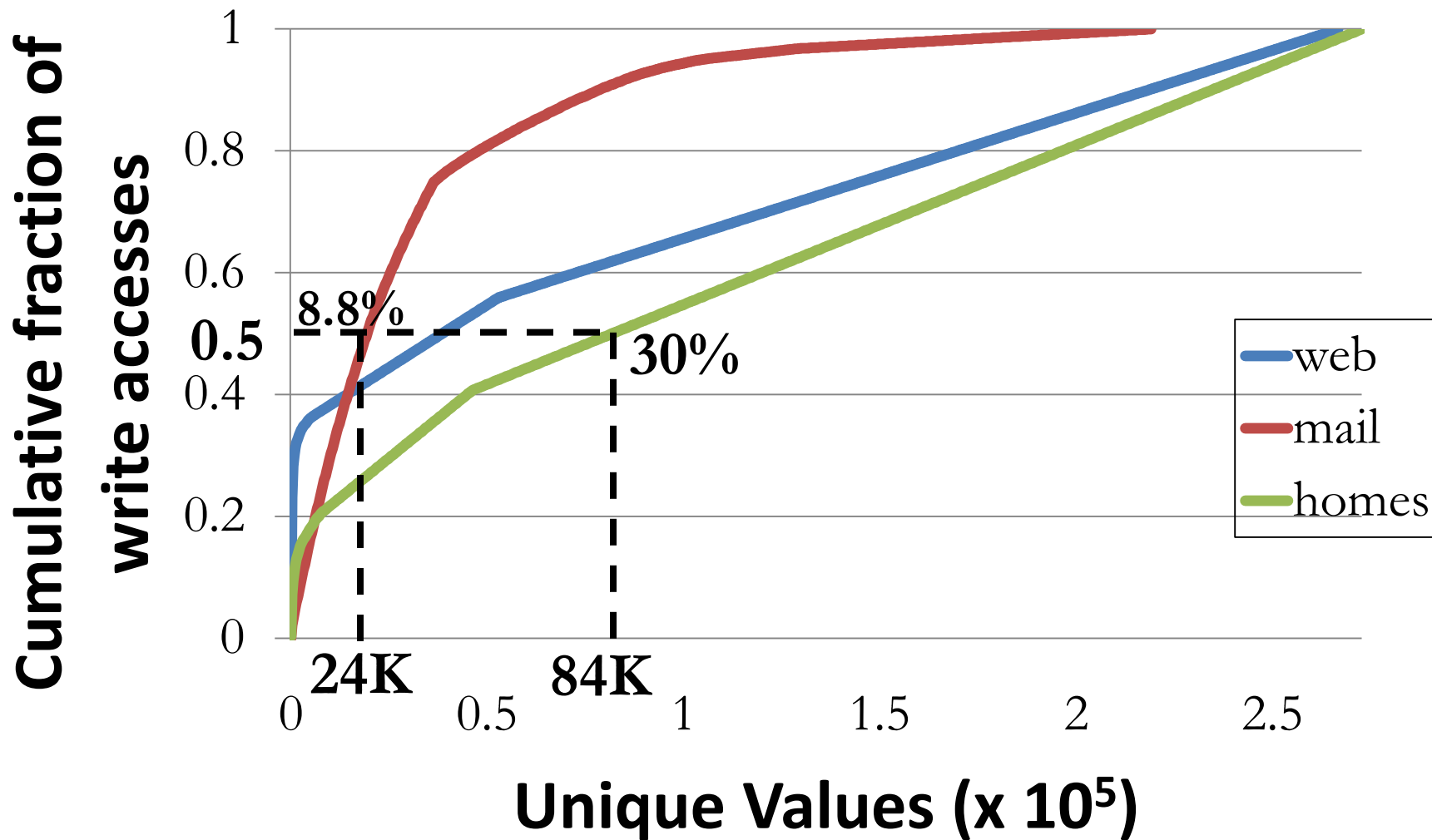
Value Popularity

- **VP** represents the number of occurrences of each unique value in a workload
 - Signifies potential for deduplication for a workload

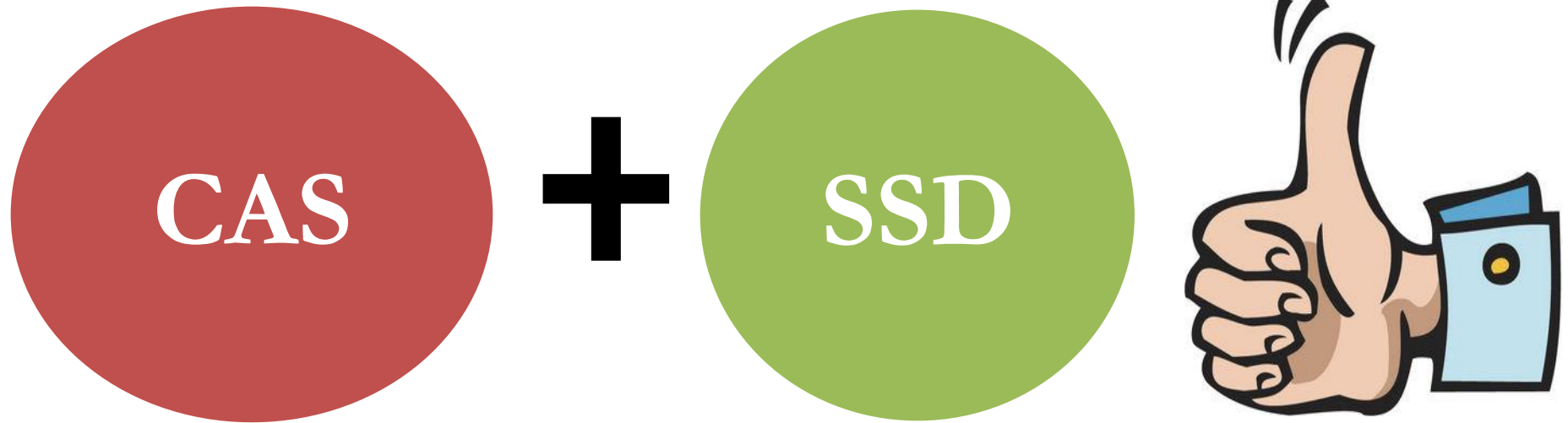
Some Values are Very Popular



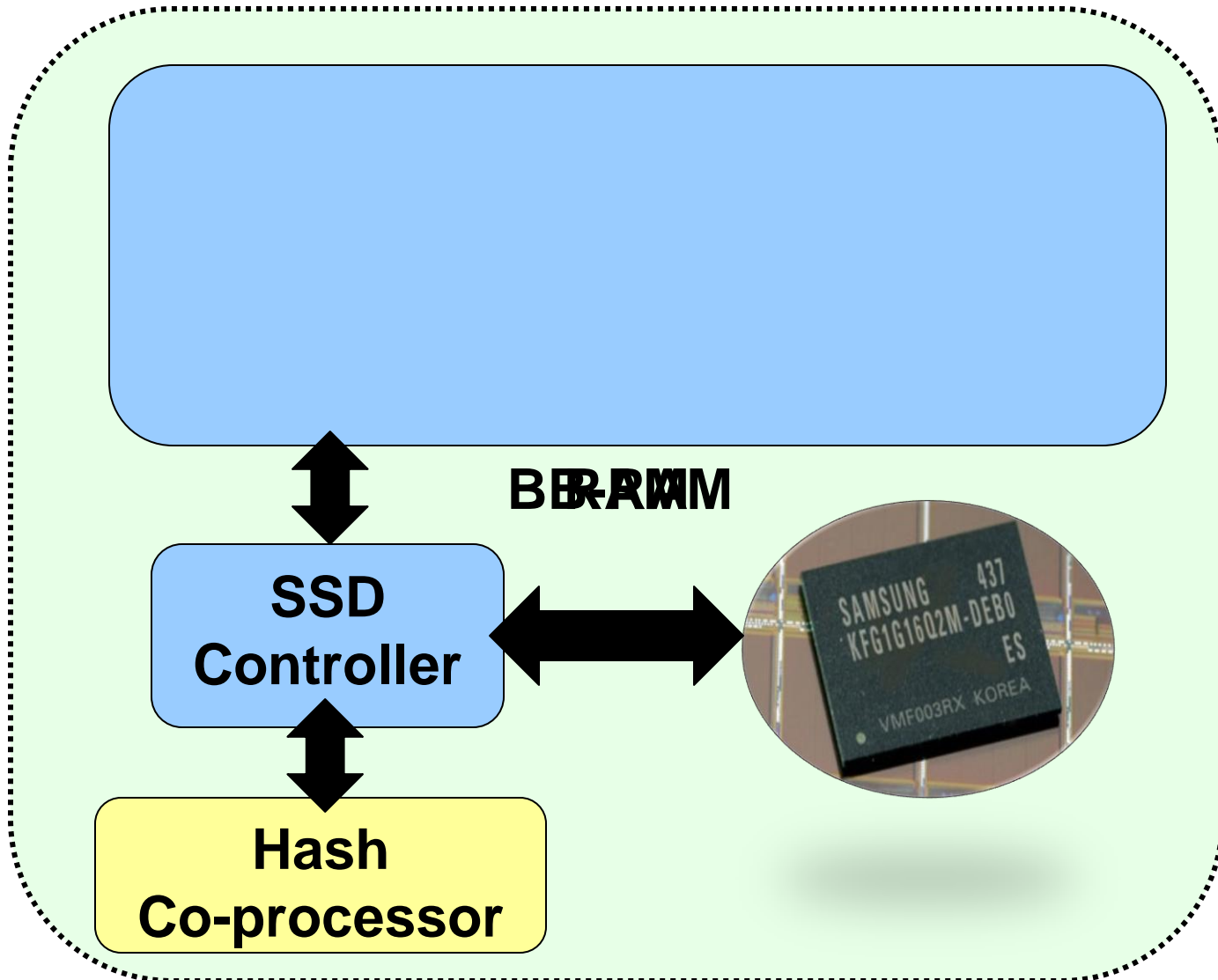
Some Values are Very Popular



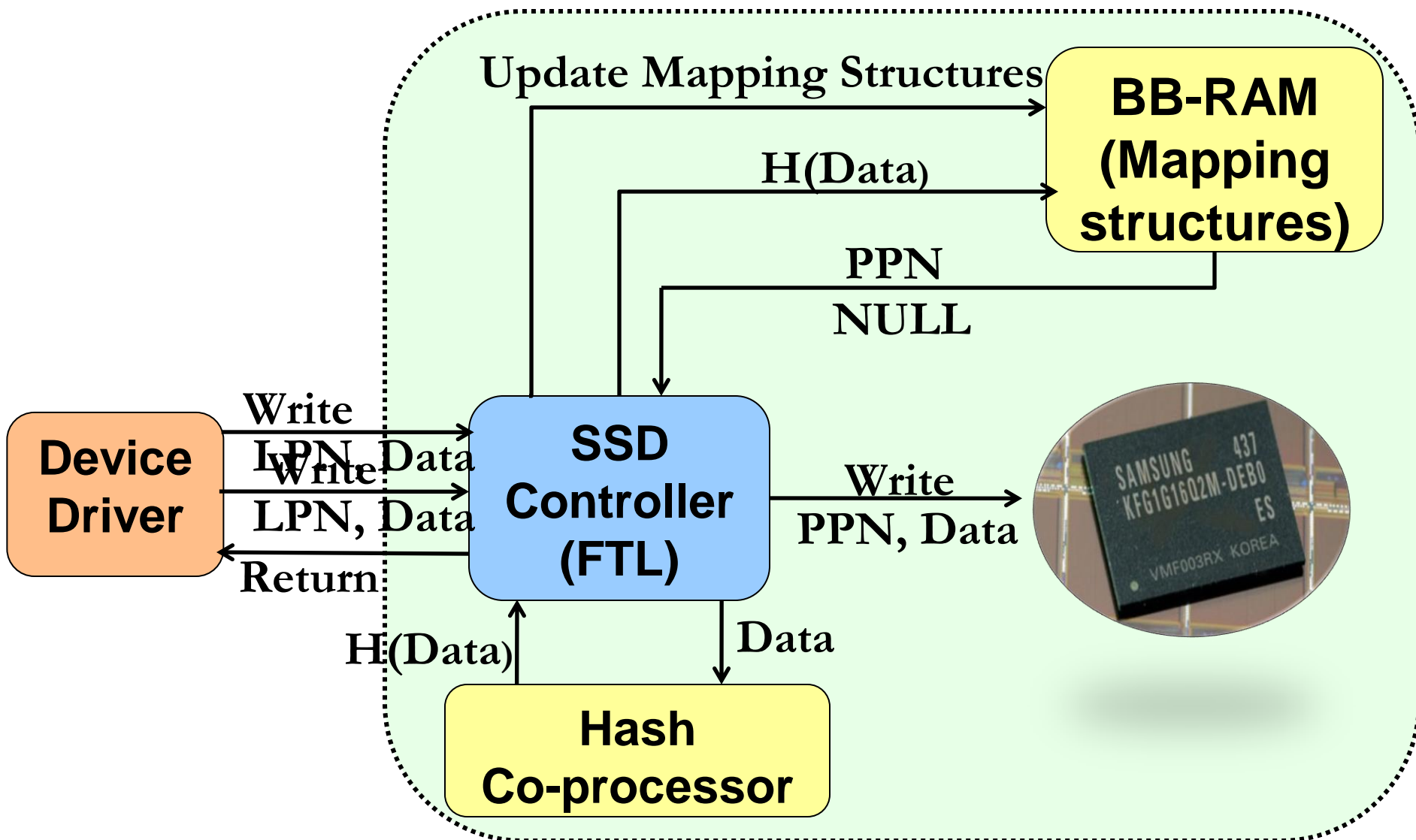
CA-SSD



CA-SSD Design



CA-SSD Design

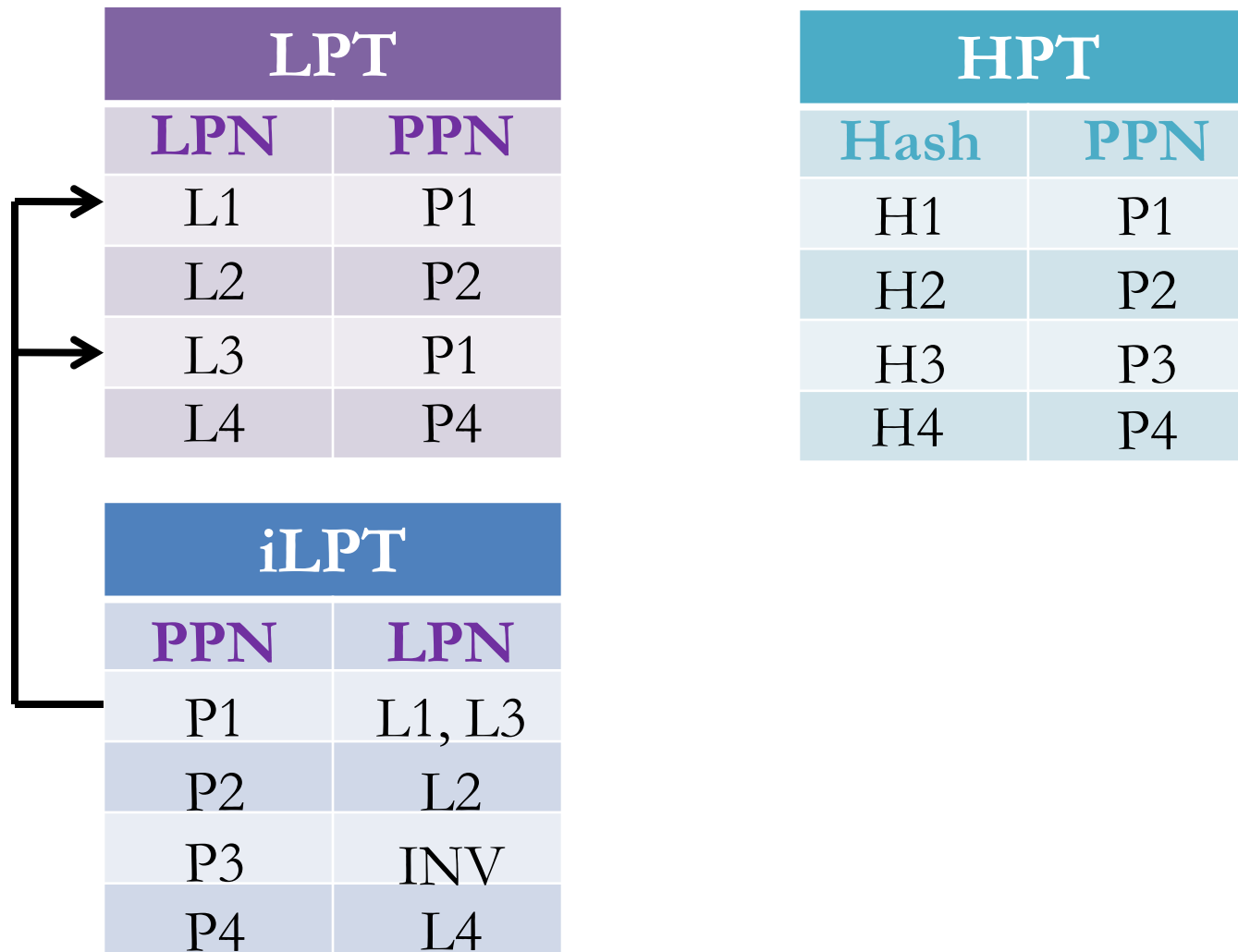


Mapping Structures: LPT & HPT

LPT	
LPN	PPN
L1	P1
L2	P2
L3	P1
L4	P4

HPT	
Hash	PPN
H1	P1
H2	P2
H3	P3
H4	P4

Mapping Structures: iLPT



Mapping Structures: iLPT & iHPT

LPT	
LPN	PPN
L1	P1
L2	P2
L3	P1
L4	P4

HPT	
Hash	PPN
H1	P1
H2	P2
H3	P3
H4	P4

iLPT	
PPN	LPN
P1	L1, L3
P2	L2
P3	INV
P4	L4

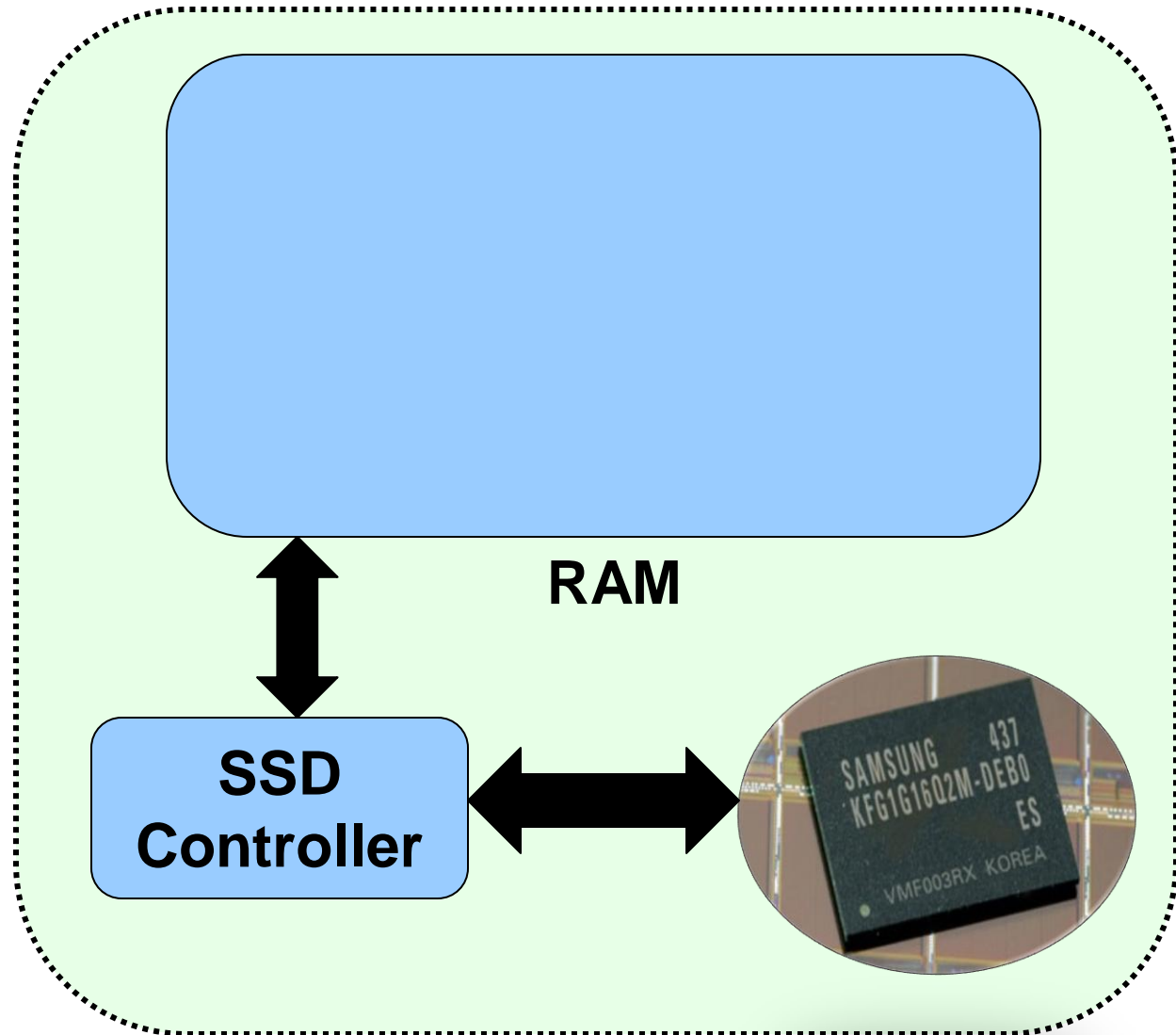
iHPT	
PPN	Hash
P1	H1
P2	H2
P3	H3
P4	H4

Remove



Metadata: Traditional SSD

LPT	
LPN	PPN
L1	P1
L2	P2
L3	P3
L4	P4



Motivation for Large SRAM



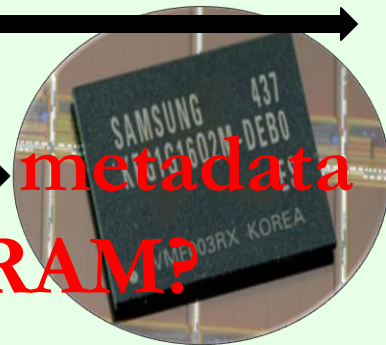
HPT	
Hash	PPN
H1	P1
H2	P2
H3	P3
H4	P4

LPT	
LPN	PPN
L1	P1
L2	P2
L3	P1
L4	P4

iLPT	
PPN	LPN
P1	L1,L3
P2	L2
P3	INV
P4	L4

iHPT	
PPN	Hash
P1	H1
P2	H2
P3	H3
P4	H4

BB-RAM
 SSD Controller
How do we fit the metadata in CA-SSD's RAM?
Not Scalable!!



Hash
 Co-processor

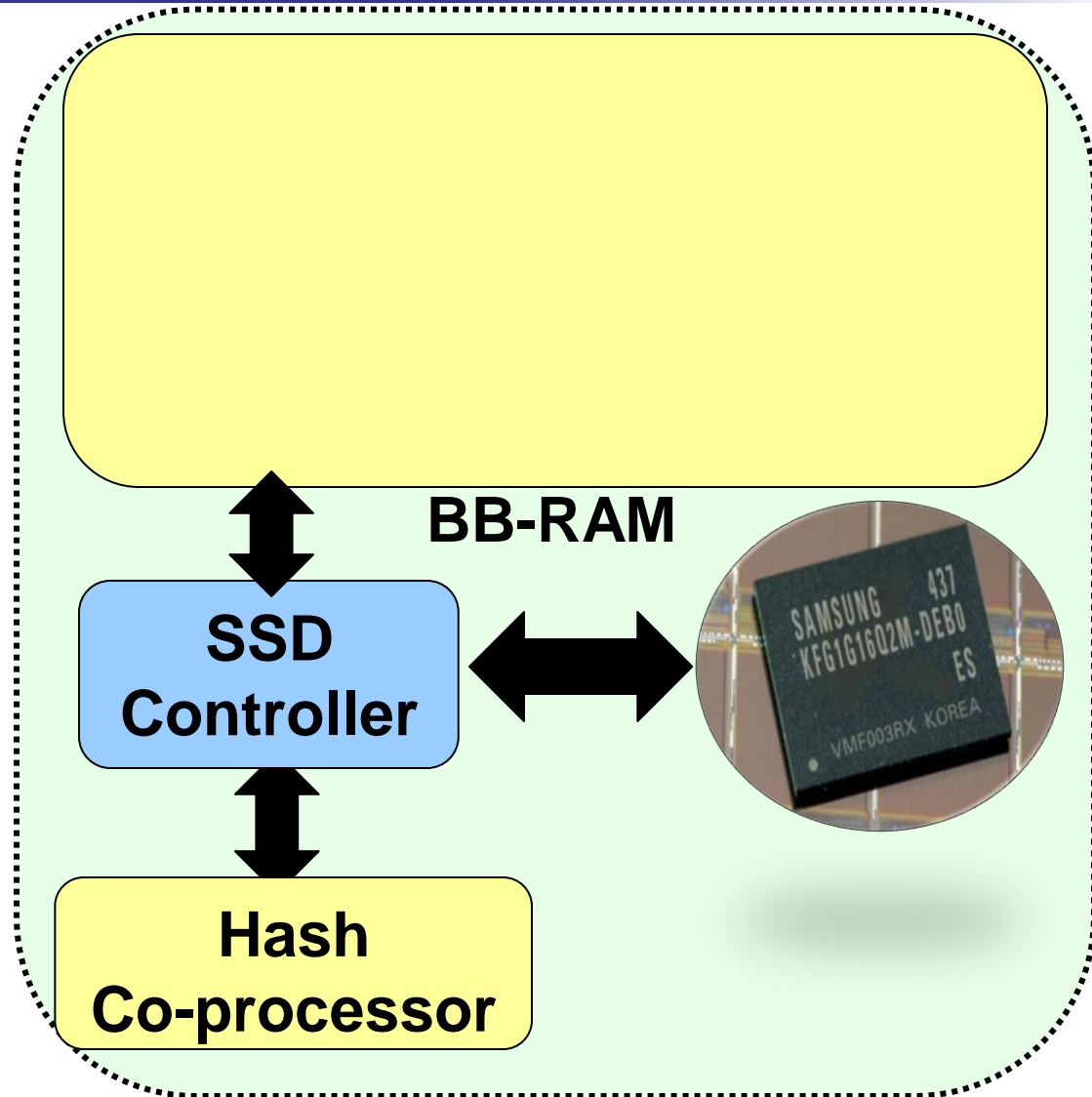
Option 2 : Shrink Metadata

LPT	
LPN	PPN
L1	P1
L2	P2
L3	P1
L4	P4

HPT	
Hash	PPN
H1	P1
H2	P2
H3	P3
H4	P4

iLPT	
PPN	LPN
P1	L1,L3
P2	L2
P3	INV
P4	L4

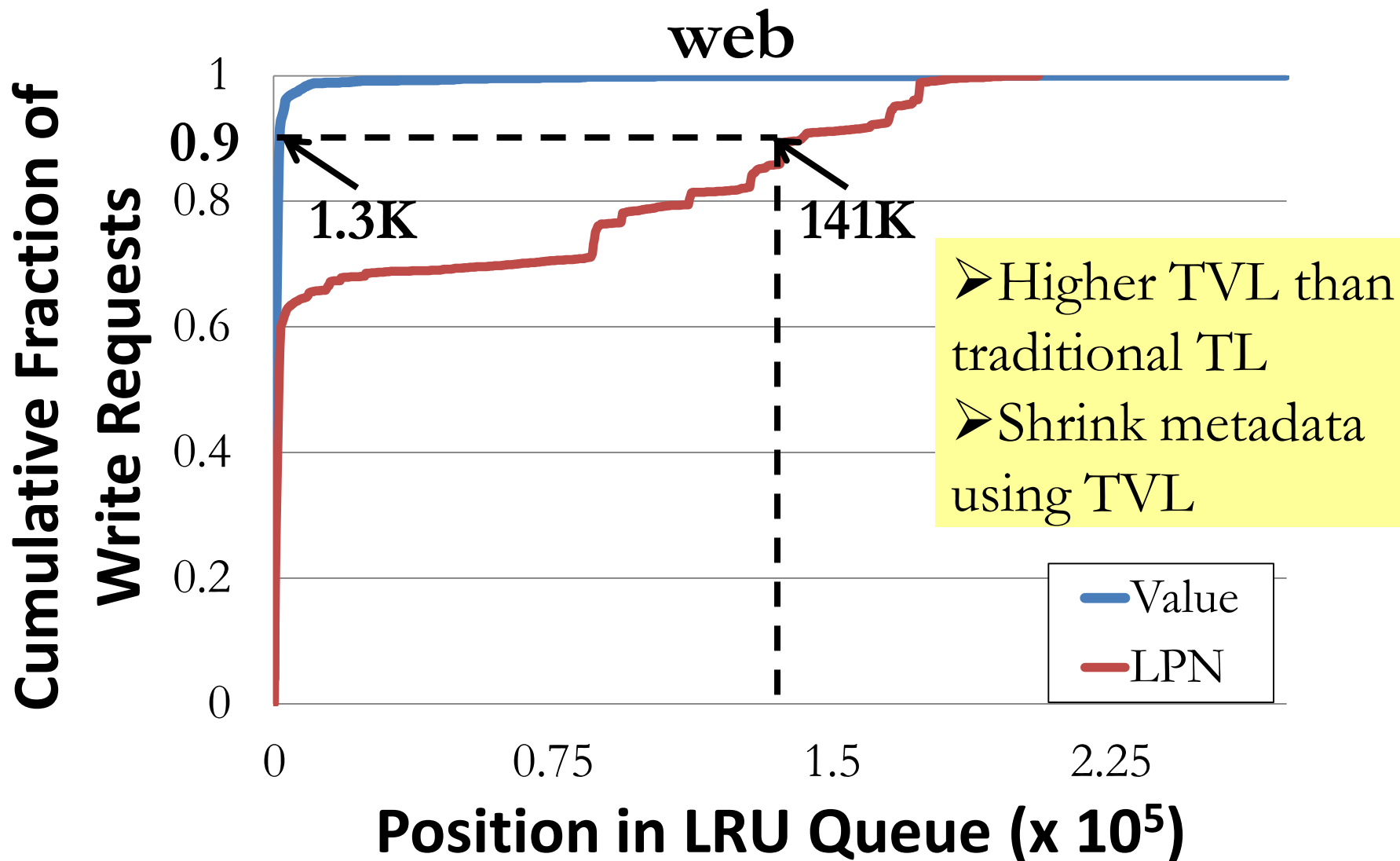
iHPT	
PPN	Hash
P1	H1
P2	H2
P3	H3
P4	H4



Temporal Value Locality

- TVL implies that if a certain *value* is accessed now, it is likely to be accessed again in the near future *not necessarily from the same address*

Temporal Value Locality: Writes



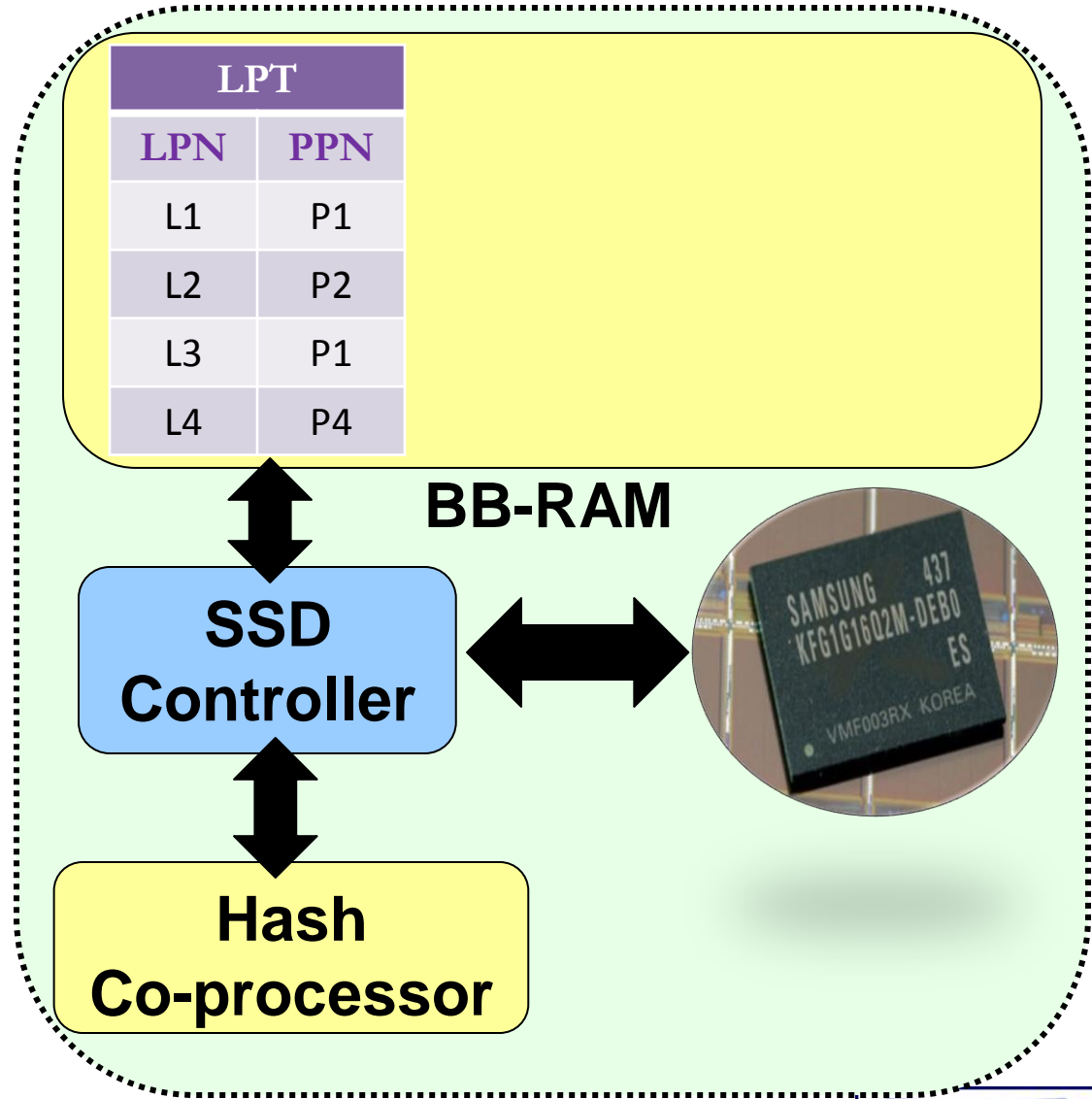
Metadata Management: TVL

HPT	
Hash	PPN
H1	P1
H2	P2
H3	P3
H4	P4

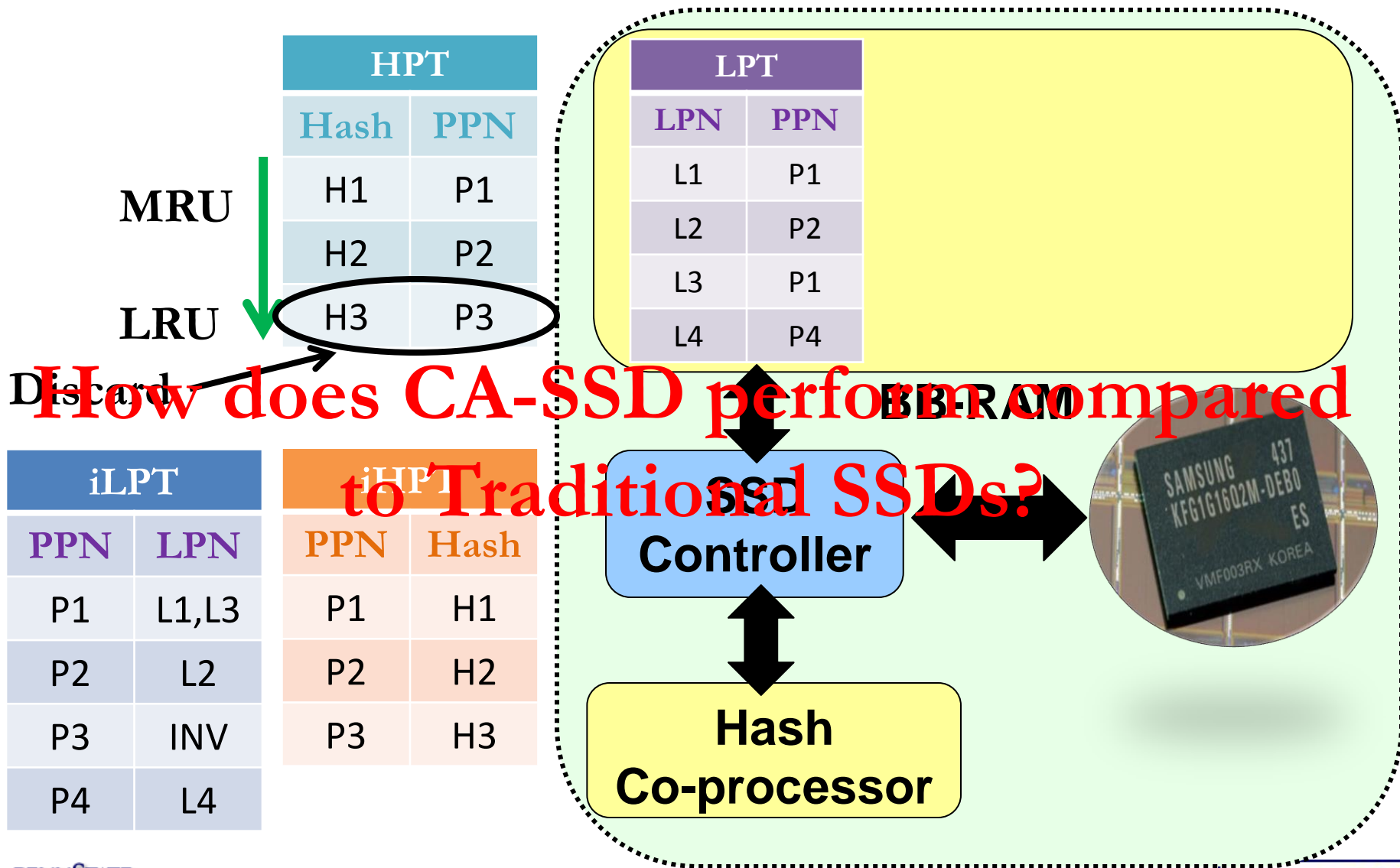
LPT	
LPN	PPN
L1	P1
L2	P2
L3	P1
L4	P4

iLPT	
PPN	LPN
P1	L1,L3
P2	L2
P3	INV
P4	L4

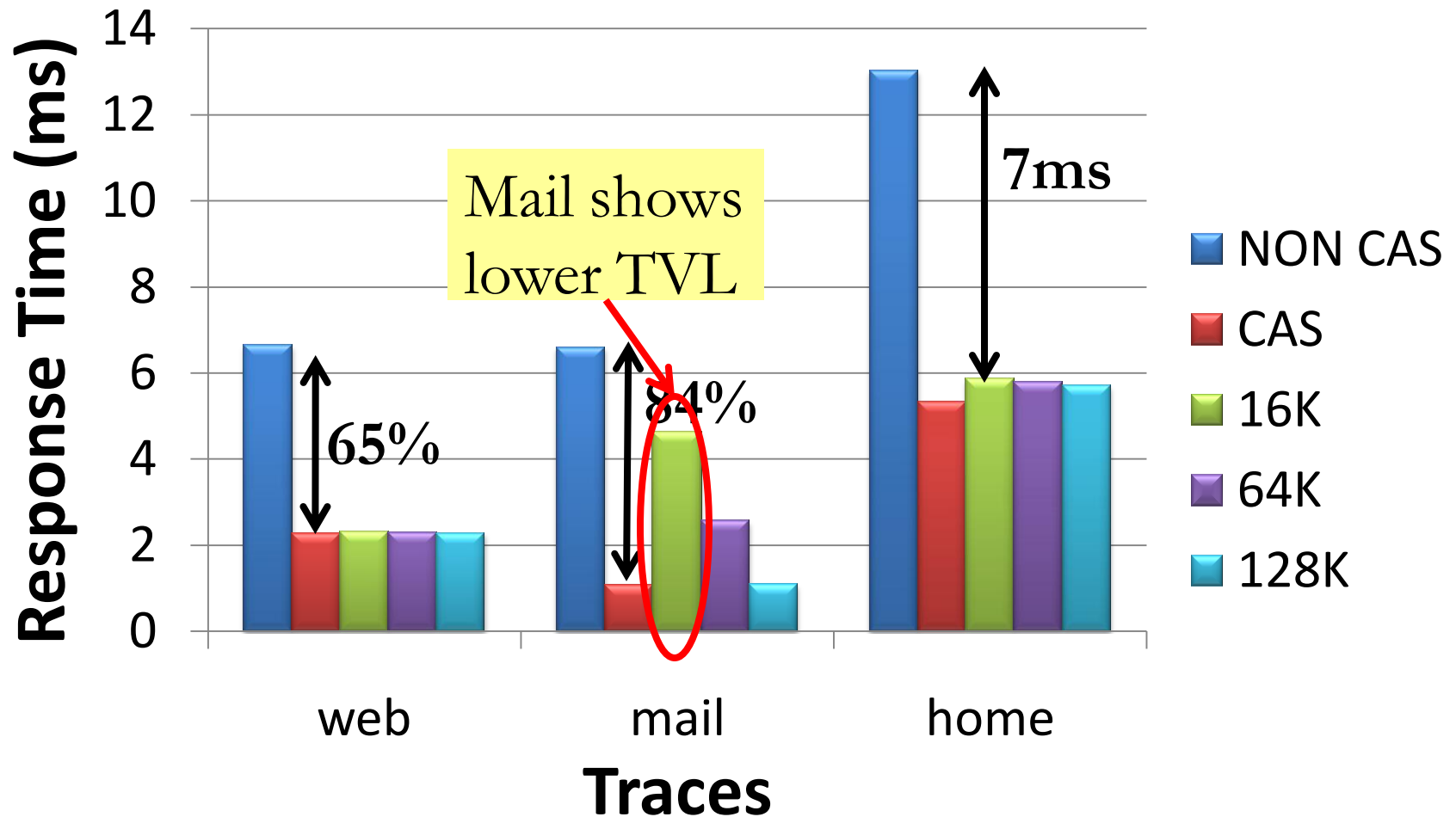
iHPT	
PPN	Hash
P1	H1
P2	H2
P3	H3
P4	H4



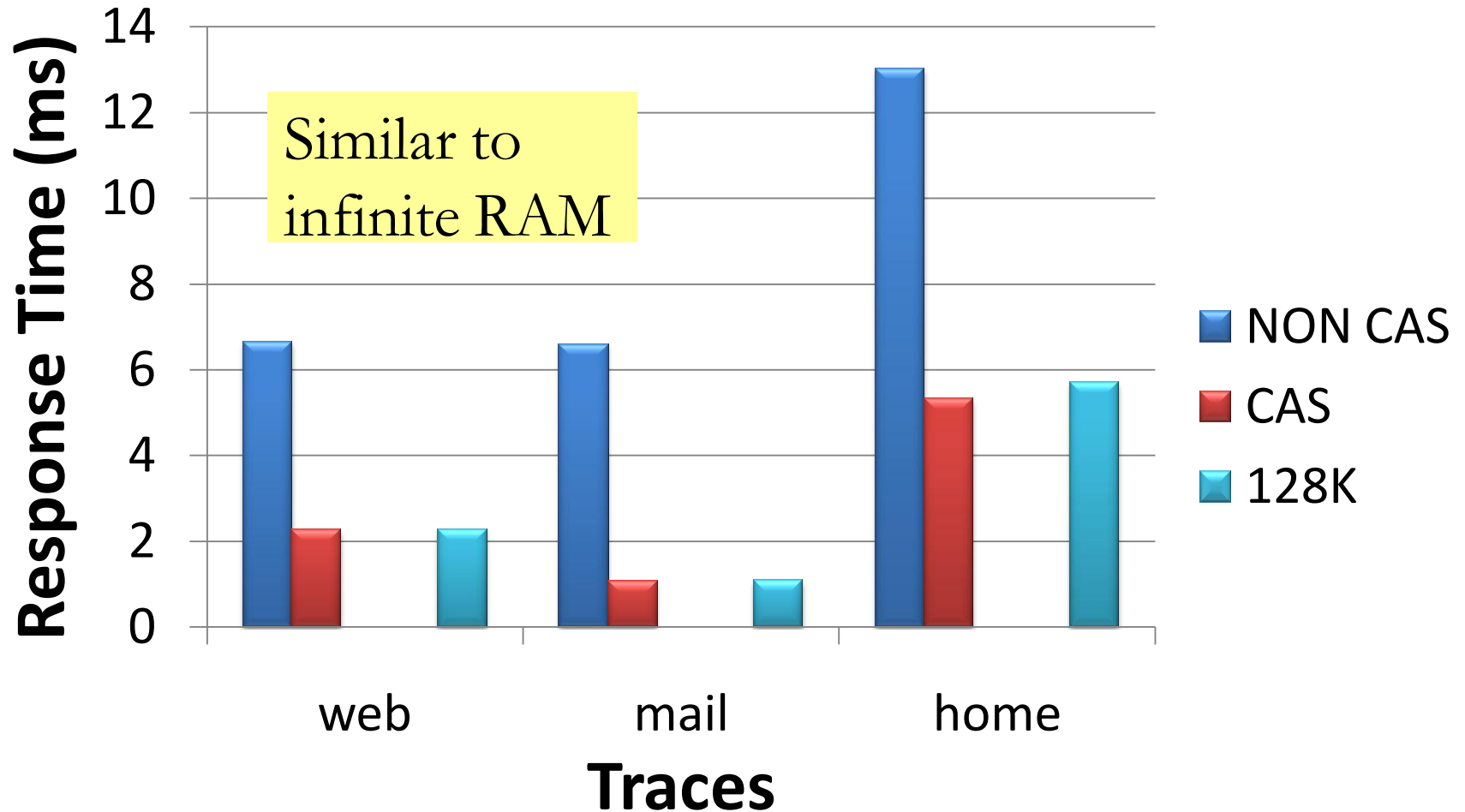
Metadata Management: TVL



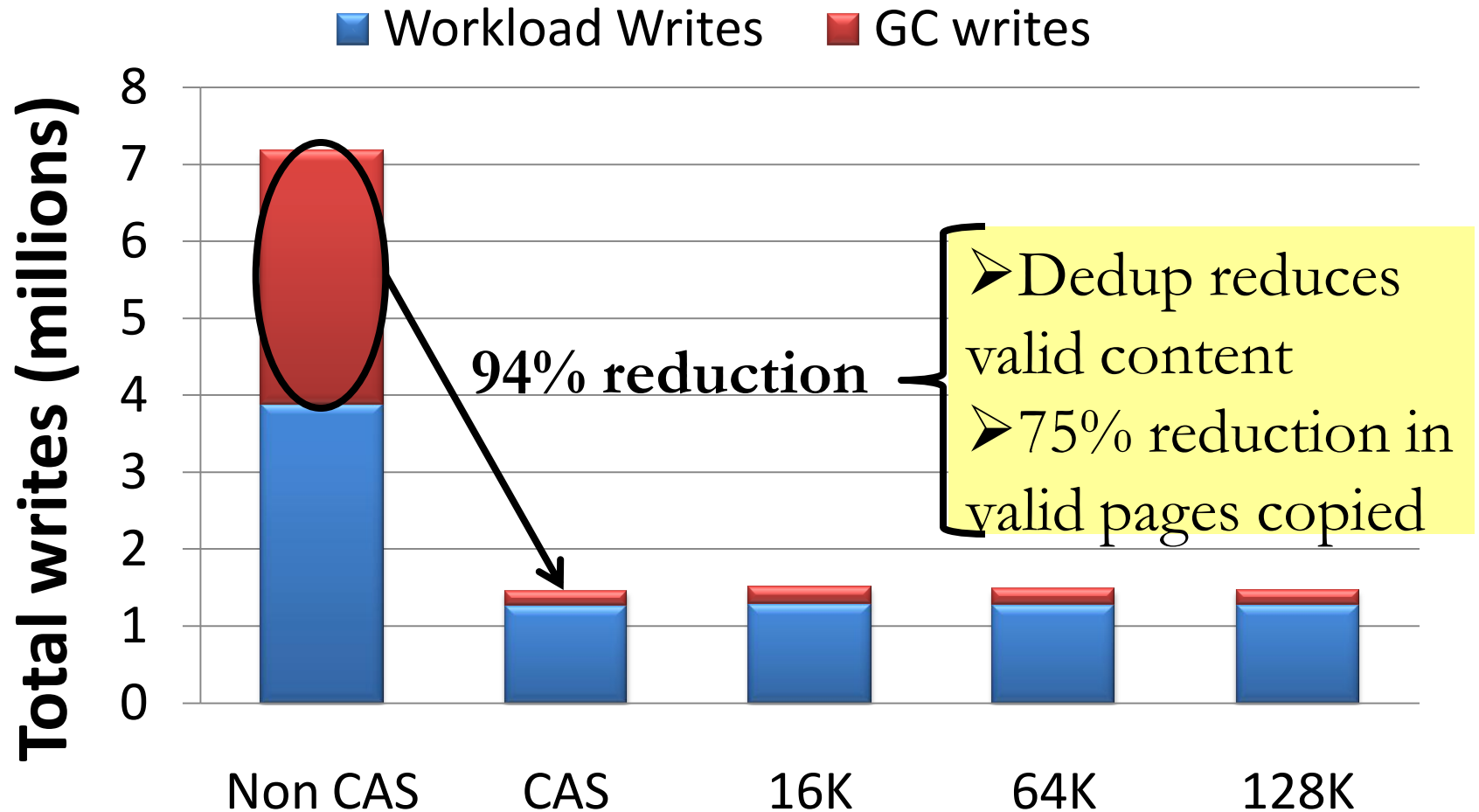
Evaluation : Response Time



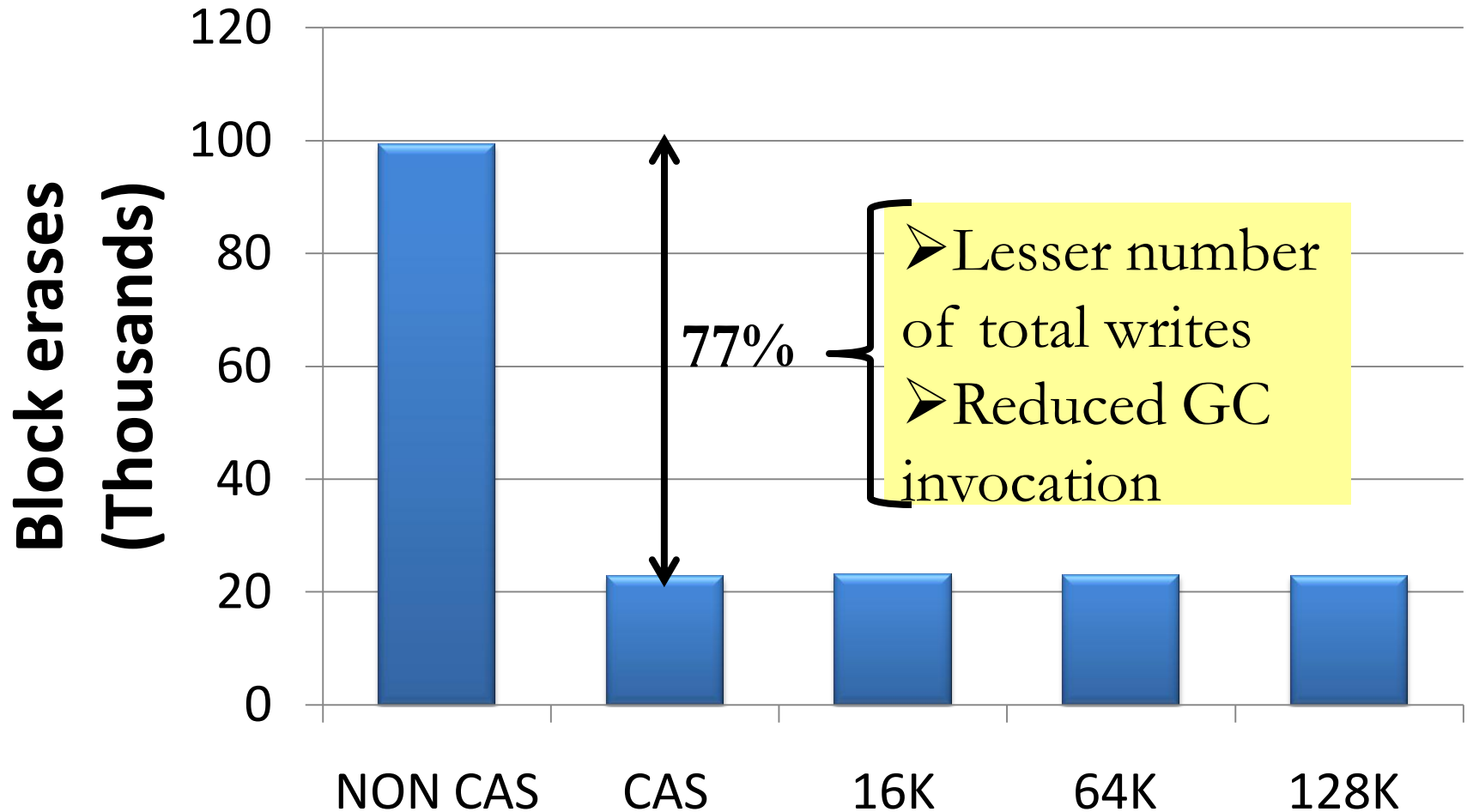
Evaluation : Response Time



Total Writes : *web*



Total Erases : *web*



Conclusions

- **Workloads exhibit significant value locality**
 - **Characterization of Value Popularity and Temporal Value Locality**
- **CAS and SSDs complement each other**
- **Certain implementation challenges need to be addressed**
 - **Mapping structures**
 - **Metadata Management**

Thank You

Questions???