PeerMon: A Peer-to-Peer Network Monitoring System

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Target: General Purpose NWs

Usually single LAN systems

Each machine's resources controlled by local OS

• NFS, but little other system-wide resource sharing

No central scheduler of NW-wide resources

- Users tend to statically pick node(s) to use (ex) write MPI hostfile once, use every time
- Users may not have a choice (ex) ssh cs.swarthmore.edu: target is chosen from static set
- Often large imbalances in NW-wide resource usage

Imbalances Cause Poor Performance

- Swapping on some while lots of free RAM on others
- Large variations in CPU loads
- Variations in contention for NIC, disk, other devices
- Parallel applications (ex. MPI)
 - Usually performance determined by slowest node
 - Picking one overloaded node can result in big performance hit
- Sequential applications
 - Low response rate for interactive jobs
 - Longer execution times for batch jobs

Want to do better load balancing

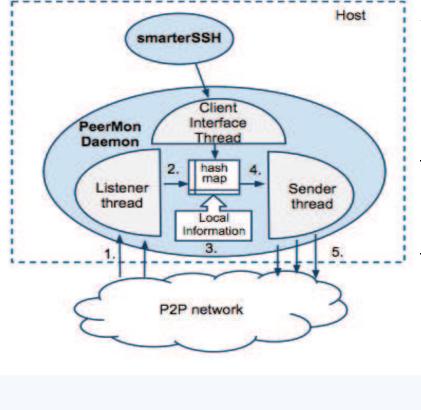
- Tool to easily and quickly discover "good" nodes
 - low CPU load, enough free RAM, fewest number of processes, total # CPUs, ...
 - use to make better job/process placement
 - get better load balancing
 - avoid problems with load imballances
- But has to fit with constraints of target system
 - Still General Purpose system where each OS manages it local node's resources
 - Not implementing a global resource scheduler

PeerMon

- P2P Resource Monitoring System
 - Scalable, fault tolerant, low overhead system
 - No central authority, so no single bottleneck nor single point of failure
 - Each node runs equal peer that provides system-wide resource usage data to local users on its node
 - Fast local access to system-wide resource usage data
- Layered Architecture:
 - PeerMon does the system-wide data collection part
 - Higher-level services use PeerMon data to do load balancing, job placement, ...

PeerMon Architecture

Peer-to-Peer Resource Monitoring System

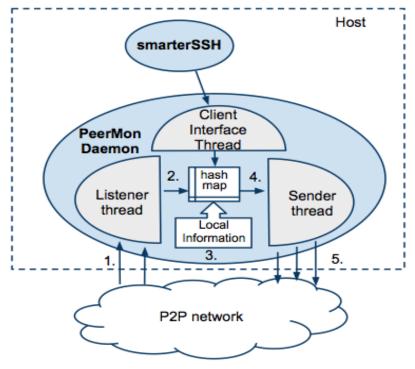


Every node runs equal peer that collects system-wide resource usage data

Sender and Listener Threads: communicate over P2P NW

<u>Client Interface Thread</u>: exports PeerMon data to higher-level services that use it (communicate with local peermon daemon only!)

Listener and Sender Threads



Listener Thread:

- receives resource usage data from other peers
- updates its system-wide resource usage data (stored in hashMap)

Sender Thread:

periodically wakes up & sends its data about whole system to 3 peers

Both use UDP/IP

- Fast, don't need reliable delivery
- Single UDP socket vs. one per connection w/TCP

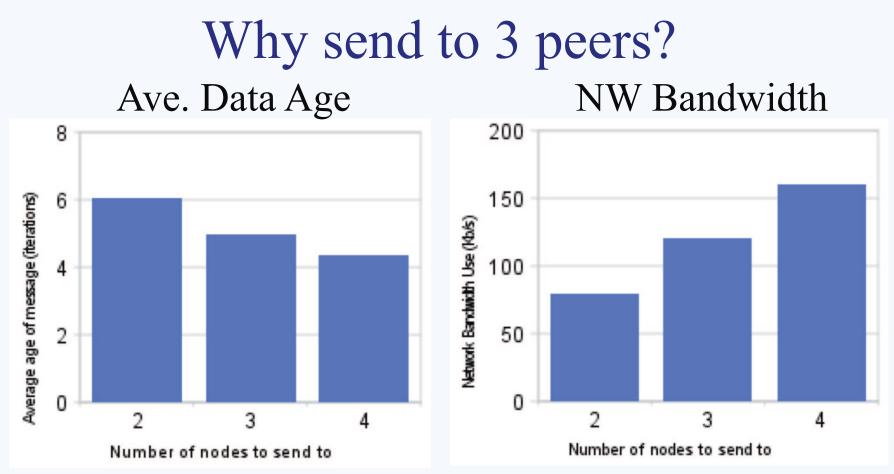
Resource Usage Data

Each PeerMon peer:

- Collects info about its own node
- Sends its full hashMap data to 3 peers
 - Cycle through different heuristics to choose 3 to ensure full conectivity & that new nodes get quickly integrated
- Receives info about other nodes from some of its peers

Constraints on PeerMon Peer's Data:

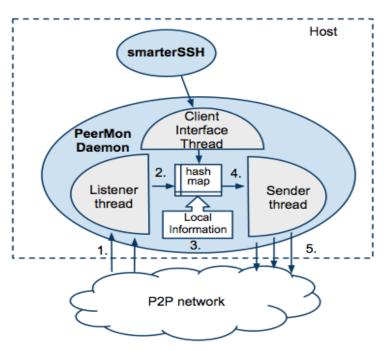
- Doesn't need to be consistent across peers
 - With good messaging heuristics it is close to consistent
- If higher-level service requires an absolute authority, then it can choose 1 PeerMon node to be that authority
 - No different from centralized SNMP systems



Results for a 500 node system

Sending to 3 peers is good trade-off in Data Age vs. NW overheads

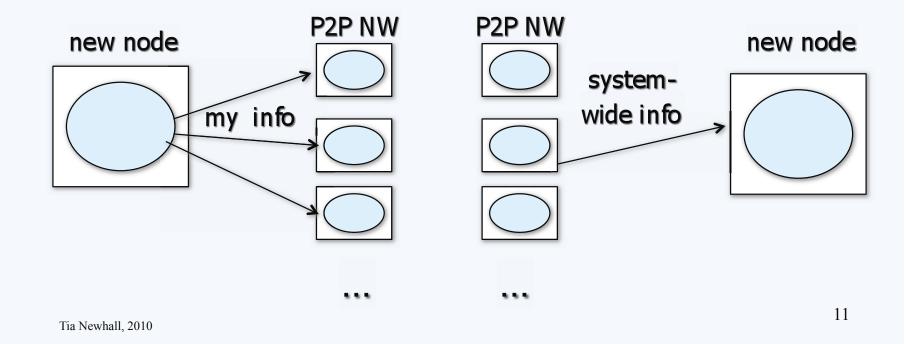
Client Thread



- Local PeerMon daemon provides all system-wide data to local users
 currently TCP interface
- If a higher-level service requires an absolute authority, then it can interact with exactly one PeerMon daemon or implement distributed consensus w/more than one
- For services that don't need absolute agreement, interact with local PeerMon daemon
 => purely distributed interaction

System start-up

New peermon process gets 3 peer IPs config file Sender thread sends data to 3 peers to connect to P2P NW If at least 1 of 3 eventually runs peermon, new node will enter PeerMon P2P NW



Fault Tolerance and Recovery

When a node fails or becomes unreachable, its data ages out of the system

• Users of PeerMon data at other nodes will not choose failed node as one of the "good" nodes

Recovery:

- No different from start-up
- No global state that needs to be reconstructed, new peerMon deamon will enter P2P NW and begin receiving system-wide resource usage data

Example Uses of PeerMon

- <u>SmarterSSH</u>:
 - Uses PeerMon data to pick best ssh target
- <u>autoMPIgen</u>
 - Generates MPI hostfile, choosing best nodes based on PeerMon data
- Dynamic DNS mapping
 - Dynamically binds name to one of current set of best nodes
 - Uses RR in BIND 9 to rotate through set of top N machines periodically updated by PeerMon

SmarterSSH and autoMPIgen

- Simple Python Programs, use PeerMon client TCP interface
- Can order "best" nodes based on CPU load, amount free RAM, or combination of both
- Uses a delta value in ordering nodes so small diffs in load are not significant to ordering
- smarterSSH randomizes the order of "equally" good nodes so subsequent quick invocations distribute ssh load over set of "best" nodes

Example smarterSSH commands

molasses,[~]	,11:29am%	rssh -v -n 5	5 by CPU&RAM
IP	CPU load	free memory	CPU cores
sesame	0.050	7303248	4
turmeric	0.000	6401308	4
molasses	0.000	6302120	4
lime	0.120	6771928	4
myrtle	0.730	9301760	8
molasses,[~] molasses,[~] IP	,11:29am% smarte	rssh -v -n 5 -c ^{top} free memory	
molasses	0.000	6302120	4
	0.000	6401308	4
lettuce		2207396	4
	0.040	2376356	4
sesame		7303248	4
molasses.[~]	.11:29am%		
molasses,[~]	,11:29am% smarte	rssh –c ssh into be	st by CPU load
sshing into	turmeric		w ex ité
		ome/newhall/.ssh/id	d_rsa':
Tia Newhall, 2010			15

How much does PeerMon help?

- Three benchmark programs:
 - 1. Memory Intensive sequential program
 - 2. CPU intensive OpenMP program (single node)
 - 3. RAM&CPU intensive parallel MPI program (ran on 8 of 50 nodes)
- Experiments comparing:
 - Runs on randomly selected node(s): no PeerMon
 - Nodes chosen using PeerMon data with:
 - Ordered by CPU only
 - Ordered by available RAM only
 - Ordered using both CPU load and available RAM

Speed-up of PeerMon vs Random

Node Ranking	Sequential (RAM Intensive)	OpenMP (CPU Intensive)	8 node MPI (Both)
CPU only	0.87	1.63	1.27
RAM only	4.62	2.19	1.78
CPU & RAM	4.62	2.29	1.83

- + Using PeerMon significantly improves performance random only does better when PeerMon ordering criterion is bad match for application
- + Combination of CPU&RAM best ordering criterion

Scalability of PeerMon

- Tested PeerMon NWs of 2-2,200 nodes
- Collected traces of MRTG data for CPU, RAM, NW bandwidth

Results:

- Per node CPU and RAM Usage remains constant
- Per node NW bandwidth grows slightly with size of P2P NW, but still very small
 - Up to .16 Mbit/s for 2,200 node system
 - Each node sends information about every node in NW, so as PeerMon NW grows, so does amt data

Conclusions

- PeerMon: P2P, low overhead, scalable, faulttolerant resource monitoring system for general purpose LANs
- It provides system-wide resource usage data and an interface to export data to higher-level tools and services
- Our example tools that use PeerMon data provide some load balancing in general purpose NW systems and result in significant improvements in performance

Future Work

• Release beta version under GPL we hope before end of summer

www.cs.swarthmore.edu/~newhall/peermon

- Further investigate security & scalability issues
 - PeerMon that spans multiple LANs?
- Implement easier to use client interface
- Add extensibility interface to change set of system resource monitored and how
- Implement more tools that use PeerMon